

1. **(15 points)** Let $A_1 = \{M\#w \mid M \text{ halts after at most } |w|^{|w|} \text{ steps when run with input } w\}$. Show that language A is Turing decidable.
2. **(30 points)** Let $A_2 = \{M \mid \exists w. M\#w \in A_1\}$.
 - (a) **(15 points)** Show that language A_2 is not Turing decidable.
 - (b) **(15 points)** Show that language A_2 is Turing recognizable.
3. **(40 points)**
 - (a) **(10 points)** Show that the class of Turing decidable languages is closed under complement.
 - (b) **(10 points)** Show that the class of Turing decidable languages is closed under star.
 - (c) **(10 points)** Show that the class of Turing recognizable languages is not closed under complement.
 - (d) **(10 points)** Show that the class of Turing recognizable languages is closed under star.
4. **(30 points)** A *linear bounded automaton* (LBA) is a Turing Machine with a bounded tape; it cannot move its head past *either* end of the input string. For example, you can assume that the input string has the form $\vdash v \dashv$ where \vdash is a special left endmarker (that appears nowhere in v) and \dashv is a special right endmarker (that appears nowhere in v). All transitions from \vdash preserve the \vdash and move the head to the right. All transitions from \dashv preserve the \dashv and move the head to the left. Let

$$A_{LBA} = \{M\#w \mid M \text{ is an LBA that accepts } w\}.$$

A_{LBA} is Turing decidable (see Sipser Lemma 5.8). Thus, the halting problem for LBA's is Turing decidable as well.

Prove that there is some language, B such that B is Turing decidable but B is not accepted by any LBA. (**Hint:** use diagonalization.)