

In this lecture we:

- Discussed an efficient algorithm to find the Longest Increasing Subsequence (*LIS*);
- And how to solve the Longest Common Subsequence (*LCS*) using *LIS*;

Handouts (posted on webpage):

- Group-work exercises (Wednesday March 1)

Reading: ASSIGNED READING(S).

## 1 Dynamic Programming

### 1.1 Longest Increasing Subsequence (*LIS*)

Having  $R = \{5, 3, 4, 9, 6, 2, 1, 8\}$  the Longest Increasing Subsequence is:  $\{3, 4, 6, 8\}$

$$R = \{5, \textcircled{3}, \textcircled{4}, 9, \textcircled{6}, 2, 1, \textcircled{8}\}$$


Using Longest Common Subsequence (*LCS*):

1. Sort  $R$   $\mathcal{O}(n \log n)$
2. Report  $LCS(\underbrace{R}_n, \underbrace{\text{sort}(R)}_n)$   $\mathcal{O}(n^2)$


Faster solution to *LIS*

When finding *LIS* of  $R\{1\dots k\}$  what information about  $R\{1\dots k-1\}$  would be useful?

A) *LIS*  $R\{1\dots k-1\}$

Not enough because  $R\{\overbrace{5, 3, 4, 9, 6, 2, 1}^{349 \text{ vs } 346}, 8\}$   


B) Best *LIS*  $R\{1\dots k-1\}$  (one with smallest last value)

Not enough because  $R\{1, 2, \overbrace{5}^{\uparrow}, 3, 4\}$   


C) Best Increasing Subsequence (*BIS*) of lengths  $\{1, 2, 3 \dots j\}$

Example step by step:

1.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{1\}$$

$$BIS[2] = \{3, 4\}$$

$$BIS[3] = \{3, 4, 6\}$$

$$BIS[4] = \{\emptyset\}$$

2.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{1\}$$

$$BIS[2] = \{3, 4\}$$

$$BIS[3] = \{3, 4, 5\}$$

$$BIS[4] = \{\emptyset\}$$

3.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{1\}$$

$$BIS[2] = \{3, 4\}$$

$$BIS[3] = \{3, 4, 5\}$$

$$BIS[4] = \{3, 4, 5, 7\}$$

This is probably the solution but, how can we compute this in less time?

The solution is to avoid writing everything let's see an example step by step:

1.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{8\}$$

$$BIS[2] = \{\emptyset\}$$

$$BIS[3] = \{\emptyset\}$$

$$BIS[4] = \{\emptyset\}$$

2.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{8, 3\}$$

$$BIS[2] = \{\emptyset\}$$

$$BIS[3] = \{\emptyset\}$$

$$BIS[4] = \{\emptyset\}$$

3.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{8, 3\}$$



$$BIS[2] = \{4\}$$

$$BIS[3] = \{\emptyset\}$$

$$BIS[4] = \{\emptyset\}$$

4.  $R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$



$$BIS[1] = \{8, 3\}$$



$$BIS[2] = \{4\}$$



$$BIS[3] = \{9\}$$

$$BIS[4] = \{\emptyset\}$$

$$5. R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$$

$$BIS[1] = \{8, \textcircled{3}\}$$

$$BIS[2] = \{\textcircled{4}\}$$

$$BIS[3] = \{\emptyset, \textcircled{6}\}$$

$$BIS[4] = \{\emptyset\}$$

$$6. R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$$

$$BIS[1] = \{8, \textcircled{3}, 2\}$$

$$BIS[2] = \{\textcircled{4}\}$$

$$BIS[3] = \{\emptyset, \textcircled{6}\}$$

$$BIS[4] = \{\emptyset\}$$

$$7. R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$$

$$BIS[1] = \{8, \textcircled{3}, 2, 1\}$$

$$BIS[2] = \{\textcircled{4}\}$$

$$BIS[3] = \{\emptyset, \textcircled{6}\}$$

$$BIS[4] = \{\emptyset\}$$

$$8. R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$$

$$BIS[1] = \{8, \textcircled{3}, 2, 1\}$$

$$BIS[2] = \{\textcircled{4}\}$$

$$BIS[3] = \{\emptyset, 6, \textcircled{5}\}$$

$$BIS[4] = \{\emptyset\}$$

$$9. R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$$

$$BIS[1] = \{8, \textcircled{3}, 2, 1\}$$

$$BIS[2] = \{\textcircled{4}\}$$

$$BIS[3] = \{\cancel{9}, \cancel{6}, \textcircled{5}\}$$

$$BIS[4] = \{\textcircled{7}\}$$

$$10. R\{8, 3, 4, 9, 6, 2, 1, 5, 7, 6\}$$

$$BIS[1] = \overbrace{\{8, \textcircled{3}, 2, 1\}}^{\text{decreasing order}}$$

$$BIS[2] = \{\textcircled{4}\}$$

$$BIS[3] = \{\cancel{9}, \cancel{6}, \textcircled{5}\}$$

$$BIS[4] = \{\cancel{7}, \textcircled{6}\}$$

To add the next number in  $R$  perform binary search on the last values in the  $BIS$ s and put the next number in sorted order, pointing to the end value of the previous  $BIS$  (black arrows) having a running time of  $\mathcal{O}(n \log n)$

## 1.2 Use LIS to solve LCS

To find a sequence of index pairs of matches so that sequence increases in both coordinates.

Example:

$$\{A, \textcircled{B}, \textcircled{C}, B, \textcircled{A}, \textcircled{C}, \textcircled{C}, B\}$$

$$\{\textcircled{B}, \textcircled{C}, D, \textcircled{A}, B, \textcircled{C}, \textcircled{C}\}$$

We can see the relation between the sequence in table 1

Table 1: *LCS* pairs

		1	2	3	4	5	6	7	8
		A	B	C	B	A	C	C	B
1	B		X		X				X
2	C			X			X	X	
3	D								
4	A	X				X			
5	B		X		X				X
6	C			X			X	X	
7	C			X			X	X	

In this case, we can see that the coordinates of the *LCS* pairs have increasing order in both coordinates.

$$\overbrace{\{(1, 2), (2, 3), (4, 5), (6, 6), (7, 7)\}}^{\text{increasing order}}$$

It is easy to obtain the matching pairs, and to list them by its row and column coordinates:

$$\begin{aligned} \text{Rows} &= \{1, 1, 1, 2, 2, 2\dots\} \\ \text{Columns} &= \{2, 4, 8, 3, 6, 7\dots\} \end{aligned}$$

The trick is to invert the column numbers that have the same value in row.

$$\begin{aligned} \text{Rows} &= \{1, 1, 1, 2, 2, 2\dots\} \\ \text{Columns} &= \{\underbrace{8, 4, 2}_{\text{inverted}}, \underbrace{7, 6, 3}_{\text{inverted}} \dots\} \end{aligned}$$

List coordinates of matching characters in order by row index and, within one row, in reverse order by column number, then run *LIS* on the sequence of columns order.

Another way to see this problem is to determine a value of 1 to vertical and horizontal movements, and a value of 0 to diagonal movements like in a chessboard, having this, the problem can be seen as the shortest path between the upper left corner and the lower right corner as seen in table 2.

Table 2: *LCS* as Shortest Path

		1	2	3	4	5	6	7	8
		A	B	C	B	A	C	C	B
1	B		(X)		X				X
2	C			(X)			X	X	
3	D								
4	A	X				(X)			
5	B		X		X				X
6	C			X			(X)	X	
7	C			X			X	(X)	