

1. **(50 points)** For each language below, determine whether or not the language is context-free, and justify your answer. For example, if the language **is** context-free, you can give a CFG or PDA for the language. If it is not context-free, you can use the pumping lemma for CFLs and/or other properties of CFLs to prove it. You may use relationships to other languages that we have shown to be context-free or not, and you can use the results from previous parts of this problem when solving subsequent parts.
 - (a) **(10 points)** $A_1 = \{w \in \{a, b, c\}^* \mid \exists x, y \in \{a, b\}^*. (w = x\#y) \wedge (x \text{ is a substring of } y)\}$
 - (b) **(10 points)** $A_2 = \{w \in \{a, b, c\}^* \mid \exists i, j \in \mathbb{Z}^{\geq 0}. w = a^i b^i c^j\}$.
 - (c) **(10 points)** $A_3 = \{w \in \{a, b, c, d\}^* \mid \exists i, j \in \mathbb{Z}^{\geq 0}. w = a^i b^i c^j d^j\}$.
 - (d) **(20 points)** $A_6 = \overline{E}$, where $E = \{w \in \{a, b, c\}^* \mid \exists n \in \mathbb{Z}^{\geq 0}. w = a^n b^n c^n\}$.
2. **(20 points)** Let $G = (V, \Sigma, R, S)$ be a CFG in Chomsky normal form. Show that if G generates any string, w , with $|w| > 2^{|V|-1}$, then $L(G)$ is an infinite set.
3. **(20 points)** Draw the state diagram for a Turing machine that recognizes the language of all strings in $\{a\}^*$ for which the number of a's in the string is a perfect square.