The actual exam will be open book, open notes, open homework and solutions, closed mouth. No calculators, laptop computers, PDA's etc.

1. (30 points): Each of the three languages described below is regular. Show this. For one language, you should construct a DFA; for another, you should construct an NFA; and for the remaining language, you should write a regular expression. You can choose which method you use with which language.
(a) $\left\{w \in\{0,1\}^{*} \mid w\right.$ is the binary representation of a number that is divisible by four. $\}$.
(b) $\left\{w \in\{0,1\}^{*} \mid w\right.$ is the binary representation of a number that is divisible by five. $\}$.
(c) $\left\{w \in\{0,1\}^{*} \mid \exists x, y, z .(w=x y z) \wedge(\# 1(y)=(\# 0(y)+3))\right\}$, where $\# 0(y)$ is the number of 0 's in y , and $\# 1(y)$ is the number of 1 's. In English, this says that $w$ contains a substring (i.e. $y$ ) that has three more 1's than 0's.

For parts (a) and (b), assume that the string is entered most significant bit first, i.e. 0101 represents the decimal value 5 .
2. (30 points): Let $A$ and $B$ be regular languages over some alphabet $\Sigma$. Let $C$ be the language:

$$
C=\left\{w \mid \exists x, y, z \in \Sigma^{*} .(w=x z) \wedge(x y z \in A) \wedge(y \in B)\right\}
$$

Show that $C$ is regular.
An acceptable answer will construct a DFA, NFA, or RE for $C$, or use closure properties that we've already shown. You should write a few sentences to explain why your construction is correct, but you don't need to write a full, formal proof.
3. (40 points): Let $\# 0(w)$ be the number of 0 's in $w$ and $\# 1(w)$ be the number of 1 's in $w$. For each language below, determine whether or not it is regular. Give a brief justification for each answer.
(a) (20 points): $A_{1}=\left\{w \in\{0,1\}^{*} \mid \# 0(w)-\# 1(w)\right.$ is divisible by 3$\}$
(b) (20 points): $A_{2}=\left\{w \in\{0,1\}^{*} \mid \# 0(w)-\# 1(w)<3\right\}$
(c) (10 points, Extra Credit): $A_{3}=\left\{w \in\{0,1\}^{*}| | \# 0(w)-\# 1(w) \mid>3\right\}$

