CPSC 221: Data Structures
Dictionary ADT
Hashing

Alan J. Hu

(Using mainly Steve Wolfman's Old Slides)

Learning Goals

After this unit, you should be able to:

- Define various forms of the pigeonhole principle; recognize and solve the specific types of counting and hashing problems to which they apply.
- Provide examples of the types of problems that can benefit from a hash data structure.
- Compare and contrast open addressing and chaining.
- Evaluate collision resolution policies.
- Describe the conditions under which hashing can degenerate from O(1) expected complexity to O(n).
- Identify the types of search problems that do not benefit from hashing (e.g. range searching) and explain why.
- Manipulate data in hash structures both irrespective of implementation and also within a given implementation.

Outline

- Dictionary ADT
- Hash Table Overview
- Hash Functions
- Collisions and the Pigeonhole Principle
- Collision Resolution:
  - Chaining
  - Open-Addressing
- Deletion and Rehashing

Dictionary ADT

- Dictionary operations
  - create
  - destroy
  - insert
  - find
  - delete

insert

- brownies
  - tasty

- midterm
  - would be tastier with brownies
- prog-project
  - so painful… who invented templates?
- wolf
  - the perfect mix of oomph and Scrabble value

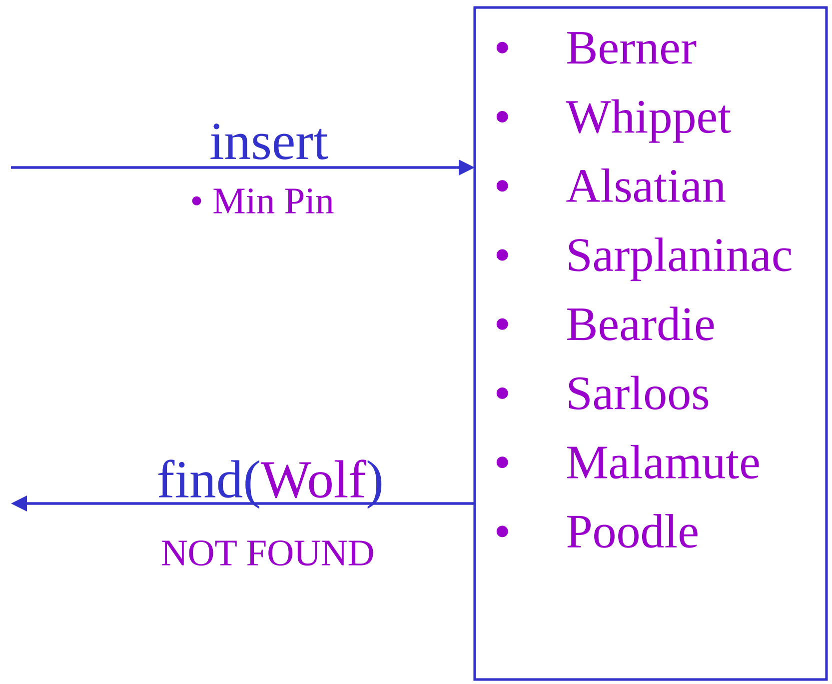
find(wolf)

- wolf
  - the perfect mix of oomph and Scrabble value

- Stores *values* associated with user-specified *keys*
  - values may be any (homogenous) type
  - keys may be any (homogenous) comparable type

Search/Set ADT

- Dictionary operations
  - create
  - destroy
  - insert
  - find
  - delete



- Stores keys
  - keys may be any (homogenous) comparable
  - quickly tests for membership

A Modest Few Uses

- Arrays and “Associative” Arrays
- Sets
- Dictionaries
- Router tables
- Page tables
- Symbol tables
- C++ Structures
- Python’s __dict__ that stores fields/methods

Naïve Implementations

	insert	find	delete
• Linked list			
• Unsorted array			
• Sorted array			

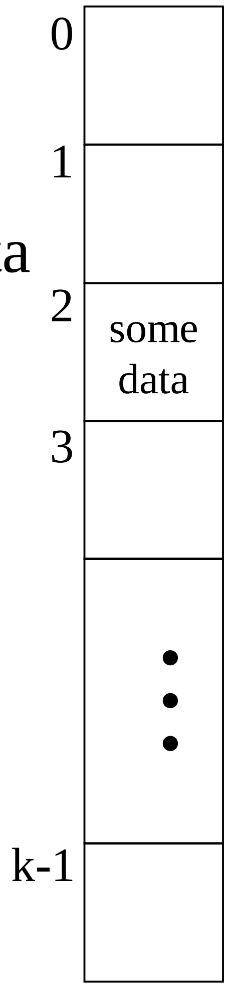
Desiderata

- Fast insertion
  - runtime:
- Fast searching
  - runtime:
- Fast deletion
  - runtime:

Hash Table Goal

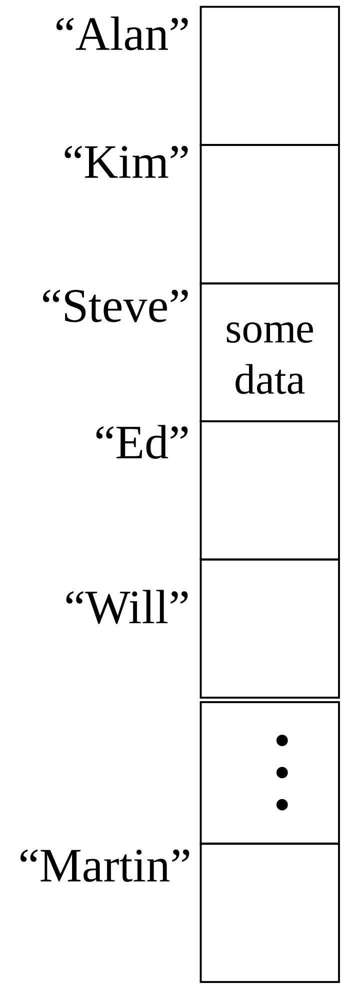
We can do:

a[2] = some data



We want to do:

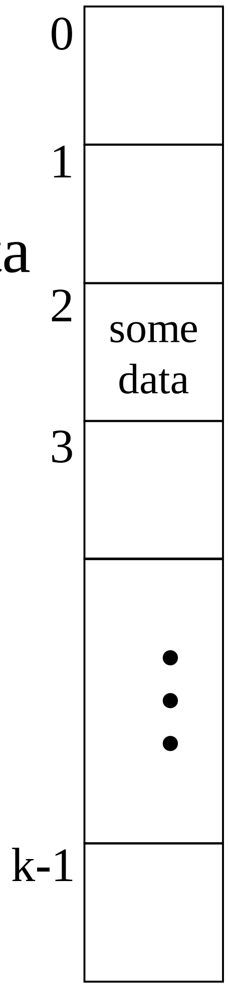
a[“Steve”] = some data



Aside: How do arrays do that?

We can do:

a[2] = some data



Q: If I know houses on a certain block in Vancouver are on 33-foot-wide lots, where is the 5th house?
A: It's from (5-1)*33 to 5*33 feet from the start of the block.

element_type a[SIZE];

Q: Where is a[i]?
A: start of a + i*sizeof(element_type)

Aside: This is why array elements have to be the same size, and why we start the indices from 0.

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  - Chaining
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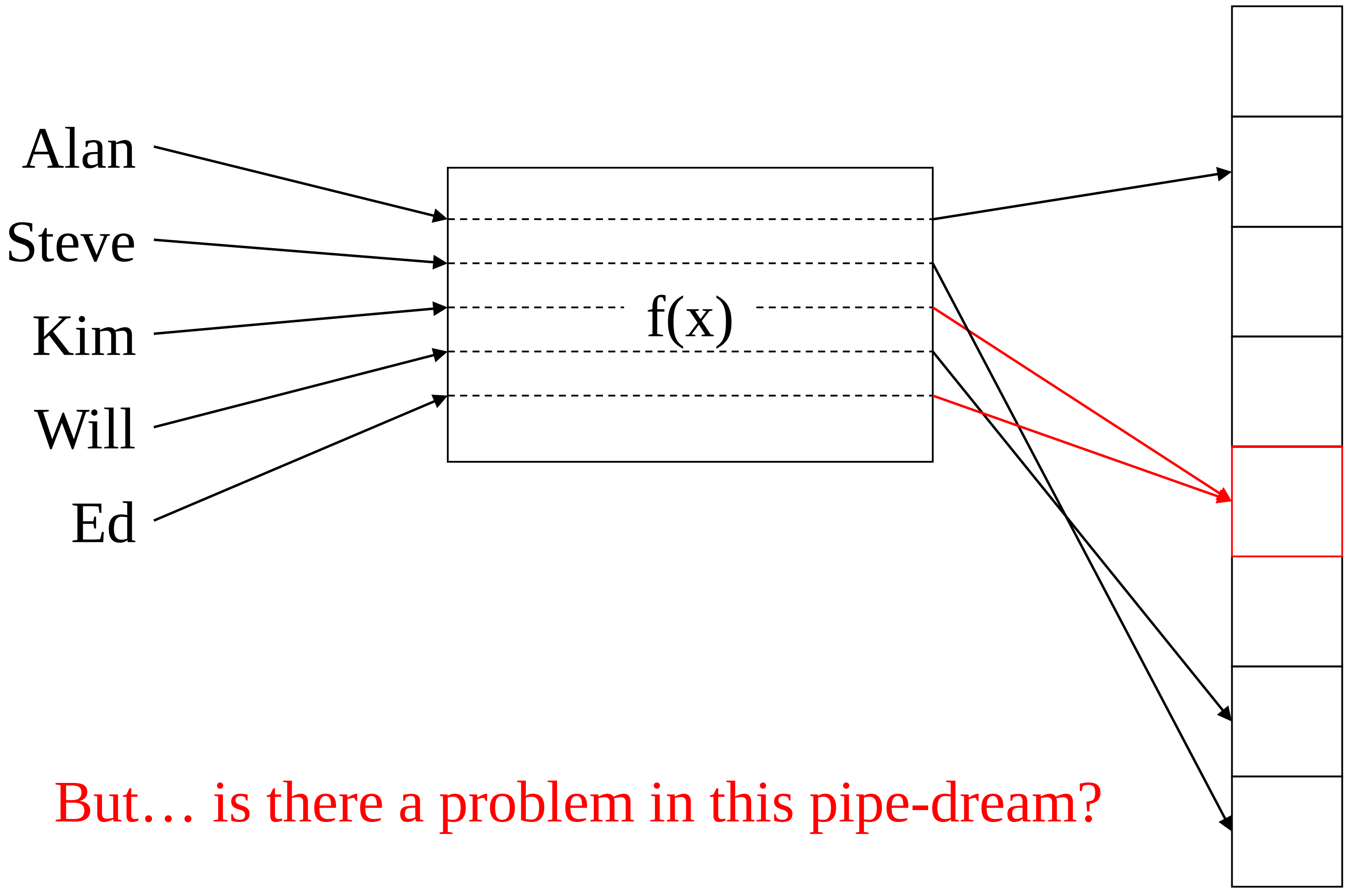
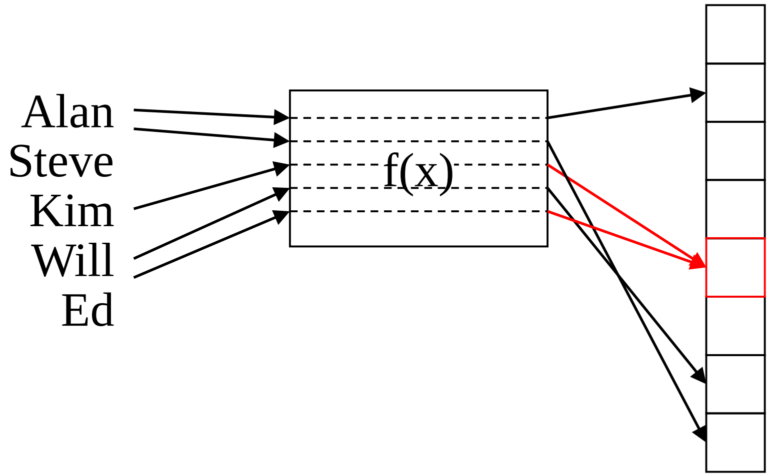
Hash Table Approach



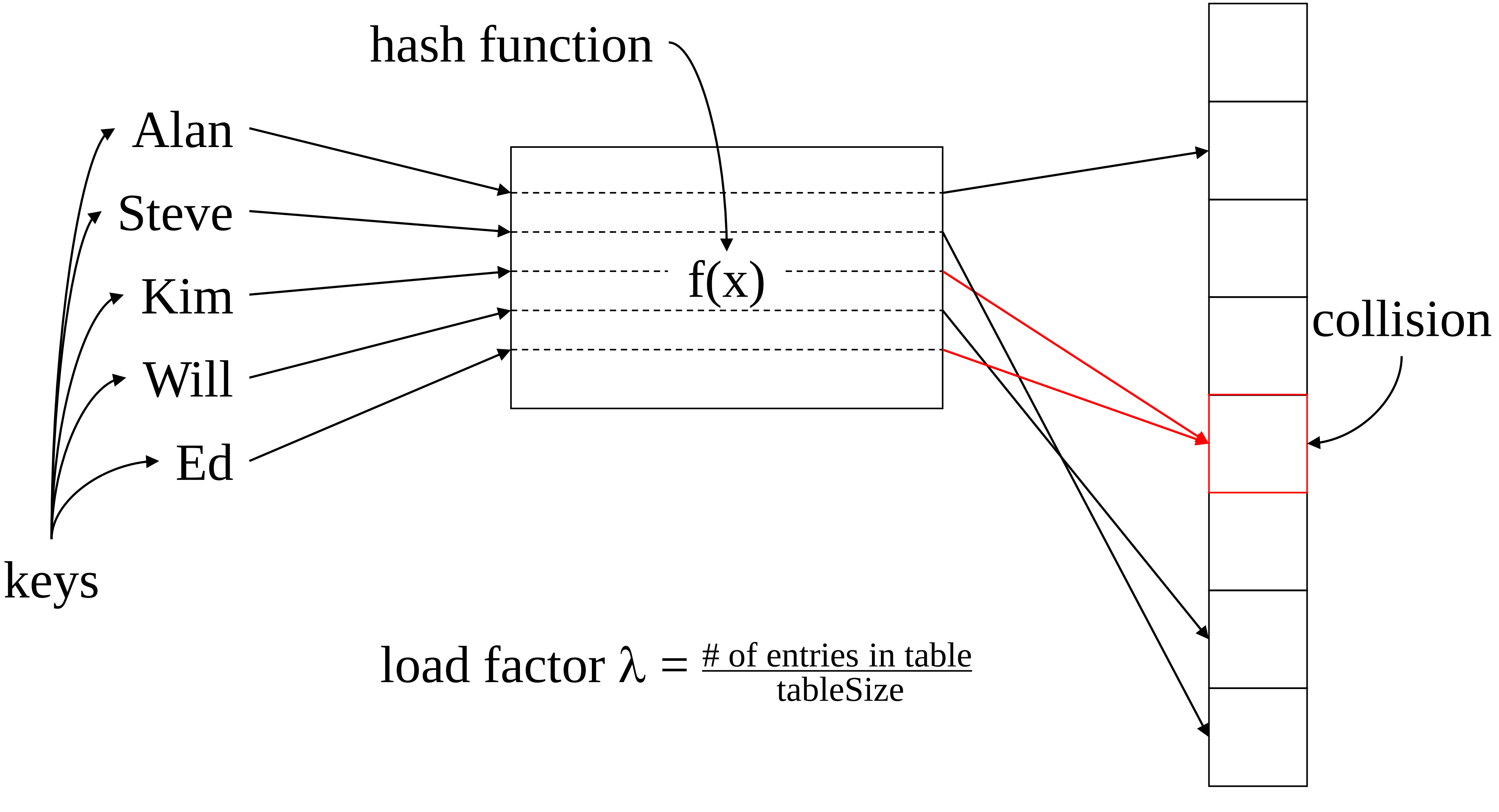
But… is there a problem in this pipe-dream?

Hash Table
Dictionary Data Structure

- Hash function: maps keys to integers
  - result: can quickly find the right spot for a given entry
- Unordered and sparse table
  - result: cannot efficiently list all entries, ***definitely*** cannot efficiently list all entries in order or list entries between one value and another (a “range” query)



Hash Table Terminology



load factor $\lambda = \frac{\text{\# of entries in table}}{\text{tableSize}}$

Hash Table Code
First Pass

```
Value & find(Key & key) {
  int index = hash(key) % tableSize;
  return Table[index];
}
```

What should the hash function be?

How should we resolve collisions?

What should the table size be?

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- Constant-Time Dictionaries?
- Hash Table Overview
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A Good (Perfect?) Hash Function…

…is easy (fast) to compute (O(1) *and* fast in practice).

…distributes the data evenly (hash(a) % size ≠ hash(b) % size).

…uses the whole hash table (for all 0 ≤ k < size, there's an i such that hash(i) % size = k).

Aside: a Bit of 121 Theory

…is easy (fast) to compute (O(1) *and* fast in practice).

…distributes the data evenly (hash(a) % size ≠ hash(b) % size).

…uses the whole hash table (for all $0 \leq k <$ size, there's an i such that hash(i) % size = k).

Onto (surjective)

Ideally, one-to-one (injective)

Good Hash Function for Integers

- Choose
  - tableSize is prime
  - hash(n) = n
- Example:
  - tableSize = 7

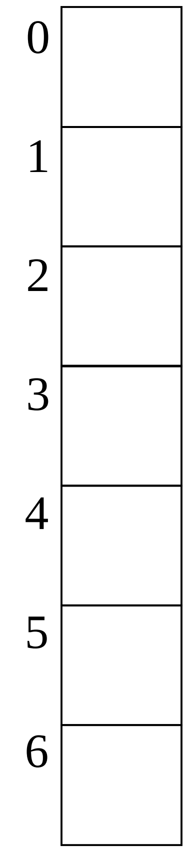
insert(4)

insert(17)

find(12)

insert(9)

delete(17)



Good Hash Function for Strings?

- Let $s = s_0 s_1 s_2 s_3 \ldots s_{n-1}$: choose
  - $hash(s) = s_0 + s_1 31 + s_2 31^2 + s_3 31^3 + \ldots + s_{n-1} 31^{n-1}$

Think of the string as a base 31 number.

- Problems:
  - hash(“really, really big”) = well… something really, really big
  - hash(“one thing”) % 31 = hash(“other thing”) % 31

Why 31? It’s prime. It’s **not** a power of 2. It works pretty well.

Making the String Hash Easy to Compute

- Use Horner's Rule

```
int hash(String s) {
  h = 0;
  for (i = s.length() - 1; i >= 0; i--) {
    h = (s_i + 31*h) % tableSize;
  }
  return h;
}
```

Making the String Hash Cause Few Conflicts

- Ideas?

Making the String Hash Cause Few Conflicts

- Ideas?

    Make sure tableSize is not a multiple of 31.

Hash Function Summary

- Goals of a hash function
  - reproducible mapping from key to table entry
  - evenly distribute keys across the table
  - separate commonly occurring keys (neighboring keys?)
  - complete quickly
- Sample hash functions:
  - h(n) = n % size
  - h(n) = string as base 31 number % size
  - *Multiplicative Hash: multiply key by a constant*
  - *Universal Hashing: functions with random parameters*
  - *Cryptographically Secure Hashing (e.g., MD5, SHA-1, etc.)*

How to Design a Hash Function

- Know what your keys are *or*
- Study how your keys are distributed.
- Try to include all important information in a key in the construction of its hash.
- Try to make “neighboring” keys hash to very different places.
- Prune the features used to create the hash until it runs “fast enough” (application dependent).

How to Design a Hash Function

- Know what your keys are *or*

In real life, use a standard hash function that people have already shown works well in practice!

different places.

- Prune the features used to create the hash until it runs “fast enough” (application dependent).

Extra Slides:
Some Other Hashing Methods

Good Hashing: Multiplication Method

- Hash function is defined by some positive number *A*

  $h_A(k) = (A * k) \% size$

- Example: A = 7, size = 10

  $h_A(50) = 7*50 \mod 10 = 350 \mod 10 = 0$

  - choose A to be relatively prime to size
  - more computationally intensive than a single mod
  - (This is simplified from a more general, theoretical case.)

Universal Hash Functions

- A family of hash functions is called universal if the probability that hash(x)=hash(y) is at most 1/size, if hash is chosen randomly from the family.

- (There are even stronger properties of families of hash functions that are sometimes useful, e.g., that the difference hash(x)-hash(y) is a uniform random variable, etc.)

Good Hashing:
A Universal Hash Function

- Parameterized by p, a, and b:
  - p is a big prime
  - a and b are arbitrary integers in [1,p-1]

$H_{p,a,b}(x) = \quad (a \cdot x + b) \bmod p$

(If p is the table size, this is universal. If you mod the result by a smaller table size (a small fraction of p), it's almost universal.)

Good Hashing: Bit-Level Universal Hash Function

- If table size is 2b, and your keys are r bits long, this is a good universal hash function:
  - Choose a random b-by-r 0/1 matrix A.
  - Compute hash(x) = Ax

$$Ax = \begin{bmatrix} 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} 1 \\ 0 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = hash(x)$$

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The Pigeonhole Principle (informal)

You can’t put k+1 pigeons into k holes without putting two pigeons in the same hole.

Image by en:User:McKay,
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Collisions

- *Pigeonhole principle* says we can't avoid all collisions
  - try to hash without collision m keys into n slots with $m > n$
  - try to put 6 pigeons into 5 holes

Collisions

- *Pigeonhole principle* says we can't avoid all collisions
  - try to hash without collision m keys into n slots with $m > n$
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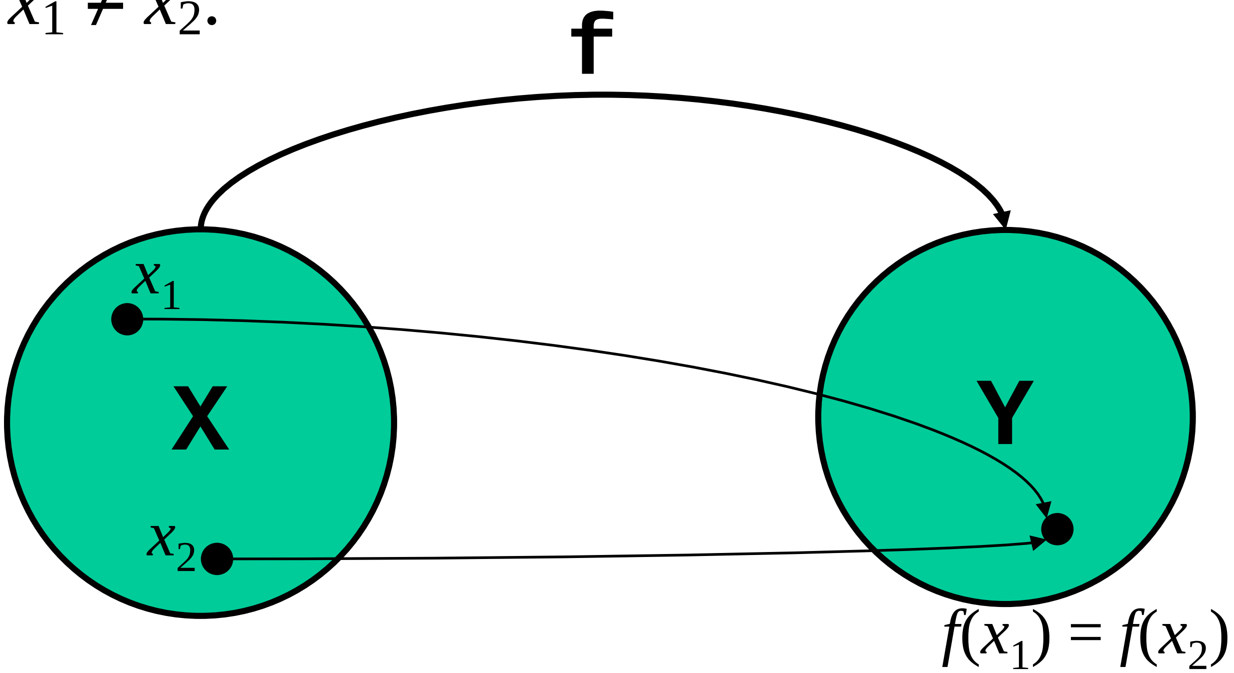
Alan's Aside: This is actually somewhat misleading.
Collisions are a problem even when m < n.
So this tie-in of collisions and the pigeonhole principle isn't really fundamental. It's just a nice chance to introduce the pigeonhole principle…

The Pigeonhole Principle (formal)

Let X and Y be finite sets where $|X| > |Y|$.

If $f : X \rightarrow Y$, then $f(x_1) = f(x_2)$ for some x_1, x_2 in X, where $x_1 \neq x_2$.



The Pigeonhole Principle (Example #1)

Suppose we have 5 colours of Halloween candy, and that there's lots of candy in a bag. How many pieces of candy do we have to pull out of the bag if we want to be sure to get 2 of the same colour?

a. 2

b. 4

c. 6

d. 8

e. None of these

The Pigeonhole Principle (?) (Example #2)

If there are 1000 pieces of each colour, how many do we need to pull to guarantee that we'll get 2 *black* pieces of candy (assuming that black is one of the 5 colours)?

a. 2
b. 6
c. 4002
d. 5001
e. None of these

The Pigeonhole Principle (No!) (Example #2)

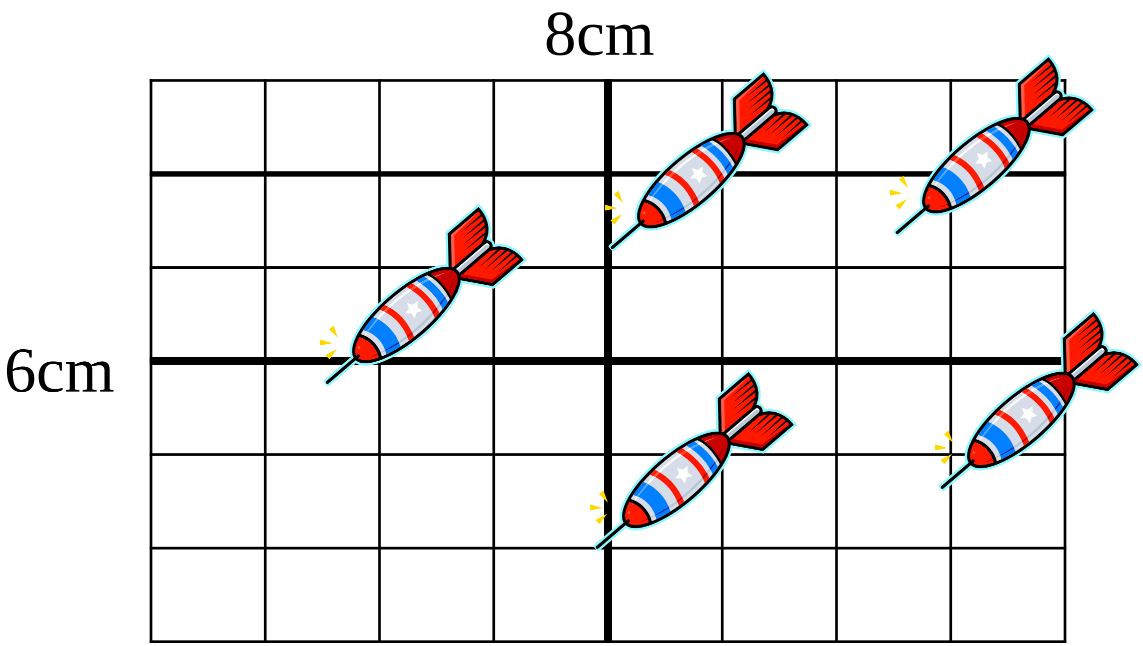
If there are 1000 pieces of each colour, how many do we need to pull to guarantee that we'll get 2 *black* pieces of candy (assuming that black is one of the 5 colours)?

a. 2
b. 6
c. 4002
d. 5001
e. None of these

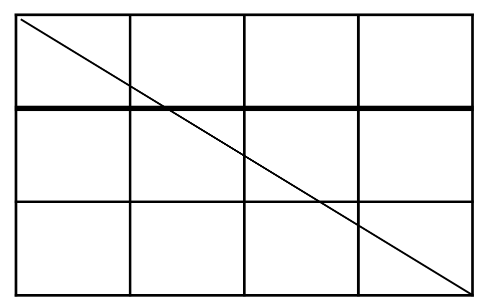
The PhP doesn't tell us *which* hole has two pigeons.

The Pigeonhole Principle (Example #3)

If 5 points are placed in a 6cm x 8cm rectangle, argue that there are two points that are not more than 5 cm apart.



Hint: How long is the diagonal?



The Pigeonhole Principle (Example #4)

For integers a, b, we write *a divides b* as $a|b$, meaning there exists integer c such that $b = ac$.

Consider n +1 distinct positive integers, each $\leq 2n$. Show that one of them must divide one of the others.

For example, if $n = 4$, consider the following sets:

{1, 2, 3, 7, 8} {2, 3, 4, 7, 8} {2, 3, 5, 7, 8}

Hint: Any integer can be written as $q*2^k$ where k is a non-negative integer and q is odd. E.g., $129 = 2^0 * 129$; $60 = 2^2 * 15$.

The Pigeonhole Principle (Full Glory)

Let X and Y be finite sets with $|X| = n$, $|Y| = m$, and $k = \lceil n/m \rceil$.

If $f : X \rightarrow Y$, then there exist k values $x_1, x_2, \ldots, x_k$ in X such that $f(x_1) = f(x_2) = \ldots = f(x_k)$.

Informally: If n pigeons fly into m holes, at least 1 hole contains at least $k = \lceil n/m \rceil$ pigeons.

Proof: Assume there's no such hole. Then, there are at most $(\lceil n/m \rceil - 1)*m$ pigeons in all the holes, which is fewer than $(n/m + 1 - 1)*m = n/m*m = n$, but that is a contradiction. QED

Birthday Paradox

- Mathematically, the problem of collisions is more related to the “Birthday Paradox” rather than the Pigeonhole Principle
- What’s the probability that in a room of 23 people, at least 2 people have the same birthday?

Birthday Paradox

- Mathematically, the problem of collisions is more related to the “Birthday Paradox” rather than the Pigeonhole Principle
- What’s the probability that in a room of 23 people, at least 2 people have the same birthday?

About 50%!

So “hashing” 23 people into 365 slots has a 50% of having at least one collision…

Birthday Paradox Explained

- What's the probability that n people all have different birthdays?

$$\frac{365}{365}\times\frac{364}{365}\times\frac{363}{365}\times\frac{362}{365}\times\cdots\times\frac{365-n+1}{365}$$

Those fractions quickly drop the probability toward 0.

Birthday Paradox Approximate Rule of Thumb

- Probability of at least one collision given n keys hashed to size slots is approximately:

$$P(collision) \approx \frac{n^2}{2 \cdot size}$$

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Collision Resolution

- *Pigeonhole principle* says we can't avoid all collisions
  - try to hash without collision m keys into n slots with $m > n$
  - try to put 6 pigeons into 5 holes
- What do we do when two keys hash to the same entry?
  - chaining: put little dictionaries in each entry

    shove extra pigeons in one hole!
  - open addressing: pick a next entry to try

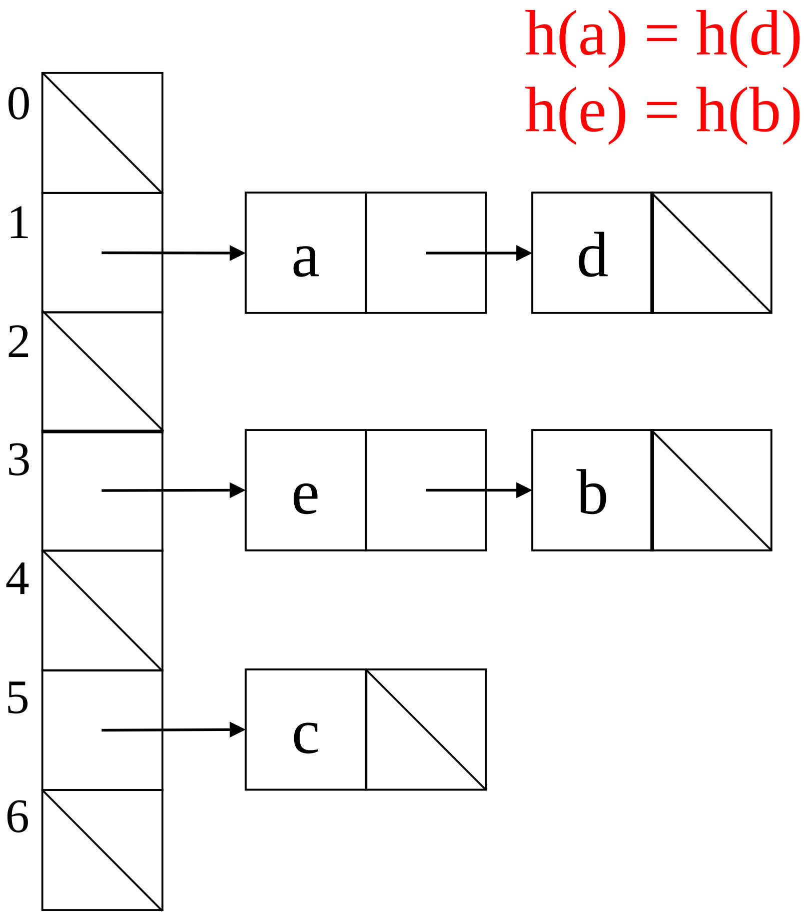
(Alan Aside) Collision Resolution

- *Pigeonhole principle* says we can't avoid all collisions
  - try to hash without collision m keys into n slots with $m > n$
  - try to put 6 pigeons into 5 holes
- What do we do when two keys hash to the same entry?
  - chaining (AKA open hashing or closed addressing): put little dictionaries in each entry

    shove extra pigeons in one hole!
  - open addressing (AKA closed hashing): pick a next entry to try

Hashing with Chaining

- Put a little dictionary at each entry
  - choose type as appropriate
  - common case is unordered linked list (chain)
- Properties
  - λ can be greater than 1
  - performance degrades with length of chains



Chaining Code

```
Dictionary & findBucket(const Key & k) {
  return table[hash(k)%table.size];
}

void insert(const Key & k,
            const Value & v)
{
  findBucket(k).insert(k,v);
}

void delete(const Key & k)
{
  findBucket(k).delete(k);
}

Value & find(const Key & k)
{
  return findBucket(k).find(k);
}
```

Load Factor in Chaining

- Search cost
  - unsuccessful search:
  - successful search:
- Desired load factor:

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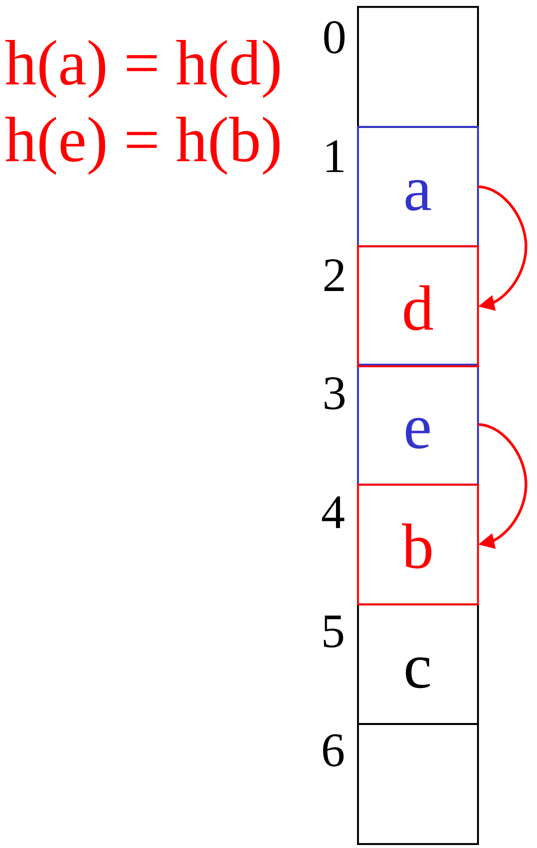
Open Addressing / Closed Hashing

What if we only allow one key at each entry?

  – two objects that hash to the same spot can't both go there
  – first one there gets the spot
  – next one must *go in another spot*

- Properties
  – $\lambda \leq 1$
  – performance degrades with difficulty of finding right spot



Probing

- Probing how to:
  - First probe - given a key k, hash to h(k)
  - Second probe - if h(k) is occupied, try h(k) + f(1)
  - Third probe - if h(k) + f(1) is occupied, try h(k) + f(2)
  - And so forth
- Probing properties
  - the i^{th} probe is to (h(k) + f(i)) mod size   where f(0) = 0
  - if i reaches size, the insert has failed
  - depending on f(), the insert may fail sooner
  - long sequences of probes are costly!

Linear Probing

$f(i) = i$

- Probe sequence is
  - h(k) mod size
  - h(k) + 1 mod size
  - h(k) + 2 mod size
  - …
- findEntry using linear probing:

```
bool findEntry(const Key & k, Entry *& entry) {
  int probePoint = hash1(k);
  int i=0;
  do {
    entry = &table[(probePoint+(i++)) % size];
  } while (!entry->isEmpty() && entry->key != k);
  return !entry->isEmpty();
}
```

Linear Probing (More Efficient Code)

$f(i) = i$

- Probe sequence is
  - h(k) mod size
  - h(k) + 1 mod size
  - h(k) + 2 mod size
  - …
- findEntry using linear probing:

```
bool findEntry(const Key & k, Entry *& entry) {
  int probePoint = hash1(k);
  do {
    entry = &table[probePoint];
    probePoint = (probePoint + 1) % size;
  } while (!entry->isEmpty() && entry->key != k);
  return !entry->isEmpty();
}
```

Linear Probing Example

	insert(76)	insert(93)	insert(40)	insert(47)	insert(10)	insert(55)
	76%7 = 6	93%7 = 2	40%7 = 5	47%7 = 5	10%7 = 3	55%7 = 6
0				47	47	47
1						55
2		93	93	93	93	93
3					10	10
4						
5			40	40	40	40
6	76	76	76	76	76	76
probes:	1	1	1	3	1	3

Load Factor in Linear Probing

- For *any* $\lambda < 1$, linear probing will find an empty slot
- Search cost (for large table sizes)
  - successful search: $$\frac{1}{2}\left(1+\frac{1}{(1-\lambda)}\right)$$
  - unsuccessful search: $$\frac{1}{2}\left(1+\frac{1}{(1-\lambda)^2}\right)$$

Values hashed close to each other probe the same slots.

- Linear probing suffers from *primary clustering*
- Performance quickly degrades for $\lambda > 1/2$

Quadratic Probing

$f(i) = i^2$

- Probe sequence is
  - h(k) mod size
  - (h(k) + 1) mod size
  - (h(k) + 4) mod size
  - (h(k) + 9) mod size
  - …
- findEntry using quadratic probing:

```
bool findEntry(const Key & k, Entry *& entry) {
  int probePoint = hash1(k), i = 0;
  do {
    entry = &table[(probePoint + i*i) % size];
    i++;
  } while (!entry->isEmpty() && entry->key != key);
  return !entry->isEmpty();
}
```

Quadratic Probing (more efficient code)

$f(i) = i^2$

- Probe sequence is
  - h(k) mod size
  - (h(k) + 1) mod size
  - (h(k) + 4) mod size
  - (h(k) + 9) mod size
  - …
- findEntry using quadratic probing:

```
bool findEntry(const Key & k, Entry *& entry) {
  int probePoint = hash1(k), i = 0;
  do {
    entry = &table[probePoint];
    i++;
    probePoint = (probePoint + 2*i - 1) % size;
  } while (!entry->isEmpty() && entry->key != key);
  return !entry->isEmpty();
}
```

Quadratic Probing Example ☺

	insert(76)	insert(40)	insert(48)	insert(5)	insert(55)
	76%7 = 6	40%7 = 5	48%7 = 6	5%7 = 5	55%7 = 6
0			48	48	48
1					
2				5	5
3					55
4					
5		40	40	40	40
6	76	76	76	76	76
probes:	1	1	2	3	3

Quadratic Probing Example ☹

	insert(76)	insert(93)	insert(40)	insert(35)	insert(47)
	76%7 = 6	93%7 = 2	40%7 = 5	35%7 = 0	47%7 = 5
0				35	35
1					
2		93	93	93	93
3					
4					
5			40	40	40
6	76	76	76	76	76
probes:	1	1	1	1	∞

Quadratic Probing Succeeds (for $\lambda \leq ½$)

- If size is prime and $\lambda \leq ½$, then quadratic probing will find an empty slot in size/2 probes or fewer.
  - show for all **0 ≤ i, j ≤ size/2** and **i ≠ j**

    $$(h(x) + i^2) \bmod size \neq (h(x) + j^2) \bmod size$$
  - this means that the size/2 probes must all land in different places, so at least one must succeed if $\lambda \leq ½$

Quadratic Probing Succeeds (for $\lambda \leq \frac{1}{2}$)

- If size is prime and $\lambda \leq \frac{1}{2}$, then quadratic probing will find an empty slot in size/2 probes or fewer.
  - show for all $\mathtt{0 \leq i, j \leq size/2}$ and $\mathtt{i \neq j}$

    $\mathtt{(h(x) + i^2) \bmod size \neq (h(x) + j^2) \bmod size}$
  - by contradiction: suppose that for some i, j:

    $\mathtt{(h(x) + i^2) \bmod size = (h(x) + j^2) \bmod size}$

    $\mathtt{i^2 \bmod size = j^2 \bmod size}$

    $\mathtt{(i^2 - j^2) \bmod size = 0}$

    $\mathtt{[(i + j)(i - j)] \bmod size = 0}$
  - but how can $\mathtt{i + j = 0}$ or $\mathtt{i + j = size}$ when

    $\mathtt{i \neq j}$ and $\mathtt{i,j \leq size/2}$?
  - same for $\mathtt{i - j \bmod size = 0}$

Quadratic Probing May Fail (for λ > ½)

- For any i larger than size/2, there is some j smaller than i that adds with i to equal size (or a multiple of size). D'oh!

Let $i = size - j$

$$i^2 = (size - j)^2 = size^2 - 2size \cdot j + j^2 \equiv j^2 (\mathrm{mod}\, size)$$

Load Factor in Quadratic Probing

- For *any* $\lambda \leq \frac{1}{2}$, quadratic probing will find an empty slot; for greater λ, quadratic probing *may* find a slot
- Quadratic probing does not suffer from primary clustering
- Quadratic probing *does* suffer from *secondary* clustering
  - How could we possibly solve this?

Values hashed to the SAME index probe the same slots.

Double Hashing

$f(i) = i \cdot hash_2(k)$

- Probe sequence is
  - $h_1(k) \bmod size$
  - $(h_1(k) + 1 \cdot h_2(k)) \bmod size$
  - $(h_1(k) + 2 \cdot h_2(k)) \bmod size$
  - …
- Code for finding the next linear probe:

```
bool findEntry(const Key & k, Entry *& entry) {
  int probePoint = hash1(k), hashIncr = hash2(k);
  do {
    entry = &table[probePoint];
    probePoint = (probePoint + hashIncr) % size;
  } while (!entry->isEmpty() && entry->key != k);
  return !entry->isEmpty();
}
```

A Good Double Hash Function…

…is quick to evaluate.

…differs from the original hash function.

…never evaluates to 0 (mod size).

One good choice is to choose a prime $R <$ size and:

$$hash_2(x) = R - (x \bmod R)$$

Double Hashing Example

	insert(76)	insert(93)	insert(40)	insert(47)	insert(10)	insert(55)
	76%7 = 6	93%7 = 2	40%7 = 5	47%7 = 5 5 - (47%5) = 3	10%7 = 3	55%7 = 6 5 - (55%5) = 5
0						
1				47	47	47
2		93	93	93	93	93
3					10	10
4						55
5			40	40	40	40
6	76	76	76	76	76	76
probes:	1	1	1	2	1	2

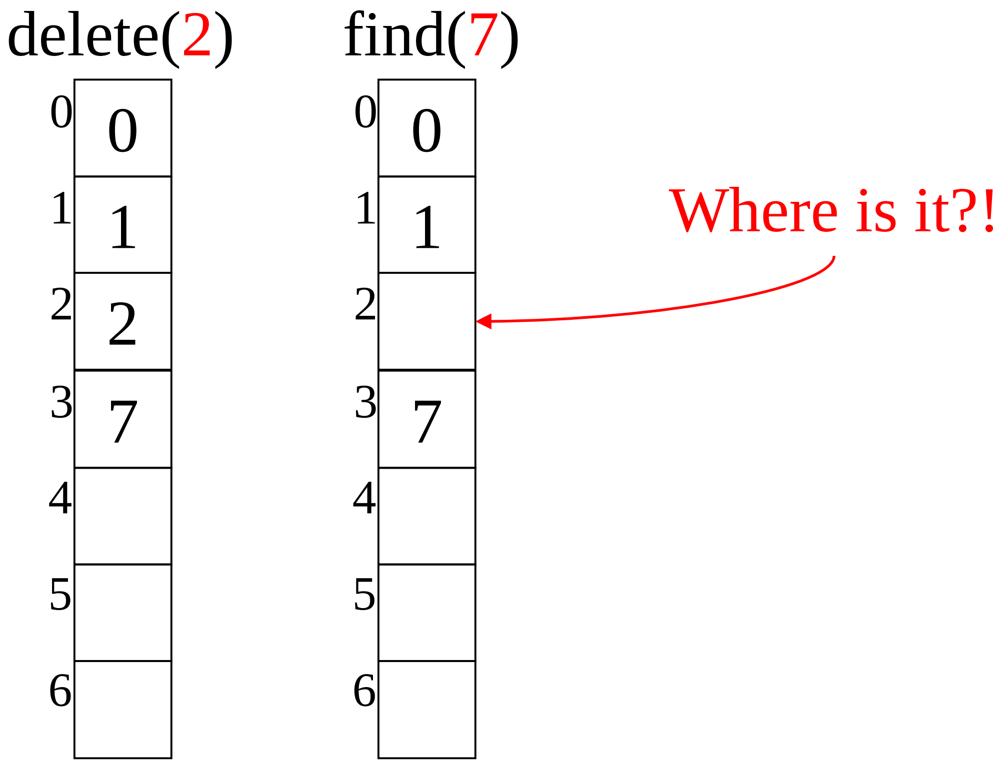
Load Factor in Double Hashing

- For *any* $\lambda < 1$, double hashing will find an empty slot (given appropriate table size and $hash_2$)
- Search cost appears to approach optimal (random hash):
  - successful search: $$\frac{1}{\lambda}\ln\frac{1}{1-\lambda}$$
  - unsuccessful search: $$\frac{1}{1-\lambda}$$
- No primary clustering and no secondary clustering
- One extra hash calculation

Outline

- Constant-Time Dictionaries?
- Hash Table Overview
- Hash Functions
- Collisions and the Pigeonhole Principle
- Collision Resolution:
  - Chaining
  - Open-Addressing
- Deletion and Rehashing

Deletion in Open Addressing



- Must use lazy deletion!
- On insertion, treat a deleted item as an empty slot

The “Squished Pigeon Principle”

- An insert using open addressing *cannot* work with a load factor of 1 or more.
- An insert using open addressing with quadratic probing may not work with a load factor of ½ or more.
- Whether you use chaining or open addressing, large load factors lead to poor performance!
- How can we relieve the pressure on the pigeons?

Hint: think resizable arrays!

Rehashing

- When the load factor gets “too large” (over a constant threshold on λ), rehash all the elements into a new, larger table:
  - takes O(n), but amortized O(1) as long as we (just about) double table size on the resize
  - spreads keys back out, may drastically improve performance
  - gives us a chance to retune parameterized hash functions
  - avoids failure for open addressing techniques
  - allows arbitrarily large tables starting from a small table
  - clears out lazily deleted items