



CPSC 424

Meshes & Data Structures

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Syllabus

Curves in 2D and 3D

Properties of Curves and Surfaces

Surfaces

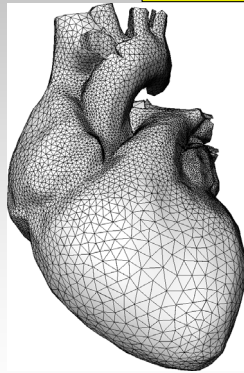
Polygonal meshes

- **Definitions & Data Structures**
- Subdivision (Loop, Butterfly, ...)
- Simplification
- Acquisition
- Analysis & Smoothing
- ...

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Meshes

Triangular mesh – mesh all of whose faces have three vertices



We will discuss connected (3-connected) manifold triangular meshes (closed or open)

Storing mesh data

Uses of mesh data:

- Rendering
- Geometry queries
 - *What are the vertices of face #3?*
 - *Are vertices i and j adjacent?*
 - *Which faces are adjacent to face #7?*
- Geometry operations
 - *Remove/add a vertex/face*



Storing mesh data (cont.)

How “good” is a data structure?

- Time to construct - preprocessing
- Time to answer a query
- Time to perform an operation (update the data structure)
- Space complexity
- Redundancy

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List of faces

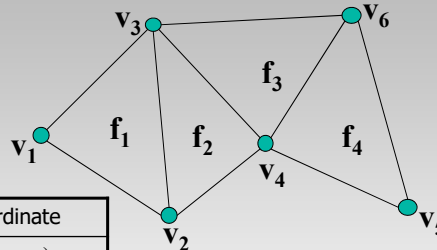
List of vertices (coordinates)

List of faces - triplets of pointers to face vertices (c_1, c_2, c_3)

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List of faces - example



vertex	coordinate
v_1	(x_1, y_1, z_1)
v_2	(x_2, y_2, z_2)
v_3	(x_3, y_3, z_3)
v_4	(x_4, y_4, z_4)
v_5	(x_5, y_5, z_5)
v_6	(x_6, y_6, z_6)

face	vertices (ccw)
f_1	(v_1, v_2, v_3)
f_2	(v_2, v_4, v_3)
f_3	(v_3, v_4, v_6)
f_4	(v_4, v_5, v_6)

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List of faces

Queries:

- What are the vertices of face #3?
 - Answered in $O(1)$ - checking third triplet
- Are vertices i and j adjacent?
 - A pass over all faces is necessary – NOT GOOD

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List of faces – pros and cons

Pros:

- Convenient and efficient (memory wise)
- Can represent non-manifold meshes

Cons:

- Too simple - not enough information on relations between vertices & faces

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Storing mesh data

Uses of mesh data:

- Rendering
- Geometry queries
- Geometry operations

Storage of generic meshes – hard to implement efficiently

Assume: orientable , manifold & triangular

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Half-Edge Data Structure



Record for each face, edge and vertex:

- Geometric information
- Topological information
- Attribute information

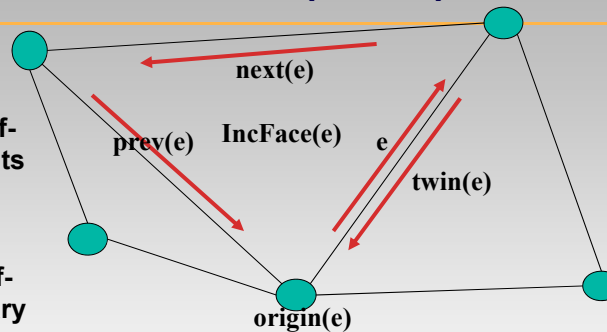
also called DCEL (Doubly-Connected Edge List) or Winged-Edge Structure

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Half-Edge Data Structure (cont.)



- **Vertex record:**
 - Coordinates
 - Pointer to one half-edge that has v as its origin
- **Face record:**
 - Pointer to one half-edge on its boundary



Half-edge record:

- Pointer to its origin, $origin(e)$
- Pointer to its twin half-edge, $twin(e)$
- Pointer to the face it bounds, $IncidentFace(e)$ (face lies to left of e when traversed from origin to destination)
- Next and previous edge on boundary of $IncidentFace(e)$

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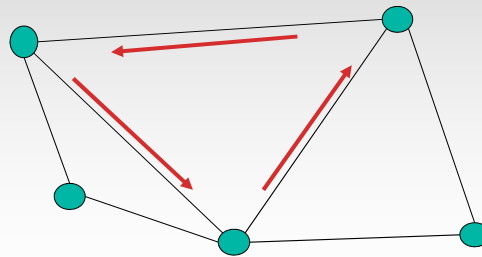
Half-Edge Data Structure (cont.)

Operations supported:

- Walk around boundary of given face
- Visit all edges incident to vertex v

Queries:

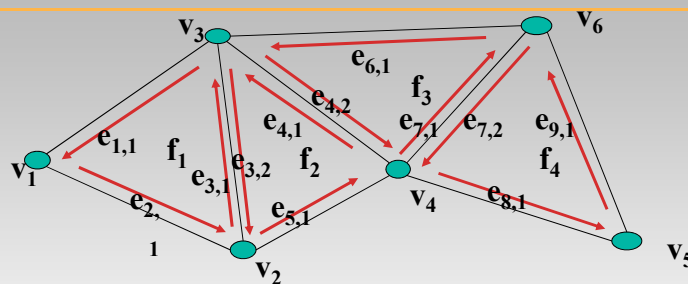
- Most queries are $O(1)$



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Half-Edge Data Structure - example

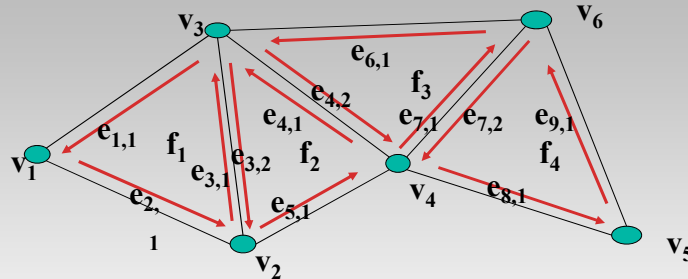


Vertex	coordinate	IncidentEdge
v_1	(x_1, y_1, z_1)	$e_{2,1}$
v_2	(x_2, y_2, z_2)	$e_{5,1}$
v_3	(x_3, y_3, z_3)	$e_{1,1}$
v_4	(x_4, y_4, z_4)	$e_{7,1}$
v_5	(x_5, y_5, z_5)	$e_{9,1}$
v_6	(x_6, y_6, z_6)	$e_{7,2}$

face	edge
f_1	$e_{1,1}$
f_2	$e_{5,1}$
f_3	$e_{4,2}$
f_4	$e_{8,1}$

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Half-Edge – example (cont.)



Half-edge	origin	twin	IncidentFace	next	prev
$e_{3,1}$	v_2	$e_{3,2}$	f_1	$e_{1,1}$	$e_{2,1}$
$e_{3,2}$	v_3	$e_{3,1}$	f_2	$e_{5,1}$	$e_{4,1}$
$e_{4,1}$	v_4	$e_{4,2}$	f_2	$e_{3,2}$	$e_{5,1}$
$e_{4,2}$	v_3	$e_{4,1}$	f_3	$e_{7,1}$	$e_{6,1}$

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Half-Edge – pros and cons

Pros:

- All queries in $O(1)$ time
- All operations are $O(1)$ (usually)

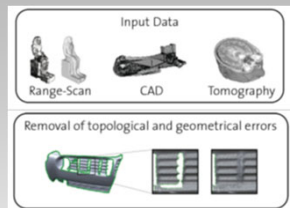
Cons:

- Represents only manifold meshes

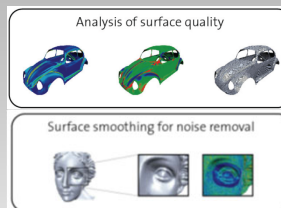
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Processing Pipeline



Acquisition (& repair)



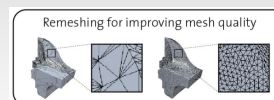
Analysis & Smoothing



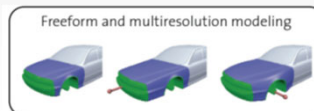
Simplification



Parameterization



Remeshing



Modeling

[Botch, 2006]

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