



CPSC 490 – Problem Solving in Computer Science

Lecture 19: Line Sweep II

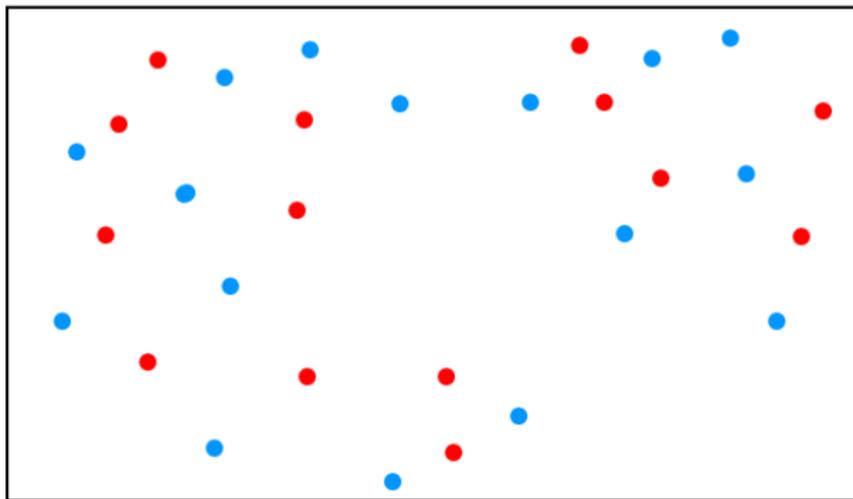
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2017/03/08

University of British Columbia

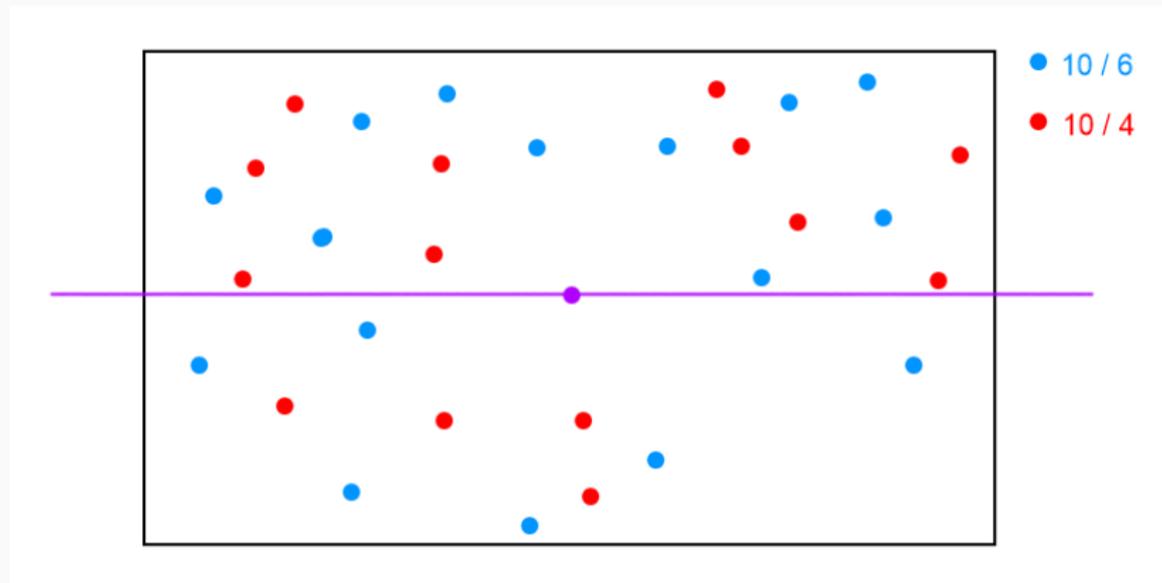
Problem 1: Angular Sweep

Divide the rectangle into two halves with equal area such that the two sides have an equal number of blue dots and red dots.



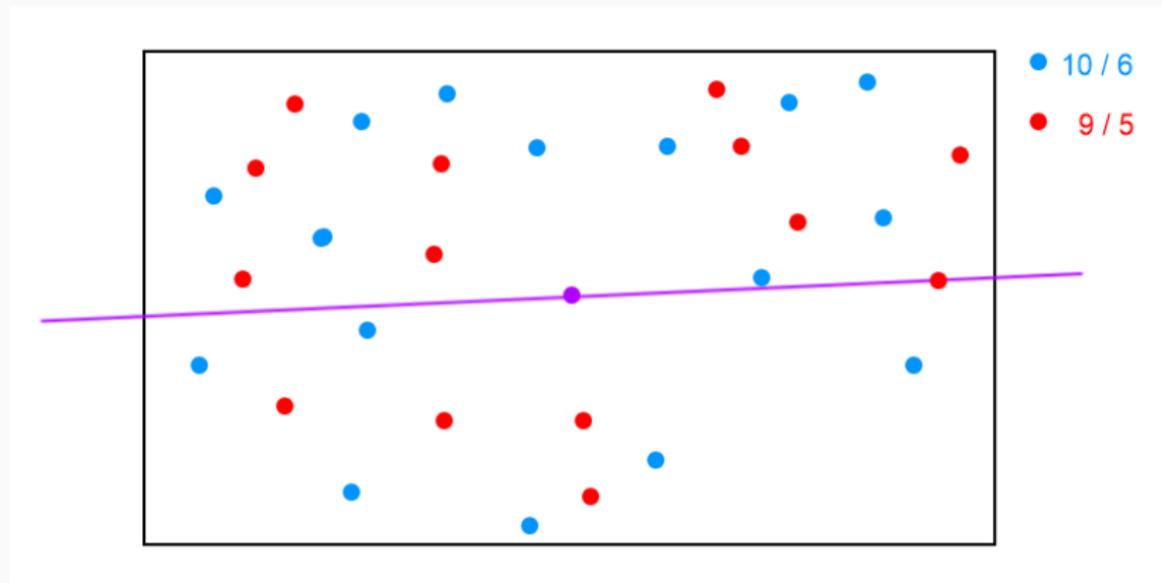
Problem 1: Solution

Observation: line must go through center \Rightarrow do an “angular sweep” by rotating a line around the center of rectangle!



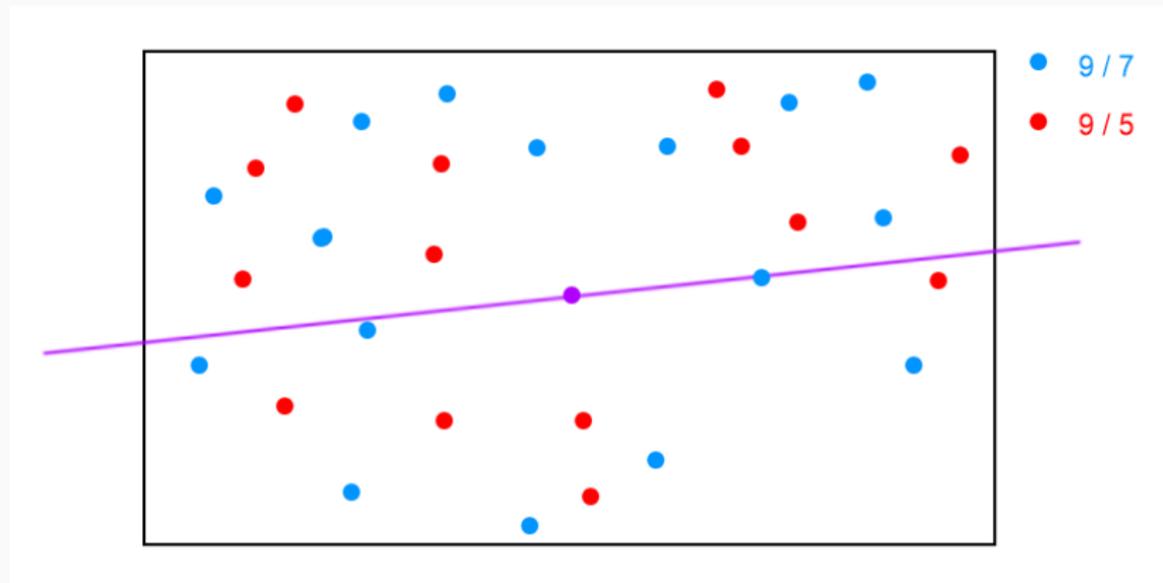
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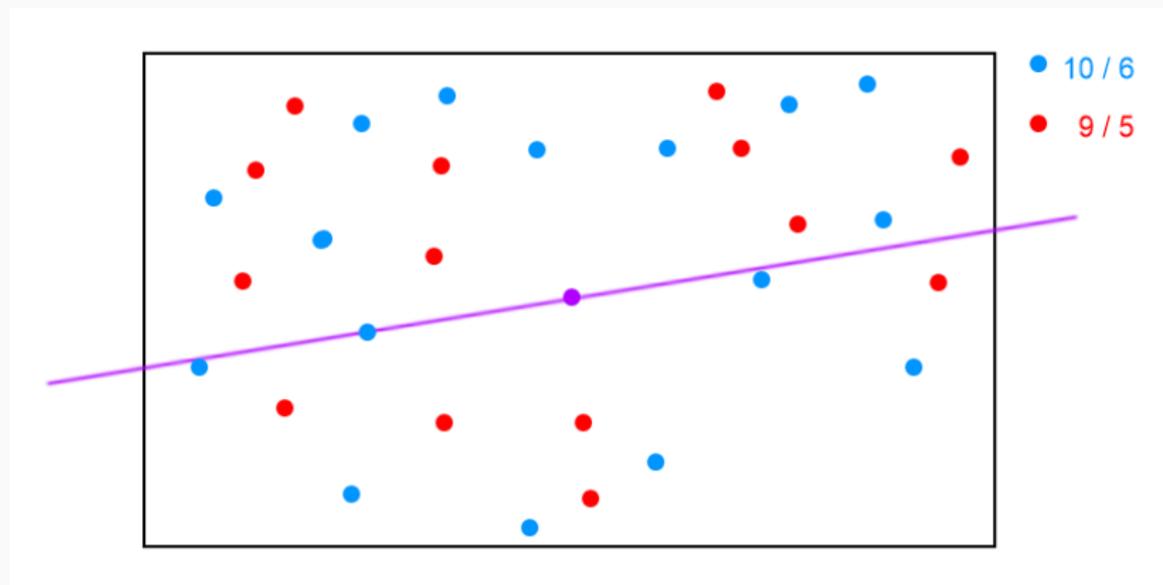
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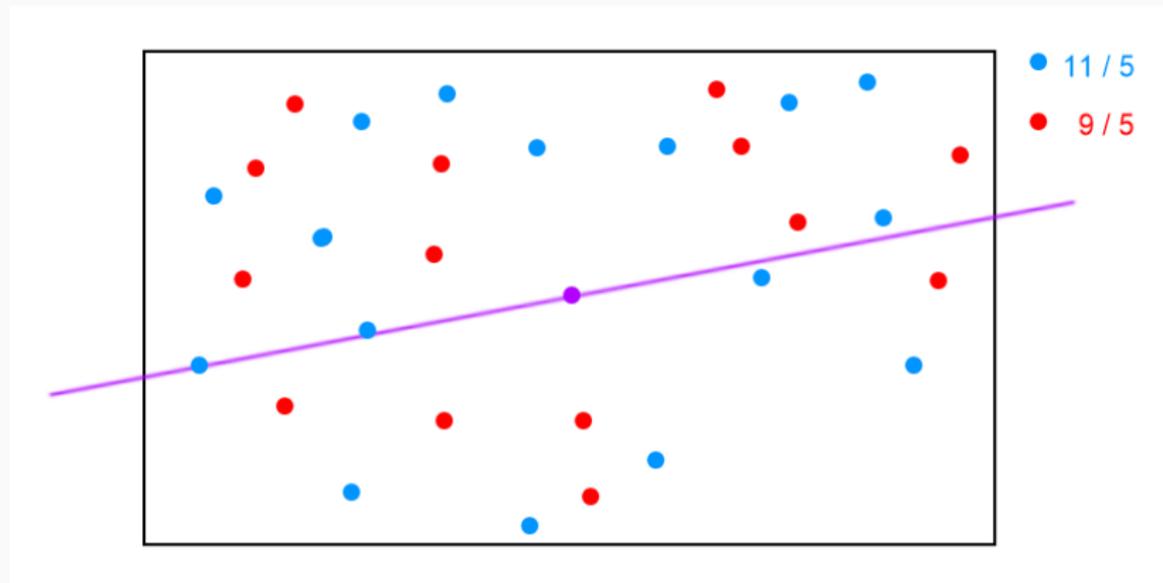
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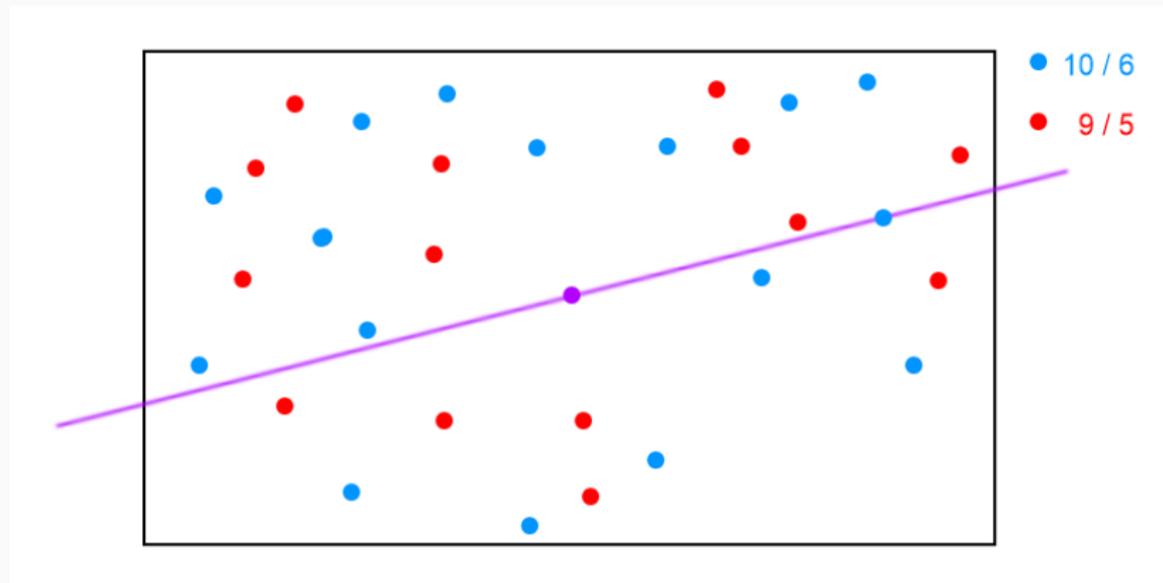
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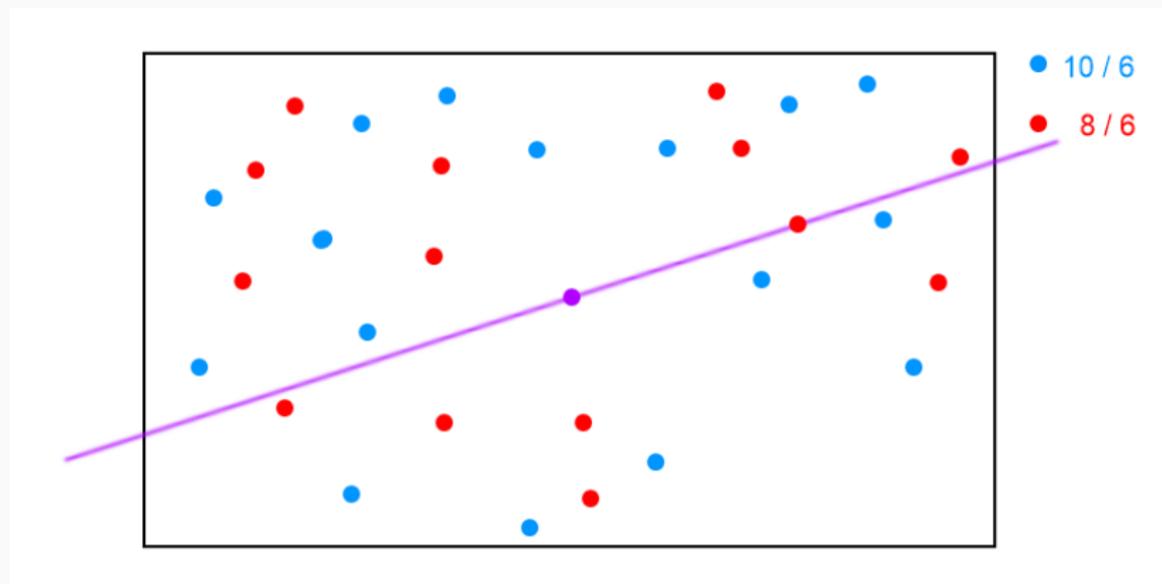
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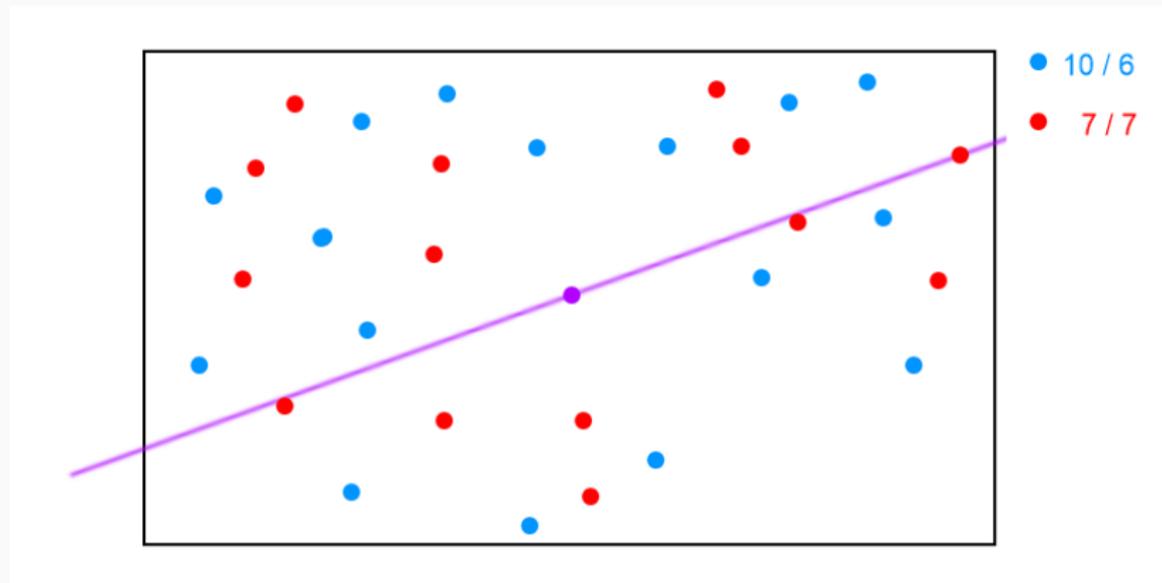
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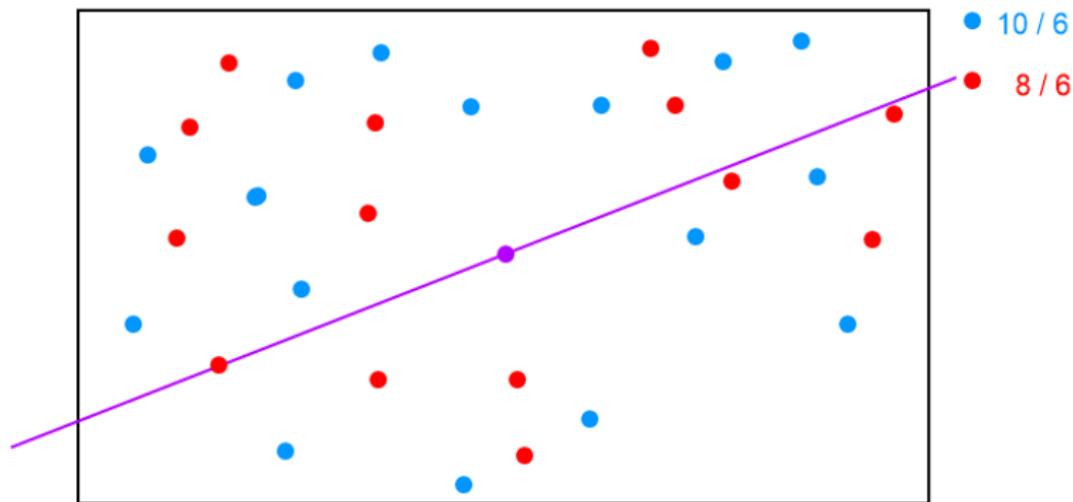
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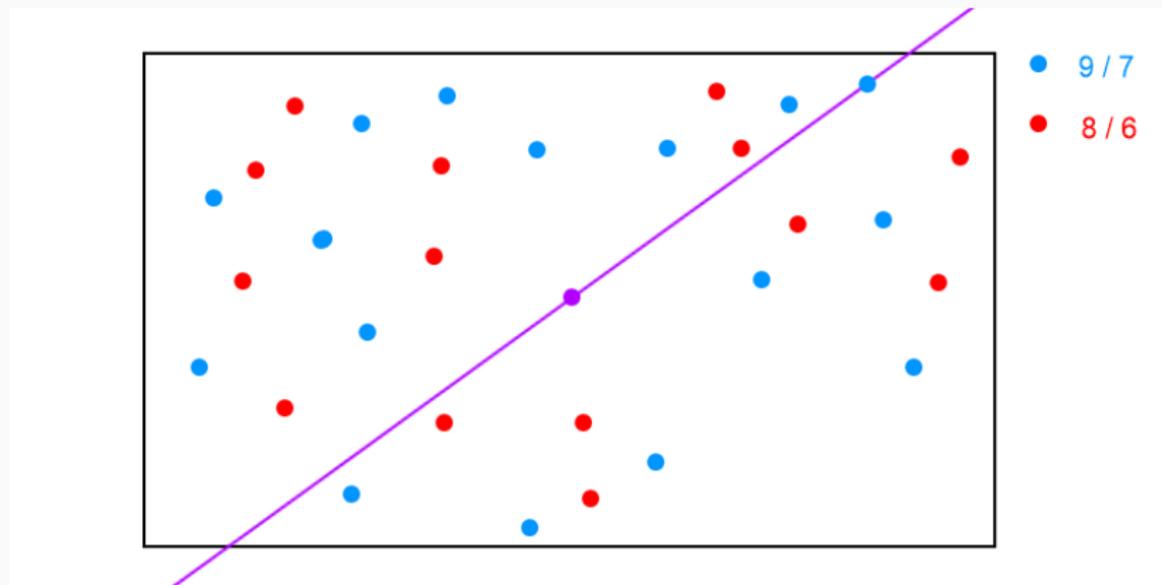
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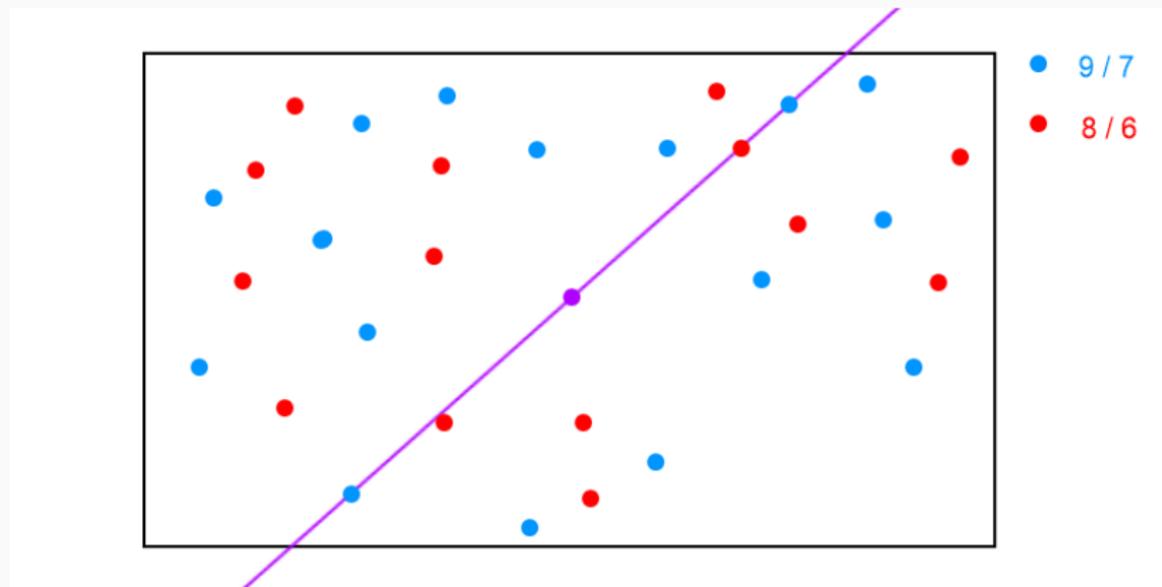
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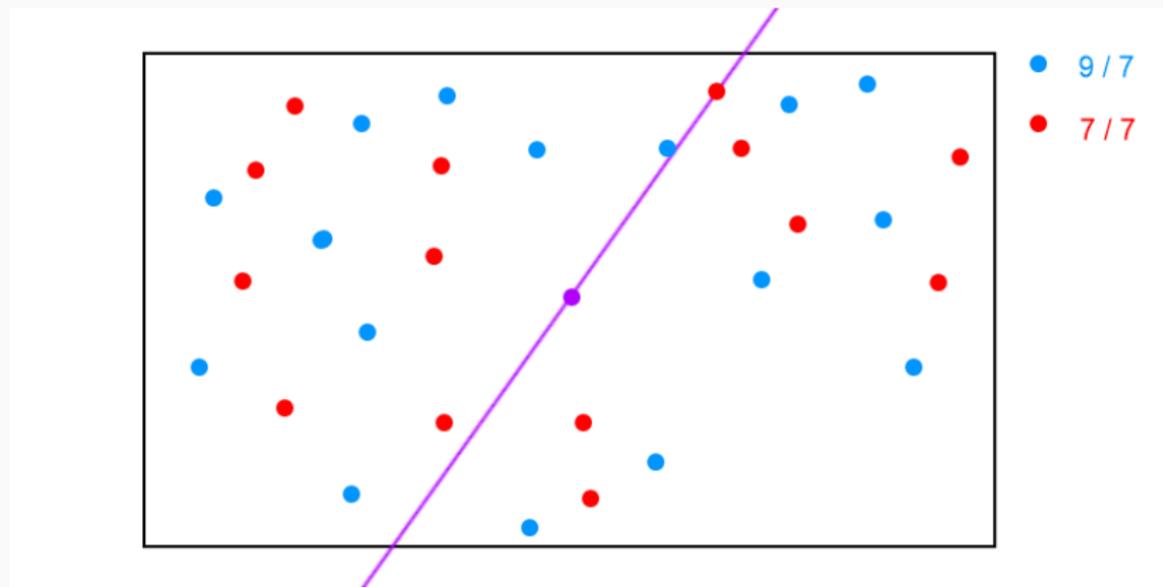
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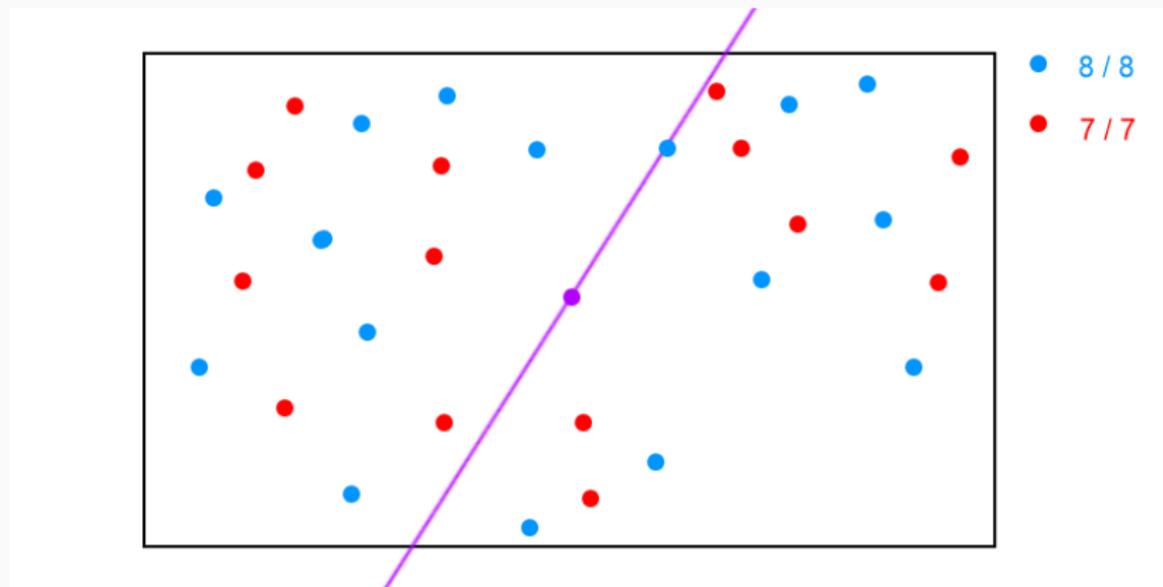
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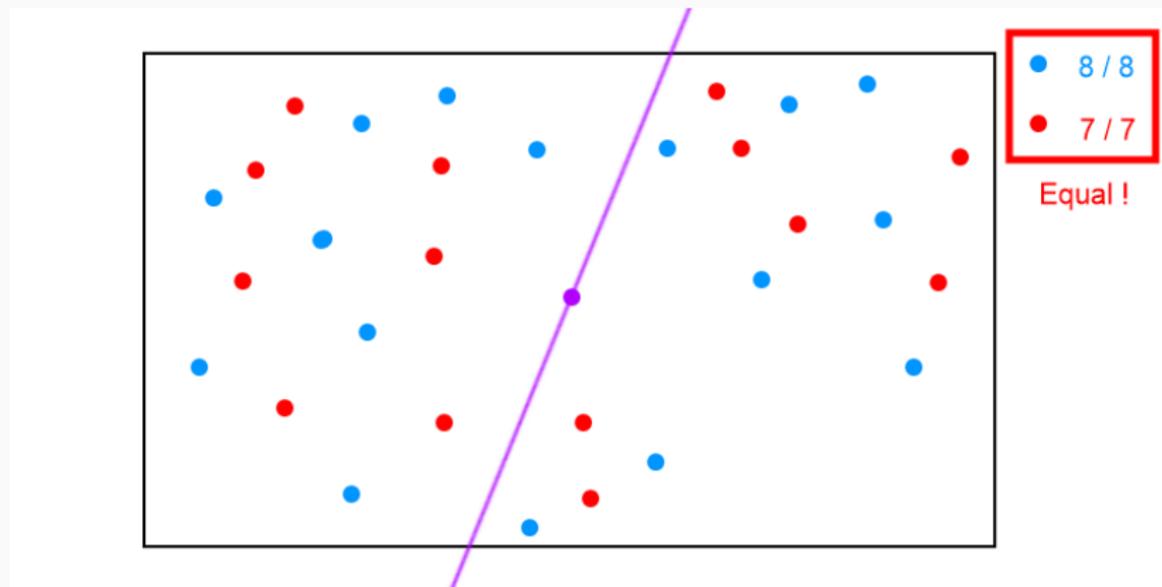
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Problem 1: Angular Sweep

State: how many of each kind of dot is on each side.

Event: line hits a dot on either side, sorted by angle around center.
(i.e. must add a “phantom event” on opposite side for dots $\geq 180^\circ$)

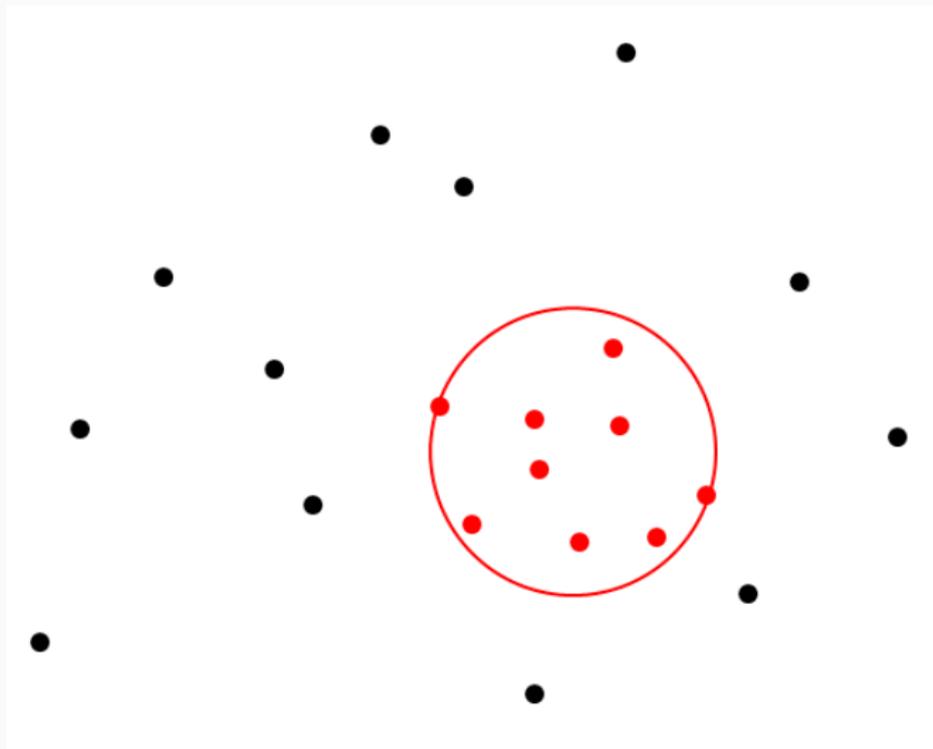
Action: when line hits a dot, +1 or -1 to the relevant state.

Time complexity: $O(n \log n)$

Problem 2: Angular Sweep with Circle

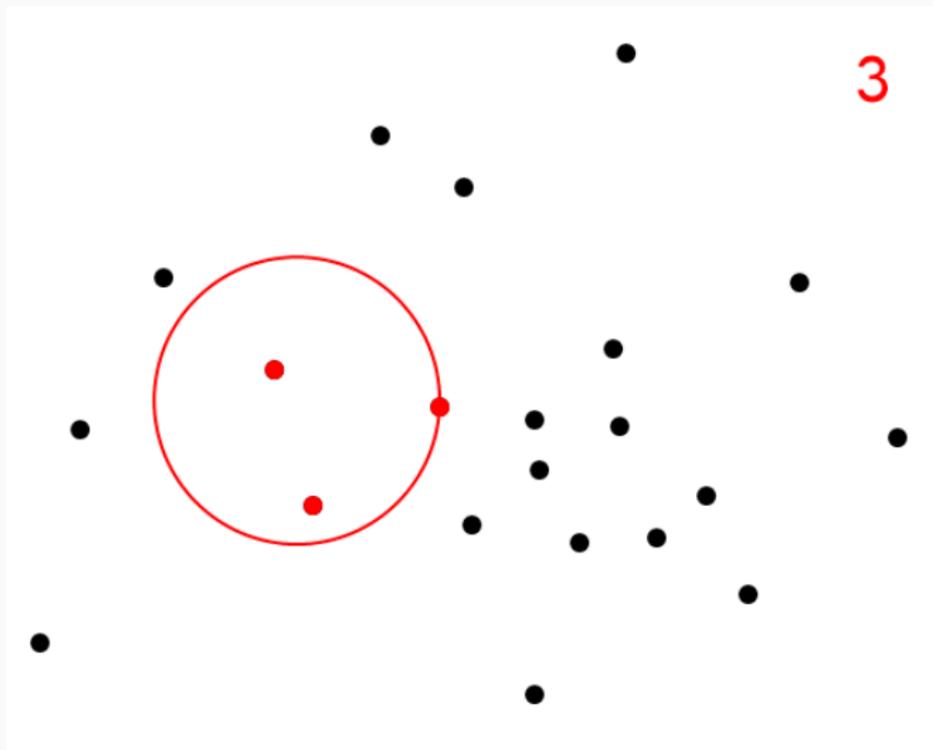
Find circle with radius R that covers the most points

Hint: rotate *what* around *where*? $O(n^2 \log n)$



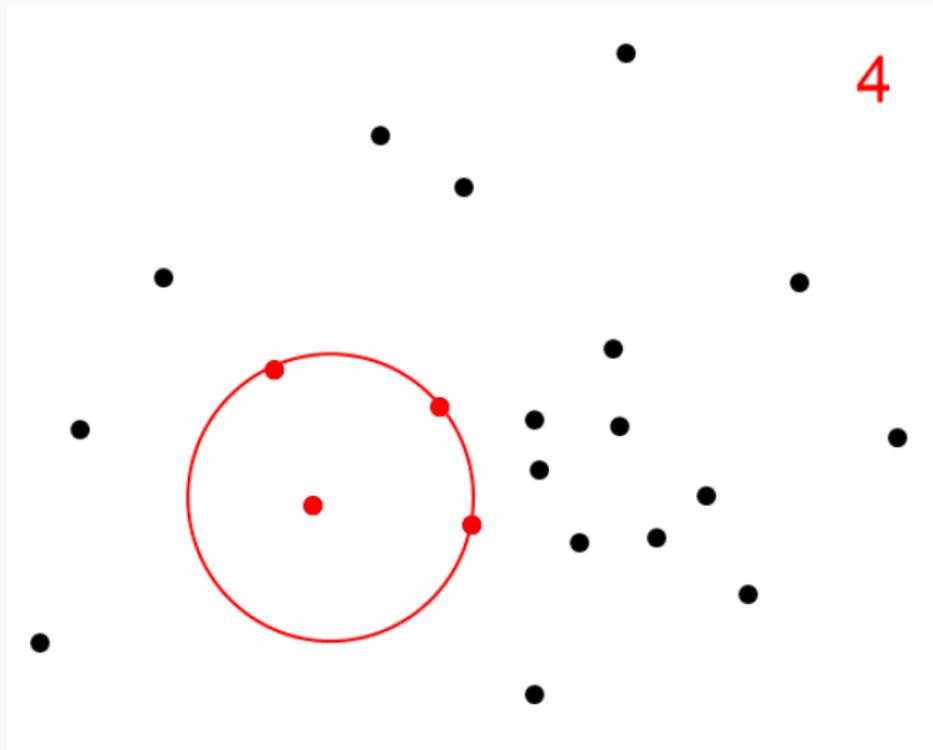
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Observation: can wiggle circle so that it tangents one point \Rightarrow try every single point, rotate a circle around that point.



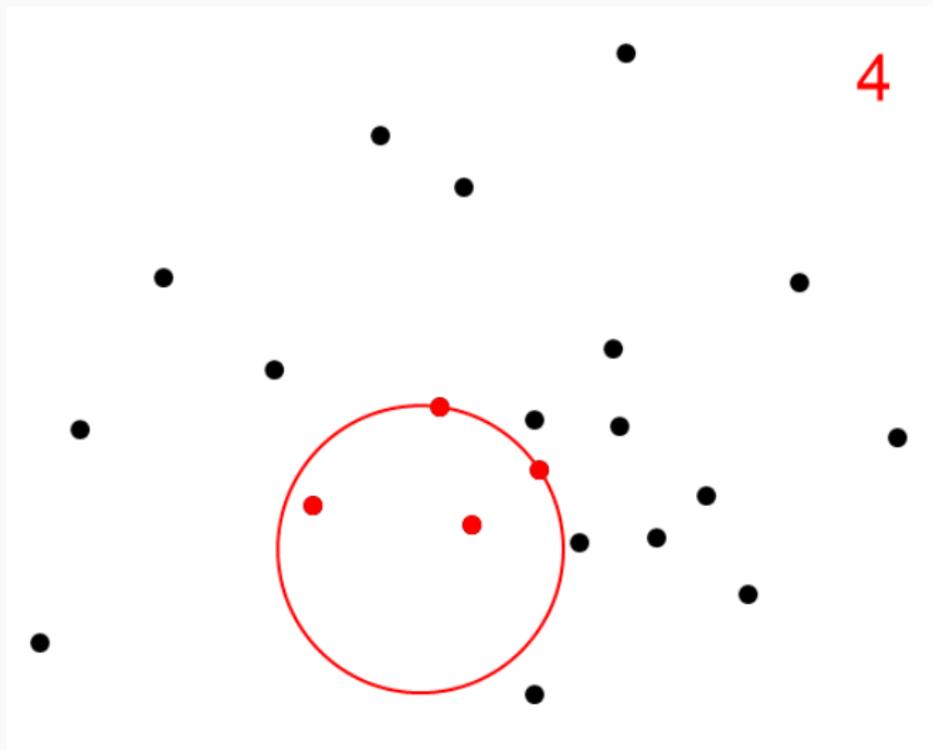
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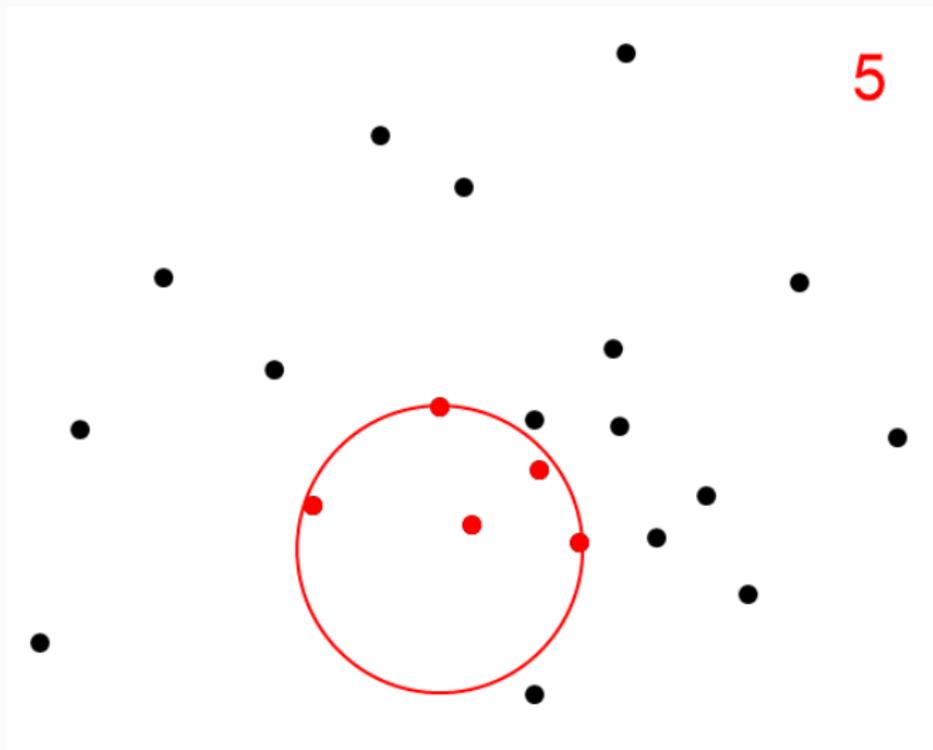
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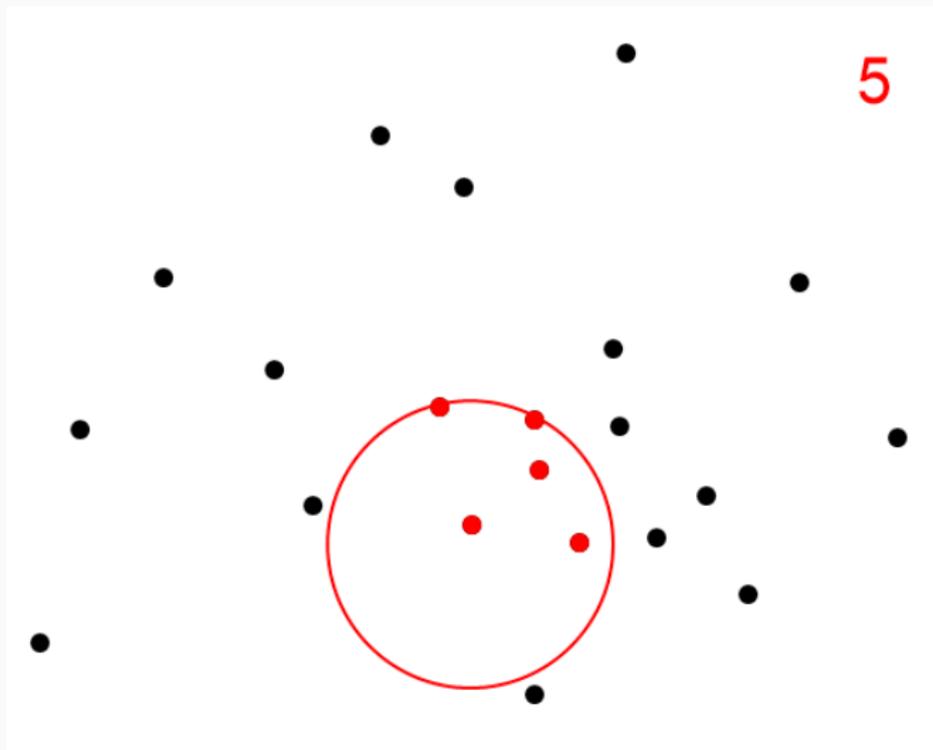
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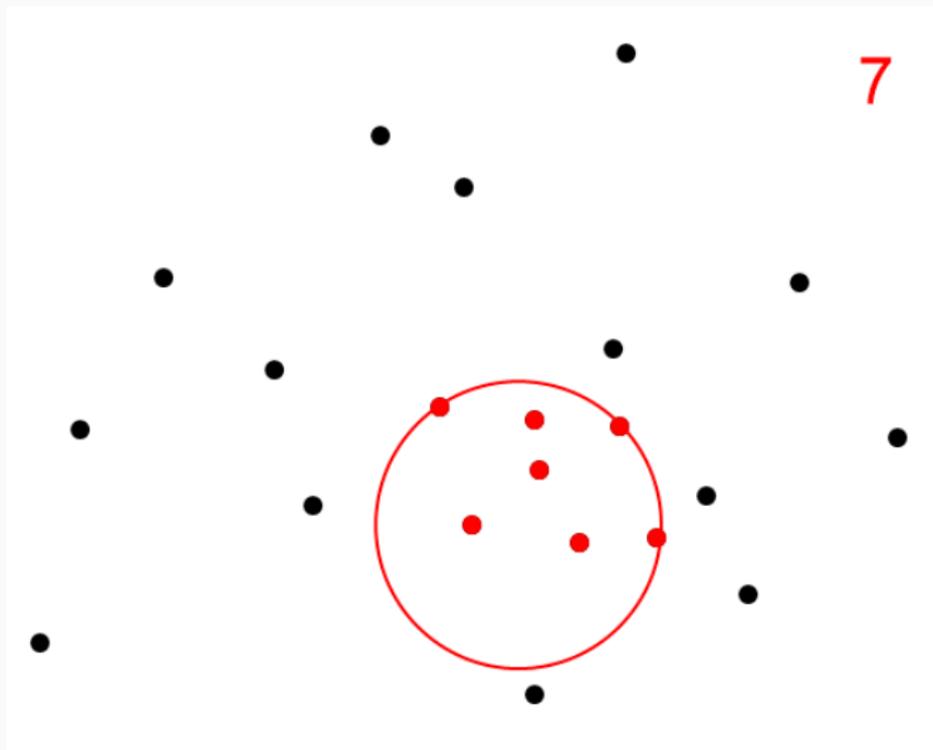
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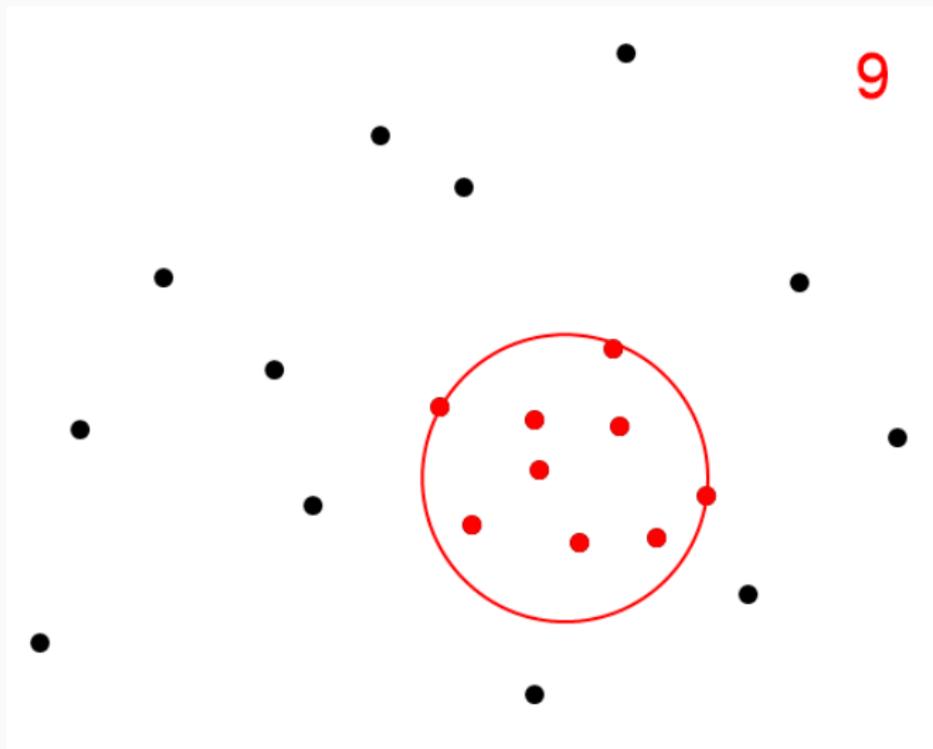
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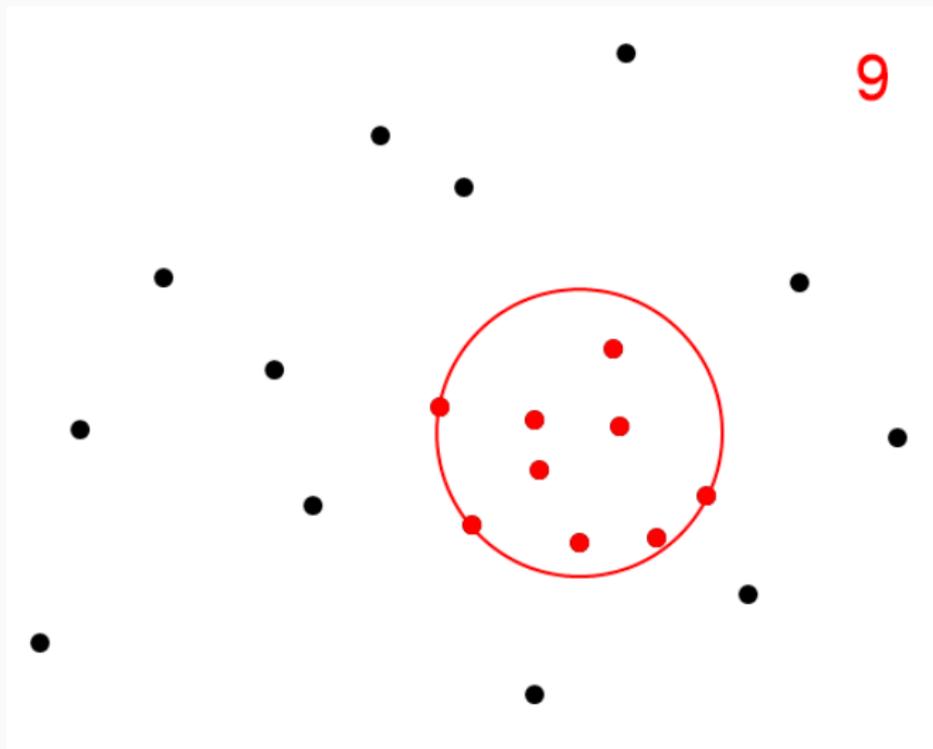
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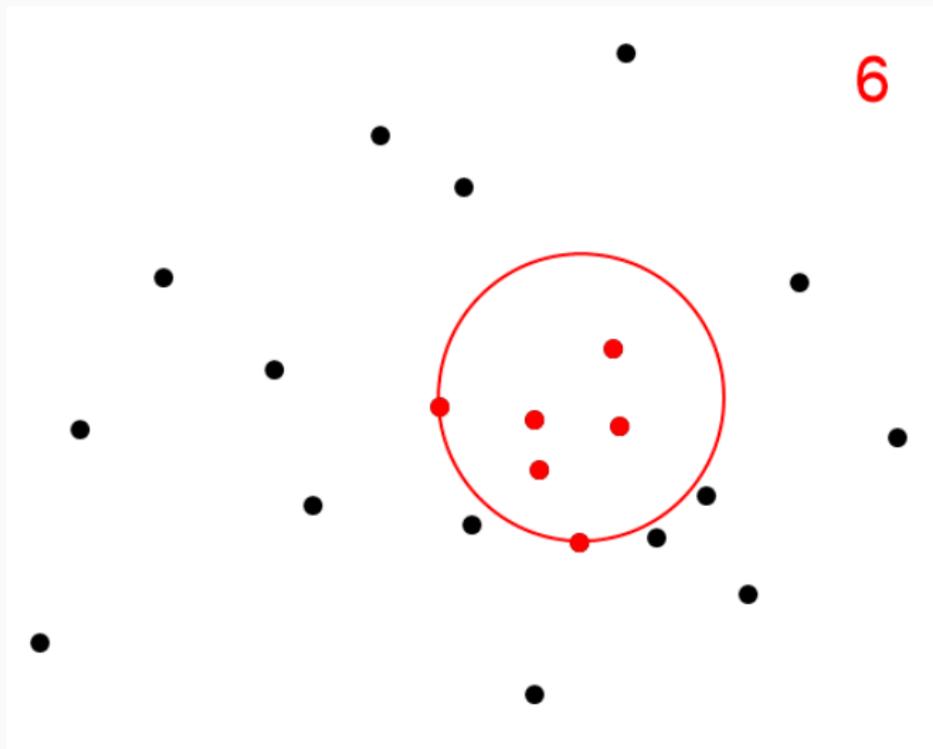
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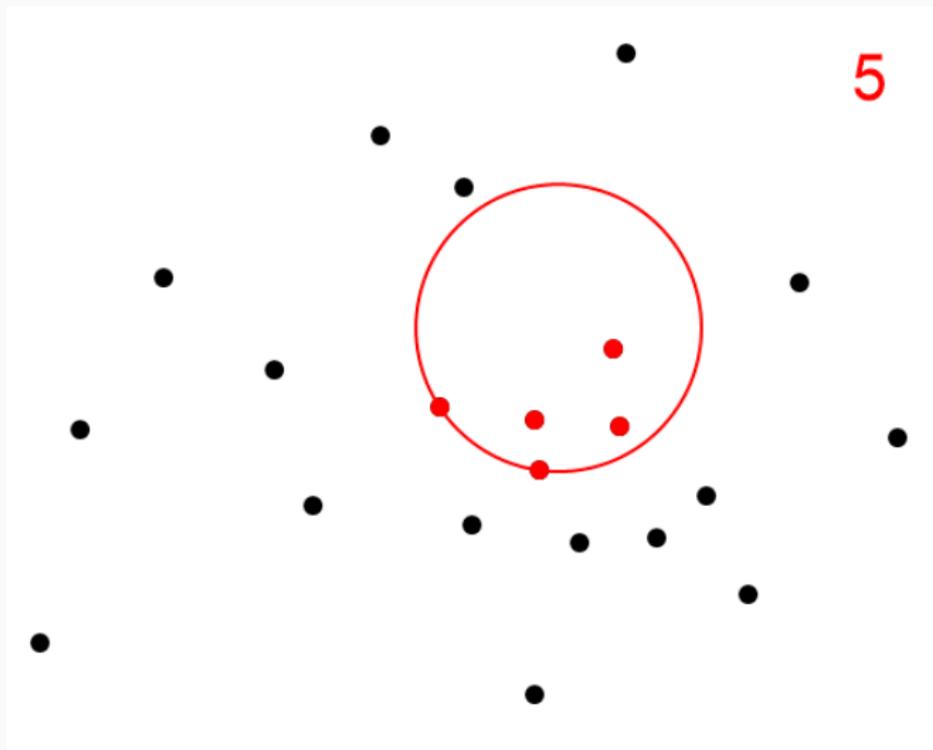
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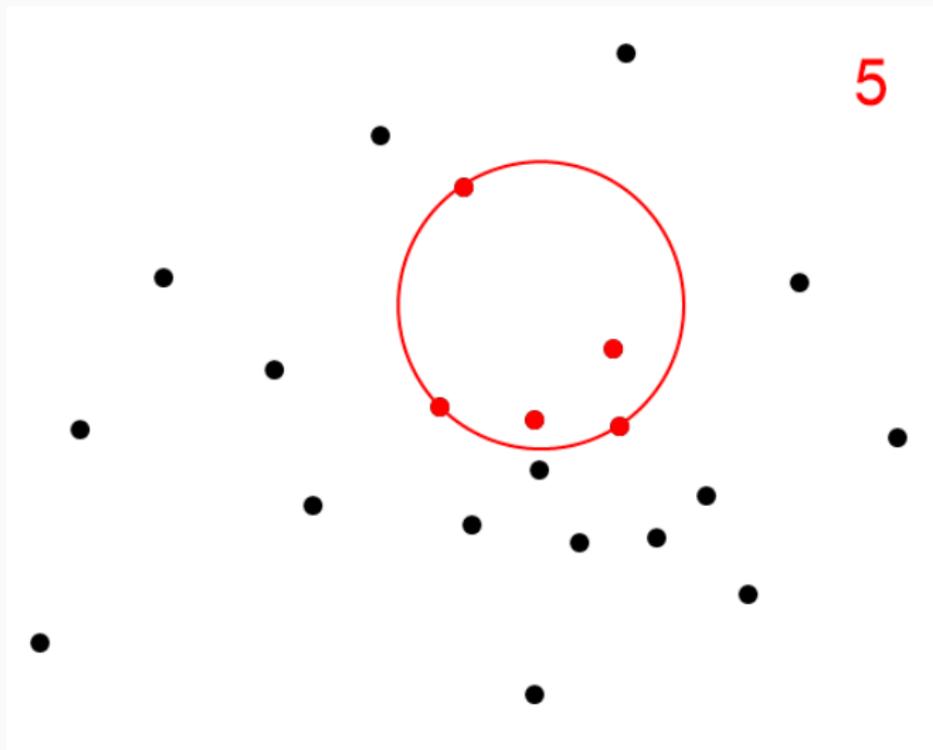
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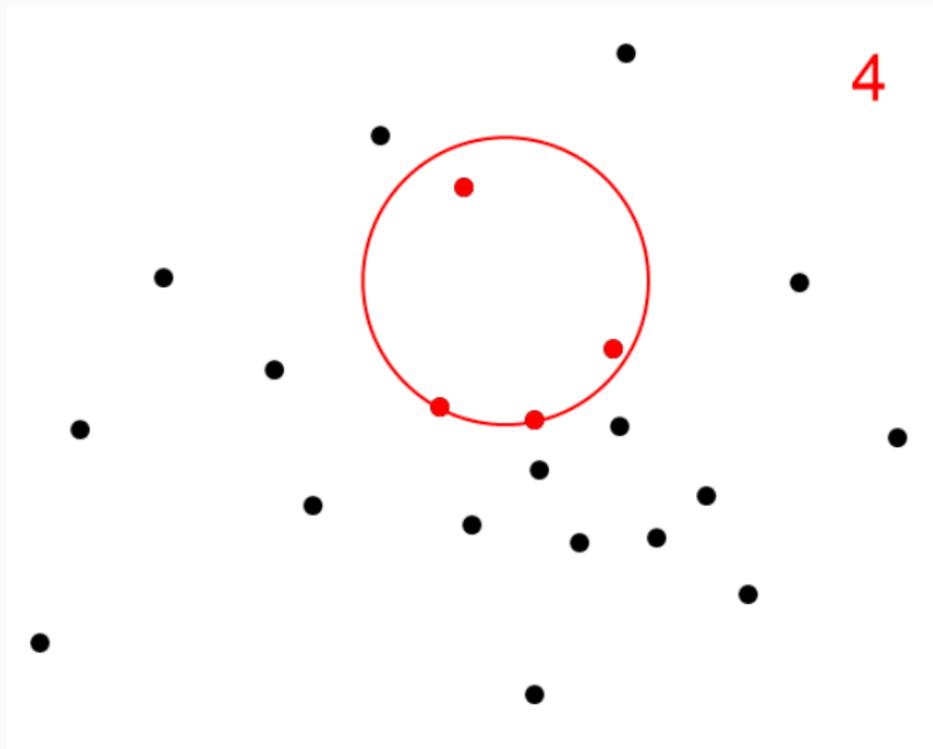
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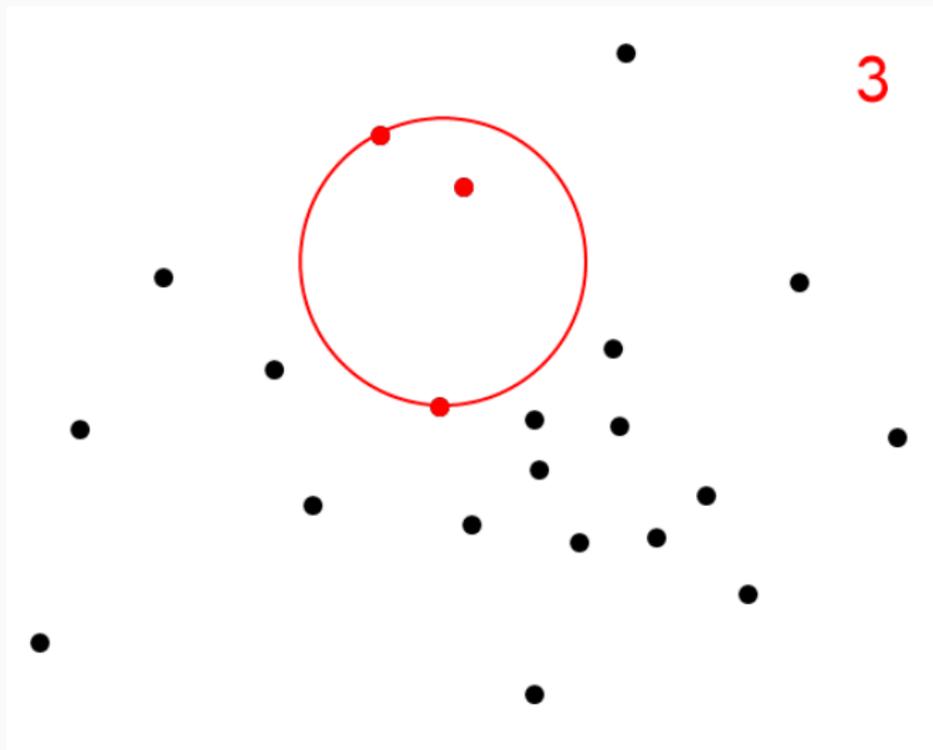
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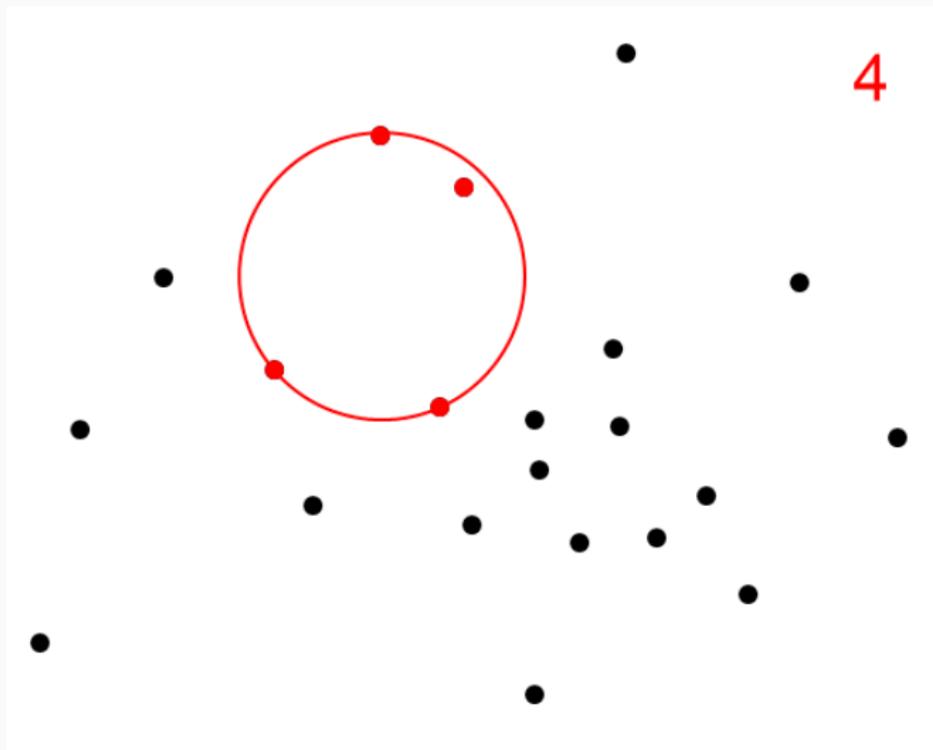
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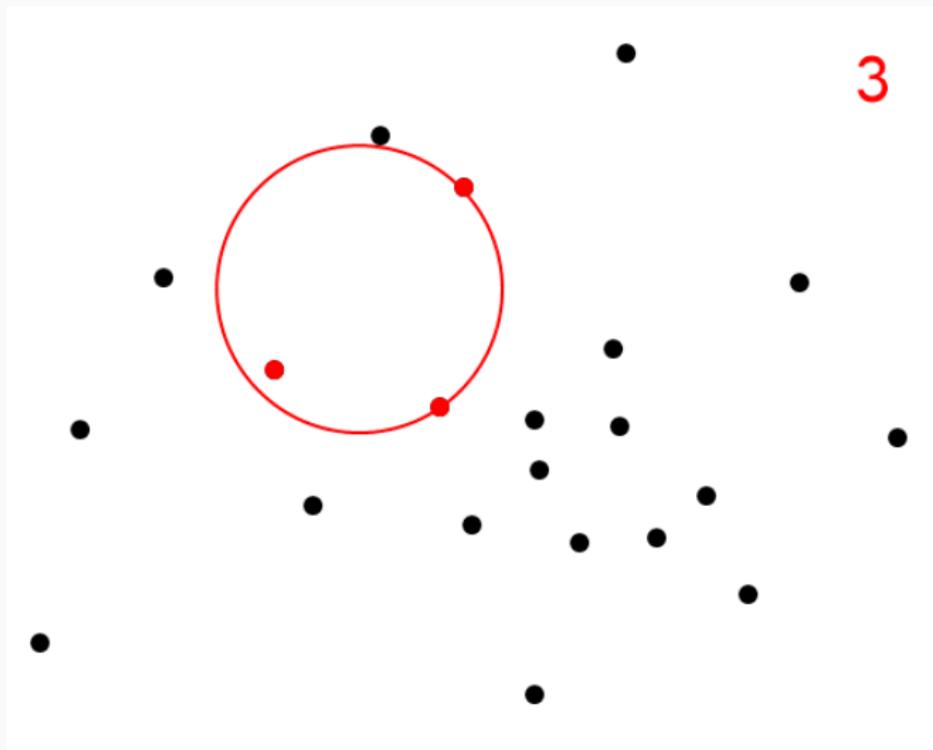
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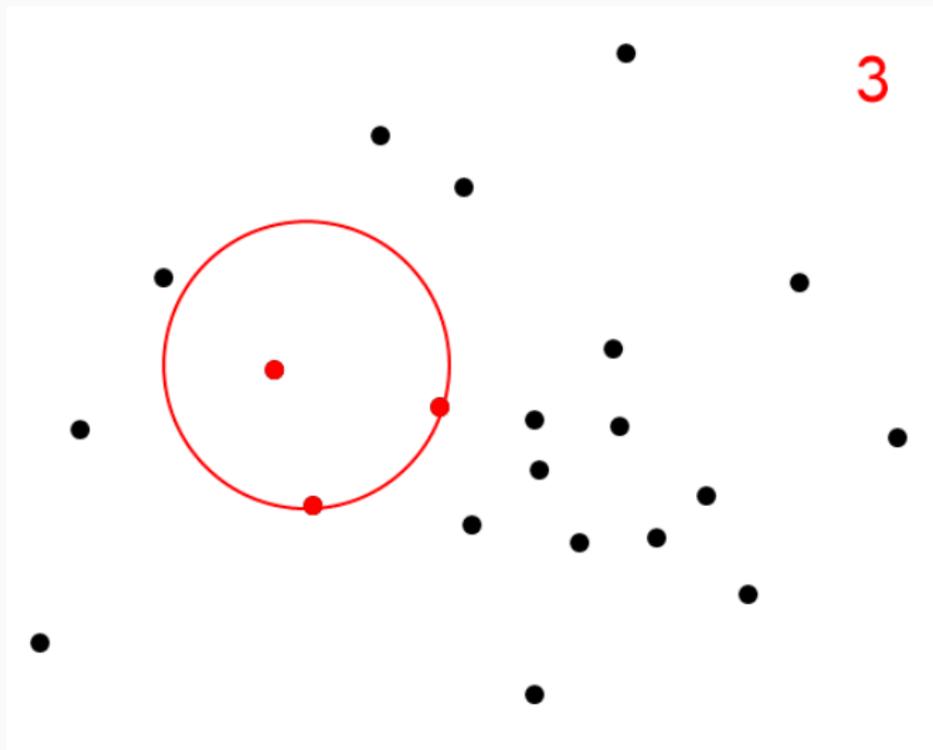
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Problem 2: Angular Sweep with Circle

Idea: for every point, rotate a circle 360° around it

State: how many points are inside the rotating circle

Event: when point enters/exits circle, sorted by rotation angle of circle, be careful about “wrapping across” 360° and initial condition

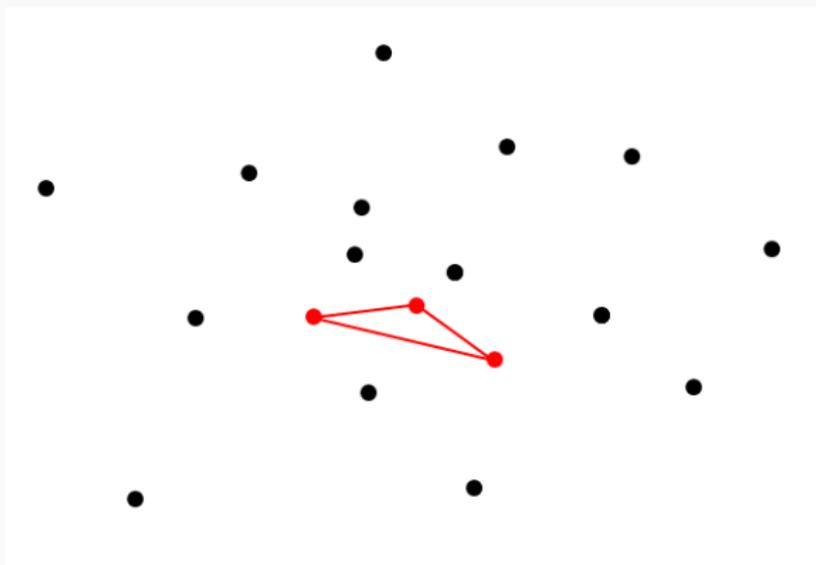
Action: when circle hits a point, $+1$ or -1 to the state.

Time complexity: $O(n^2 \log n)$

Problem 3: Duality

Find the smallest area triangle formed by using 3 of these points

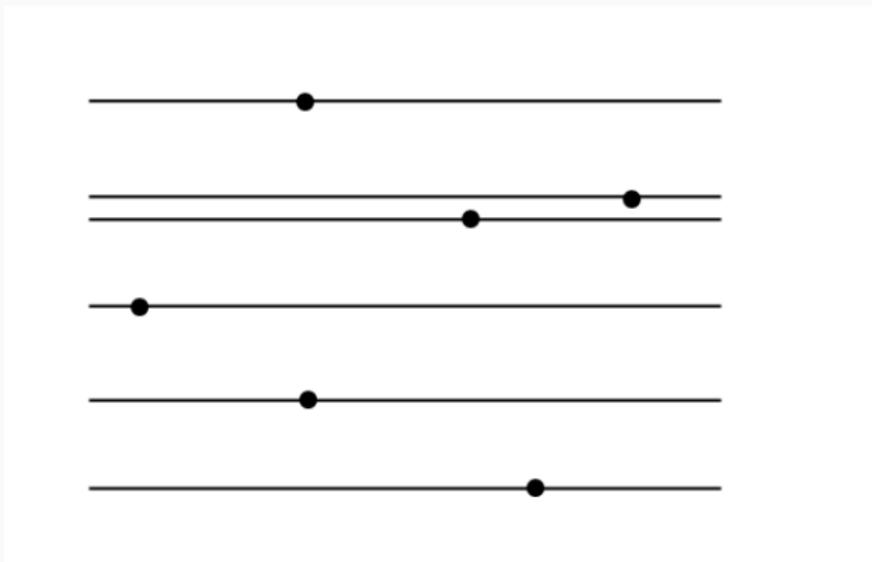
Hint: draw a set of parallel lines, one through every point, and rotate them all about their respective point at the same time.



Problem 3: Solution

Observation 1: only want to consider “adjacent points” in “vertical” ordering around any line segment “base”

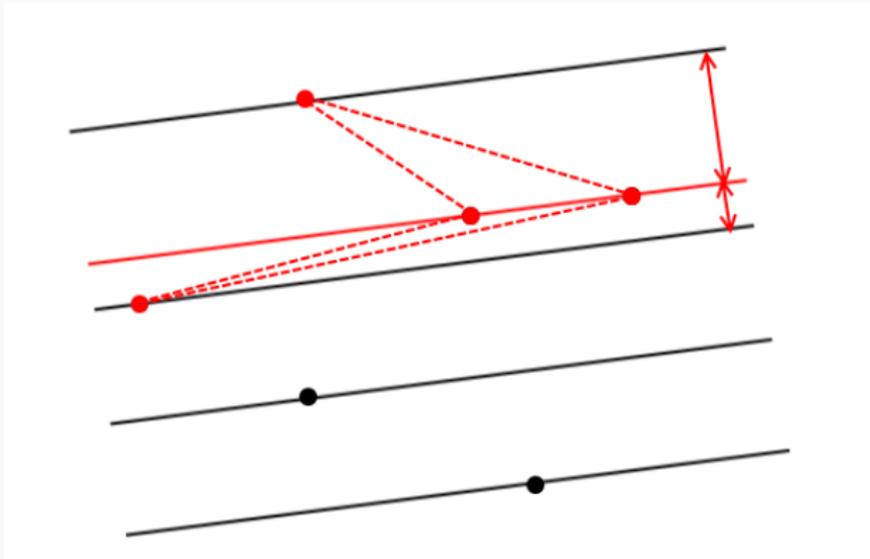
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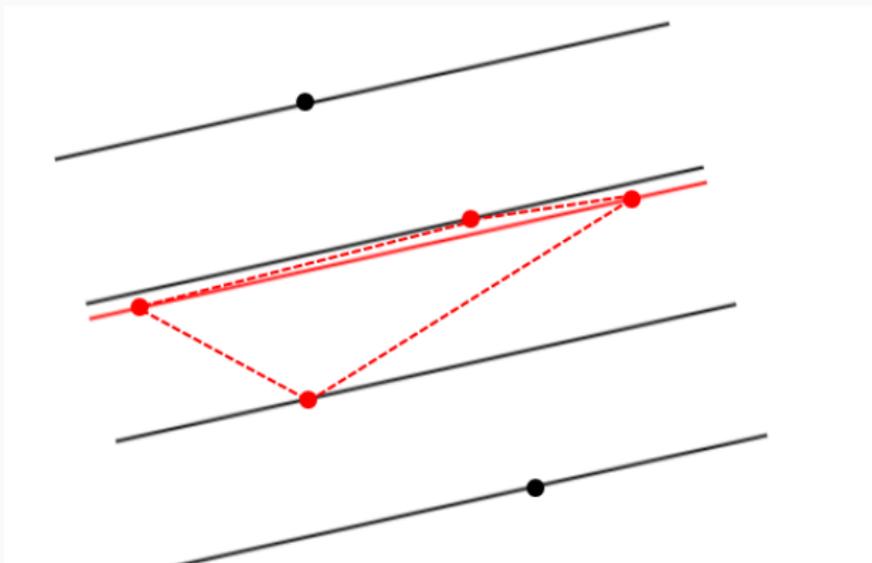
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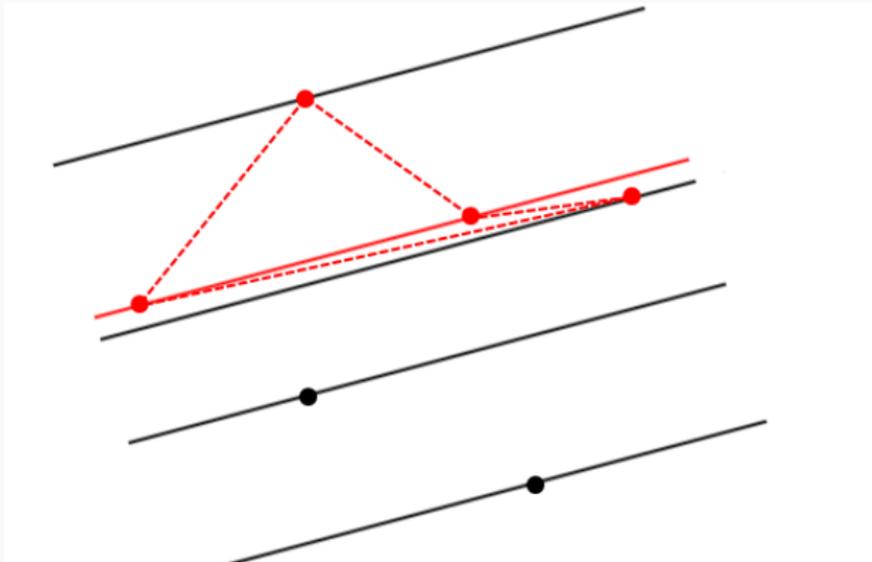
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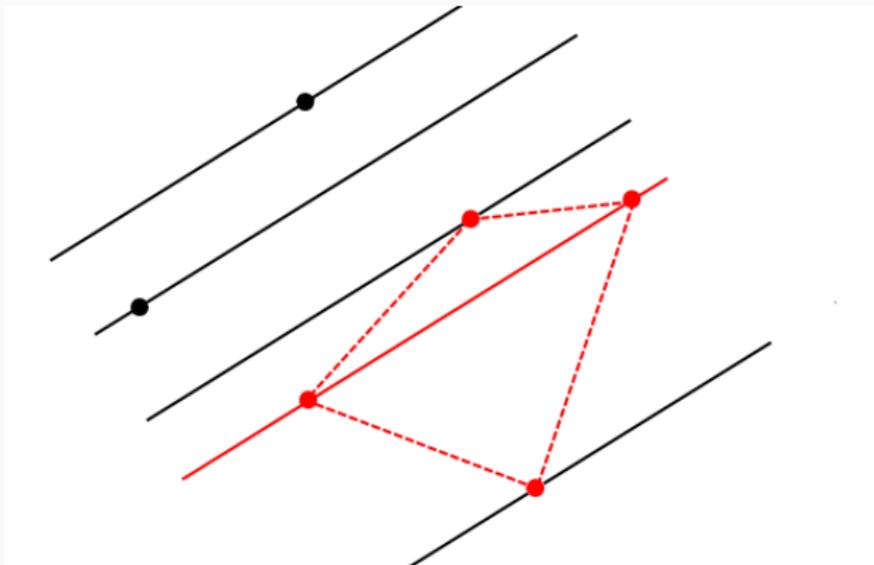
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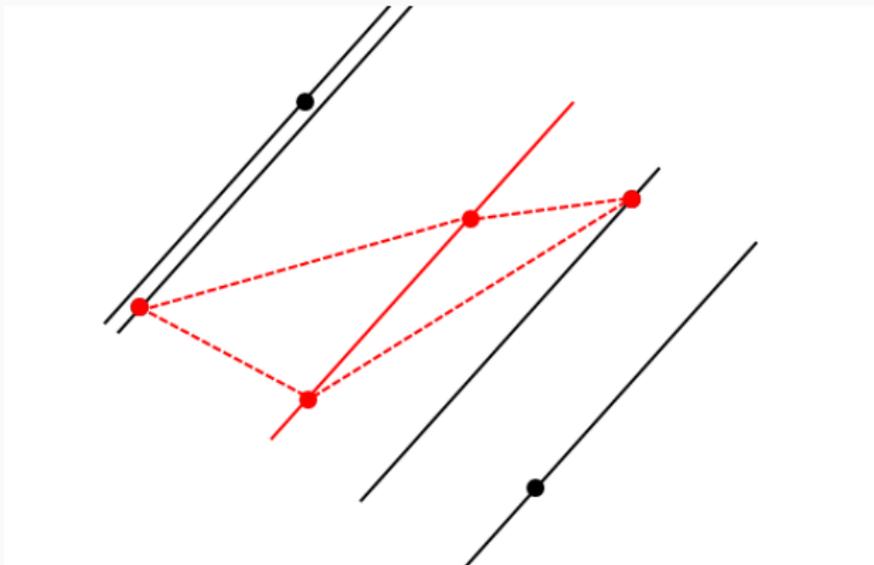
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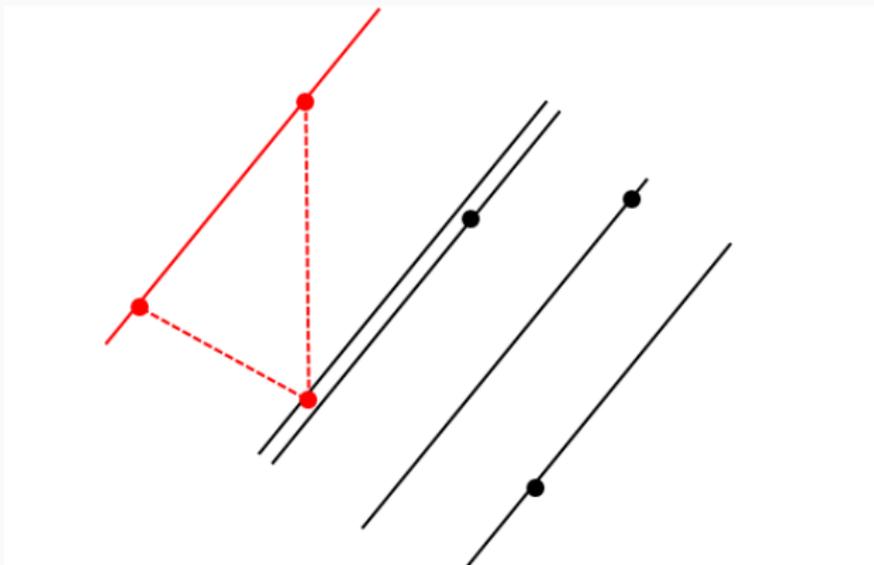
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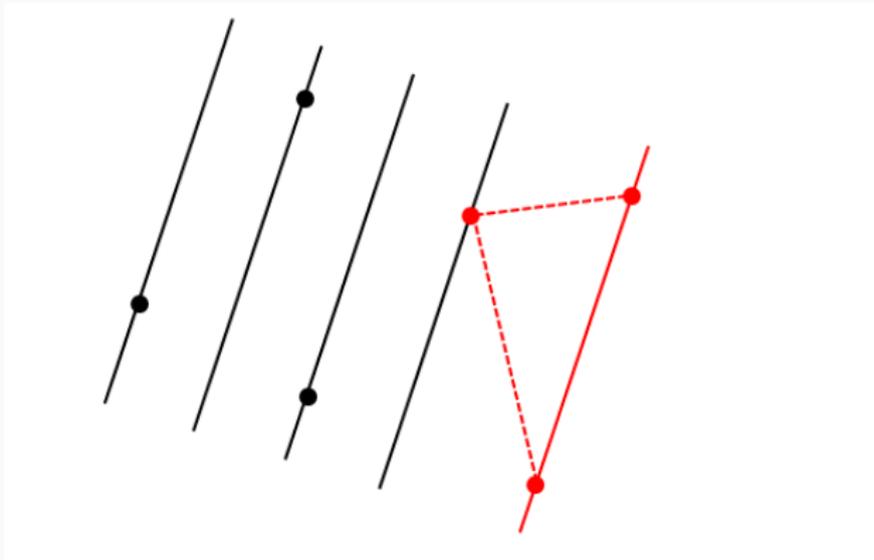
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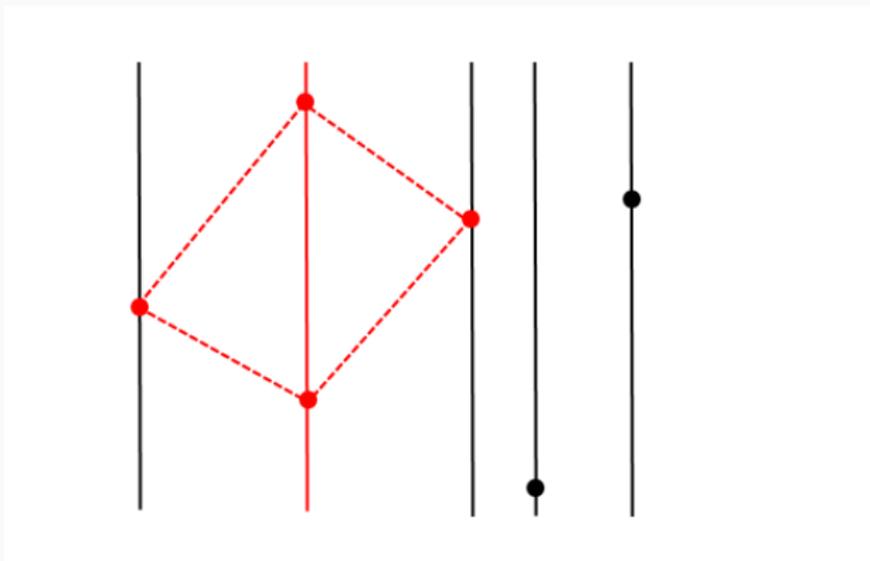
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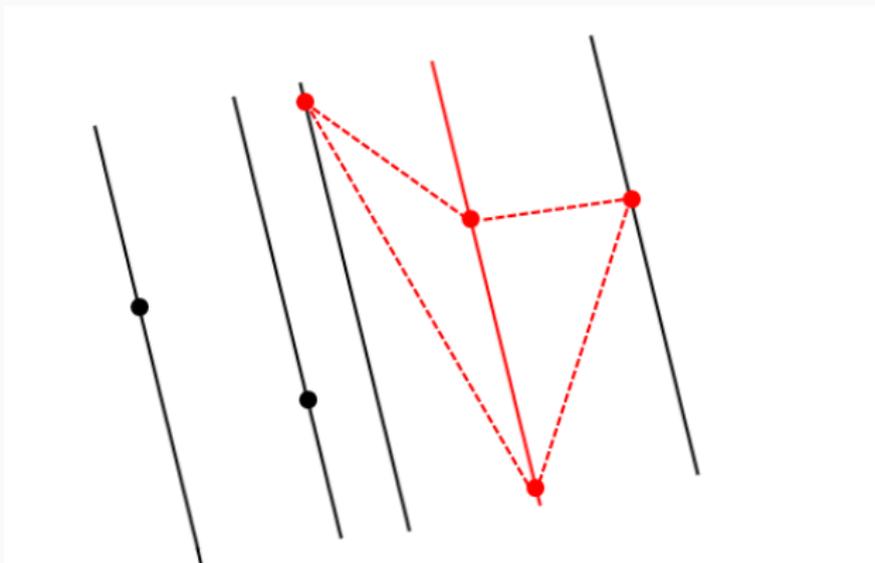
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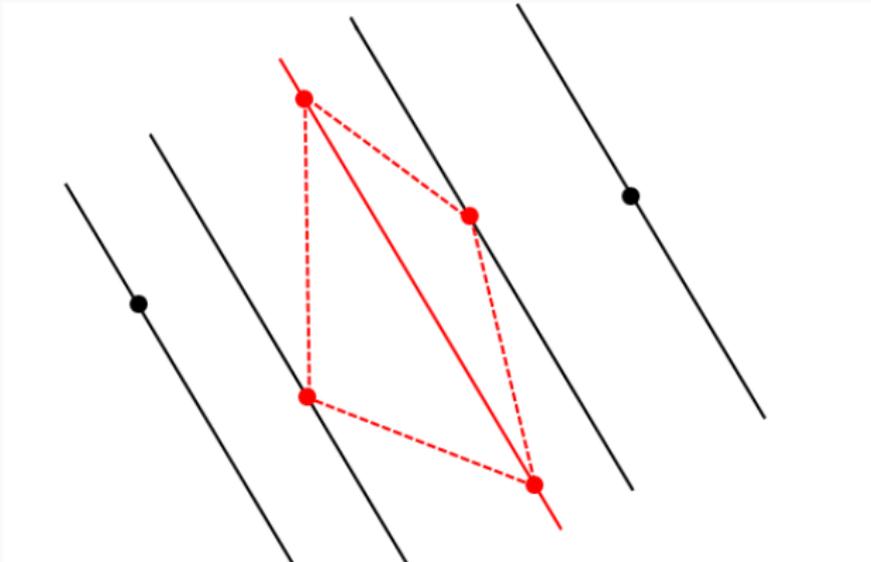
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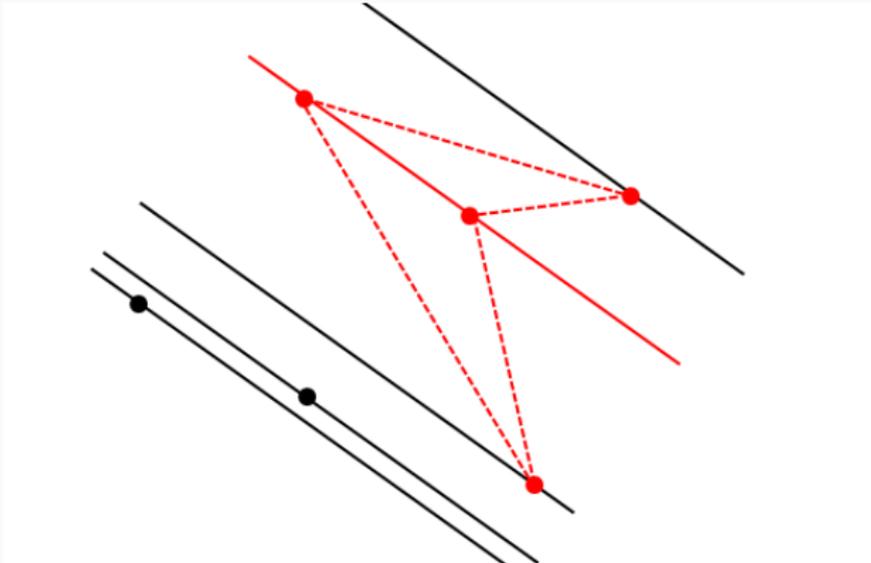
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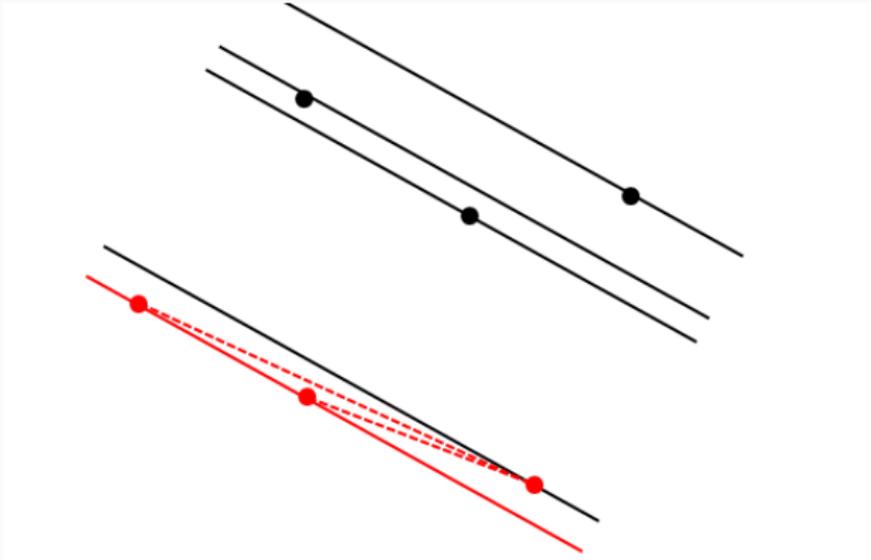
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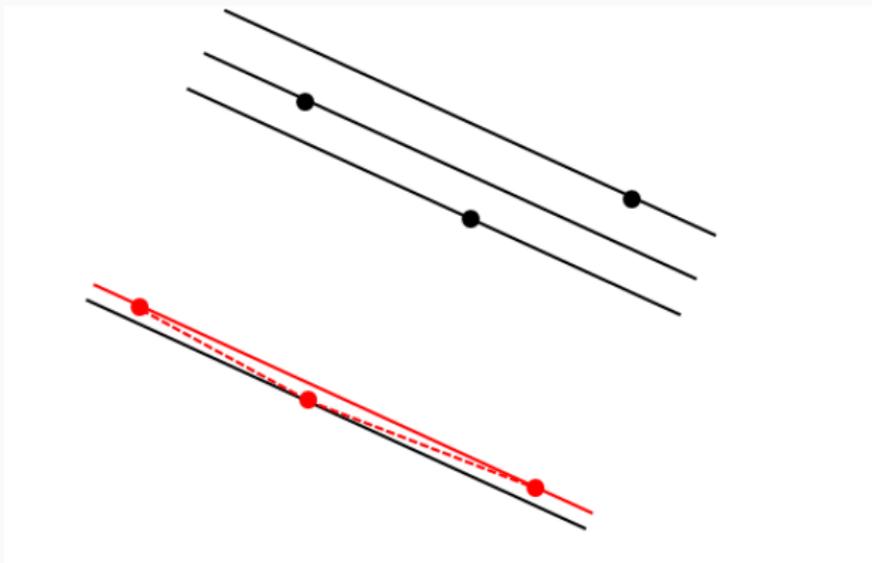
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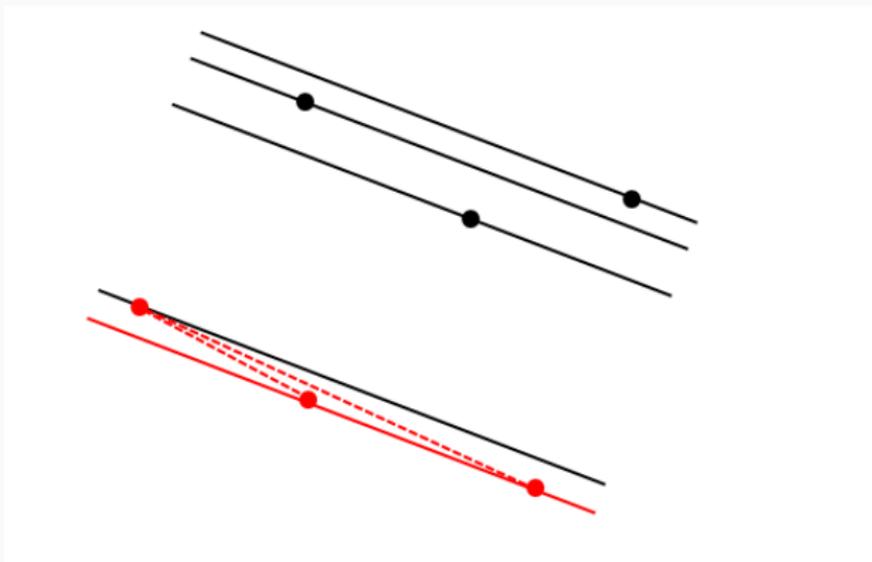
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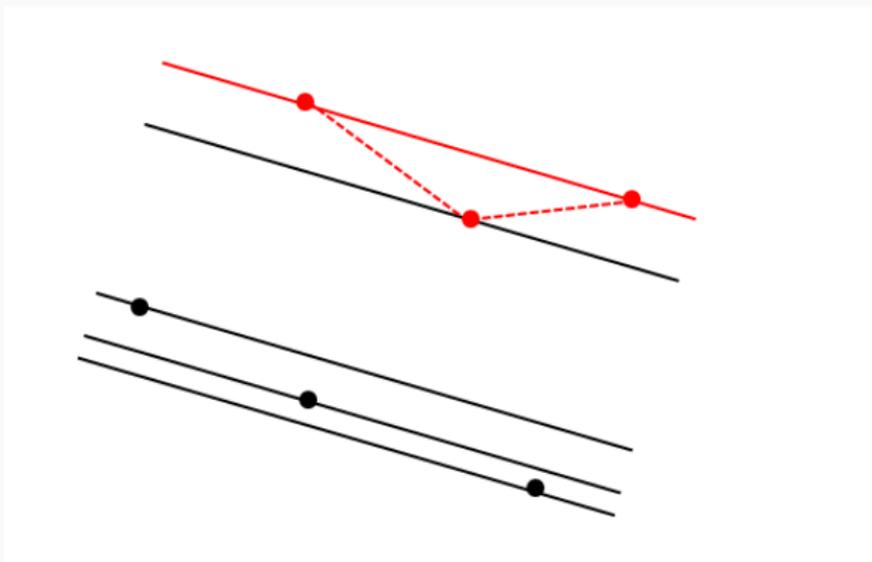
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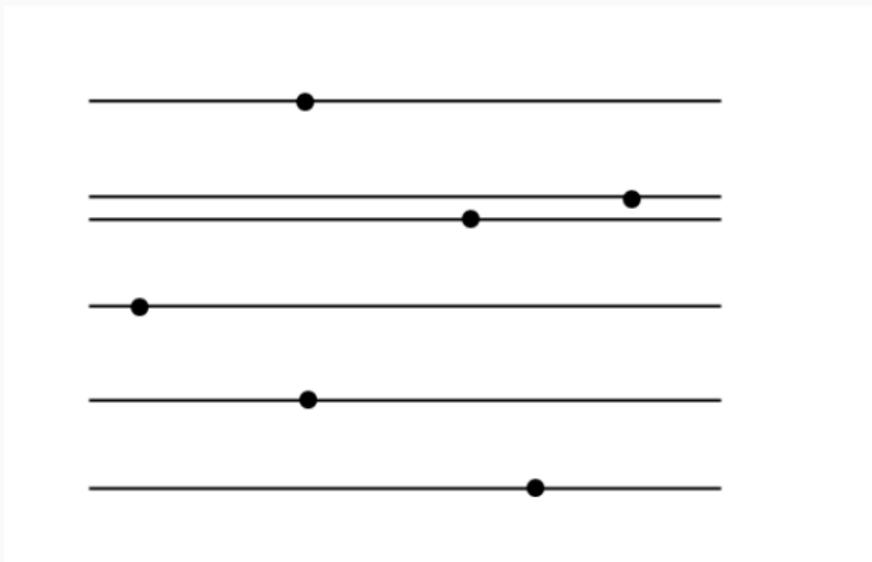
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Problem 3: Solution

State: ordering of points w.r.t. current set of parallel sweep lines

Events: when slope matches that of some lines segment
(i.e. when two of these sweep lines coincide)

Action:

- Flip the ordering of the two points whose sweep lines coincide
- Compute area of the two triangles formed using two “adjacent” points to the line segment, in the current ordering of points

Time complexity: $O(n^2 \log n)$

Duality

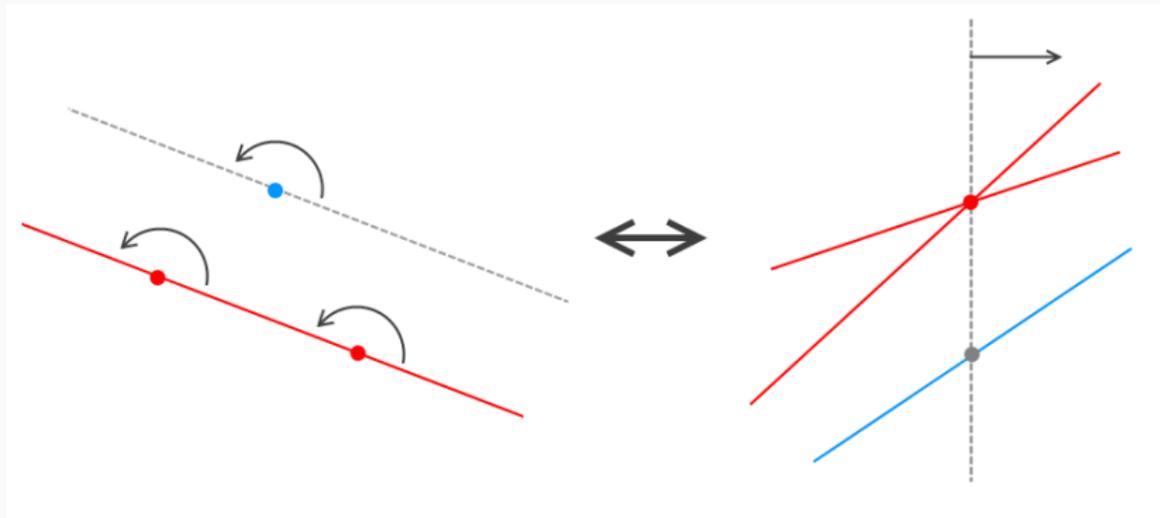
Map $(a, b) \mapsto y = ax - b$ then every point turns into a line, and every line turns into a point. This is the “dual space” representation.

Important properties

- Dual of dual of point is the same point
- Point above line \Leftrightarrow dual of point below dual of line
- 2 lines intersect at point \Leftrightarrow dual of point go thru dual of lines
- 3 points colinear \Leftrightarrow dual of 3 points intersect at same point
- Points CCW on upper hull \Leftrightarrow lines left to right on lower envelope

Smallest Triangle in Dual Space

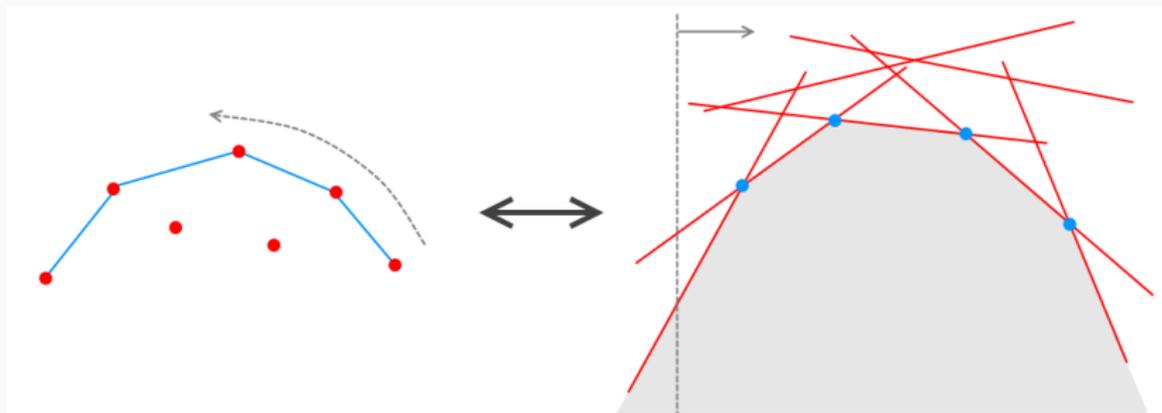
Point above line segment \Leftrightarrow Line below intersection of 2 lines.



Angular sweep in primal space \Leftrightarrow Line sweep in dual space

Convex Hull in Dual Space

Upper convex hull of points in CCW \Leftrightarrow lower envelope of lines in CW

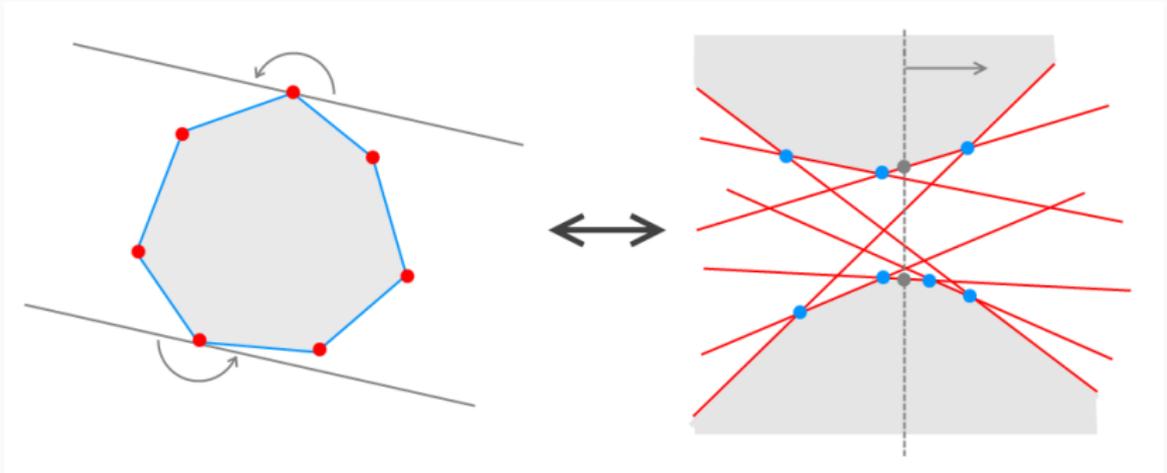


Gift wrapping in primal space \Leftrightarrow Line sweep in dual space, $O(n^2)$

Can you formulate an algorithm similar to Monotone Chain to construct lower envelope in $O(n \log n)$?

Rotating Calipers in Dual Space

Upper convex hull of points in CCW \Leftrightarrow lower envelope of lines in CW
Lower convex hull of points in CCW \Leftrightarrow upper envelope of lines in CW



Rotating calipers in primal space \Leftrightarrow Line sweep in dual space

One sweep to rule them all!

Other interesting algorithms / data structures

You are encouraged to read up on these, but we will not cover them:

- Delaunay triangulation & Voronoi diagram
 - MST on 2D, divide regions by nearest neighbor
 - They are duals of each other!
- Line arrangement & topological line sweep
 - Many applications, including visibility graph in $O(n^2)$
- K-d trees, 2D segment trees
 - 2D range query, nearest neighbor query
- Doubly Connected Edge List (DCEL)
 - Represent a planar graph / 3D polytope / line arrangement
- Dynamic convex hull / lower envelope

Divide and Conquer (with some applications to geometry)