## Notes on the CpSc 418 Course Erlang Library

Draft version: no bug-bounties (yet) for minor spelling and grammar errors.

## 1 Distributed Lists

In Erlang, lists are the primary data structure. To write efficient parallel code, we need each process to store data locally as much a possible. This means we need to be able to take a large list, and distribute across the processes. The <u>workers</u> library provides functions <u>workers:update</u> <u>workers:retrieve</u> to do this. The function <u>misc:cut</u> in the <u>misc</u> module is helpful as well.

Let's say that L101 is a long list. For this example, L101 will be the list lists:seq(1,101), i.e. [1,2,ldots,101] – that's not very "long", but it means I can write out the example without using hundreds of pages to print the list. Let W8 be a worker-tree with eight processes. We can partition L101 into eight smaller lists using misc:cut:

```
misc:cut(L101, W8) ->
[ [1,2,3,4,5,6,7,8,9,10,11,12],
                                                % 12 elements
  [13,14,15,16,17,18,19,20,21,22,23,24],
                                                % 12 elements
  [25,26,27,28,29,30,31,32,33,34,35,36],
                                                % 12 elements
  [37,38,39,40,41,42,43,44,45,46,47,48,49],
                                                % 13 elements
  [50,51,52,53,54,55,56,57,58,59,60,61,62],
                                                % 13 elements
  [63,64,65,66,67,68,69,70,71,72,73,74,75],
                                                % 13 elements
                                                % 13 elements
  [76,77,78,79,80,81,82,83,84,85,86,87,88],
  [89,90,91,92,93,94,95,96,97,98,99,100,101]] % 13 elements
```

The first argument to <u>misc:cut</u> is the list to split up. The second argument tells <u>misc:cut</u> how many pieces to make. In a bit more detail,

- If the second argument to misc: cut(L, N) is an integer, then L is divided into N segments.
- If the second argument to misc: cut(L, W) is an list, then L is divided into length(W) segments.

Why have the second case? Because worker trees such as W8 are represented as lists. The length of this list is the number of workers. Of course, I'd prefer that *you* didn't depend on knowing the internal data structures of the <u>workers</u> and <u>wtree</u> modules. In the not-too-likely event that I were to change the representation of worker trees, I would revise <u>misc:cut</u> to work with the new representation – I promise.

Now that we have divided misc: cut into pieces, we need to send these pieces to the workers. The function workers: update does just this. The general form is

```
workersupdate(W, Key, ListOfValues)
```

```
where length(ListOfValues) matches the number of workers in W. For example, workersupdate(W8, my_list, misc:cut(L101, W8))
```

will divide misc:cut(L101) into eight pieces and send each process of W8 a separate piece. The workers in the tree are ordered from left-to-right, and the values of ListOfValues are distributed in that order. In other words, the leftmost worker process receives the first value of ListOfValues and so on.

What is the Key all about? The workers need a way to access these values. For reduce, scan, and other operations, you will provide functions for these worker processes to access. How do they get the list you just sent? You will note that the *Leaf* functions for <a href="wtree:reduce">wtree:scan</a> have one parameter, <a href="ProcState">ProcState</a> is an association list that pairs keys with values. The worker processes access values in <a href="ProcState">ProcState</a> using the functions <a href="wtree:get">wtree:get</a> and <a href="wtree:put">wtree:get</a> and <a href="wtree:put">workers:get</a> and <a href="wtree:put">workers:get</a> and <a href="wtree:put">workers:get</a> and <a href="wtree:put">workers:put</a>. In particular:

wtree:get(ProcState, Key) returns the value in ProcState associated with Key. If there is no value associated with Key, then the atom undefined is returned.

wtree:get(ProcState, Key, Default) returns the value in ProcState associated with Key. If there is no value associated with Key, then the value of the parameter Default is returned.

wtree:put (ProcState, Key, Value) associate Key with Value in ProcState. Wait a second - Erlang is functional! Right. wtree:put creates a new association list, let's call it ProcState2. ProcState2 has all of the {K, V} mappings that were in ProcState plus the mapping {Key, Value}. If the mapping {Key, OldValue} was already present in ProcState, then ProcState2 has all of the mappings of ProcState except for {Key, OldValue} and it has the mapping {Key, Value}. In other words, the mapping for {Key, OldValue} in ProcState is replaced with {Key, Value} in ProcState2.

## 2 Reduce and Scan

The functions <u>wtree:reduce</u> and <u>wtree:scan</u> make use of ProcState and that typically involves distributed lists as described above. First, let's look at <u>wtree:reduce</u>(W, Leaf, Combine, Root) where W is a worker-tree produced by <u>wtree:create</u>.

Leaf (ProcState) -> Value: ProcState is an association of keys to values for this process as described above. This is how the process finds the data that it is to work on. Value is the value that Leaf computes based on ProcState and returns for the Combine tree. We typically use an Erlang fun() expression to capture information about what key to use to access ProcState and any other details that Leaf may need. I should put an example here, but I'm running out of time.

Combine (Left, Right) -> Value: Left and Right are the values from the left and right subtees of this node. Combine them into a single value. Note that Left, Right, and Value should all be of the same "type". For example, if Left and Right are numbers and Value is a tuple, you're probably doing something wrong, and the call to Combine in the next level of the tree will probably crash. Note that Combine does not take ProcState as a parameter. Anything you need to know about ProcState should be included in the values returned by Leaf and Combine. However, if you're returning the entire segment of a distributed list for that leaf-node or subtree, you probably aren't doing the reduce correctly.

Root(Value0) -> Value1: again, no ProcState.

I'll continue this document with a description of <a href="wtree:scan">wtree:scan</a> and <a href="wtree:rlist">wtree:rlist</a>. The key points are that the Leaf2 function for <a href="wtree:scan">wtree:scan</a> should return an updated ProcState. The function <a href="wtree:rlist">wtree:rlist</a> produces a random list that is distributed across the workers of a worker tree.