

$$\begin{aligned}
 T(n) &\leq T(2^{\lceil \lg n \rceil}) = 2^{\lceil \lg n \rceil} + c 2^{\lceil \lg n \rceil} \lg 2^{\lceil \lg n \rceil} \\
 &\leq 2^{\lg n + 1} + c 2^{\lg n + 1} \lg 2^{\lg n + 1} \\
 &= 2n + c 2n (\lg n + 1) \\
 &= 2n + c 2n \lg n + c 2n \\
 &= c 2n \lg n + \underline{2n(c+1)} \\
 &\leq c 2n \lg n + \underline{2n \lg n (c+1)} \quad \text{if } n \geq 2 \\
 &\quad \text{then } n \lg n \geq n \\
 &= n \lg n \left(\frac{c 2 + 2c + 2}{d} \right) \\
 &\quad d = 4c + 2
 \end{aligned}$$

REDUCE

SORTING by COMP ($\Omega(n \lg n)$)
 ↓
 reduce in $o(n \lg n)$

SKYLING

shows SKYLING $\in \Omega(n \lg n)$

Let B

INPUT: $A = [a_1, a_2, \dots, a_k]$

↓

$B = [(a_1, l, a_1 + 1), (a_2, l, a_2 + 1), \dots, (a_k, l, a_k + 1)]$

①

OUTPUT: $(x_1, 1), (x_2, 0), \dots, (x_i, 1), (x_{i+1}, 0)$

\Downarrow

$x_1, x_2, \dots, x_{i-1}, \dots, x_i, x_{i+1}, \dots, x_{i+1}$

So: SKYLINE ^{PROBLEM} $\in \Omega(n \lg n)$ and $O(n \lg n)$
 SKYLINE_ALG $\in \Theta(n \lg n)$

DANGERS OF GUESS-AND-TEST

GUESS INCORRECTLY $T(n) \in O(n)$

$$T(1) = 1$$

$$T(n) = T(\lfloor \frac{n}{2} \rfloor) + T(\lceil \frac{n}{2} \rceil) + \Theta(1)$$

Guess: $T(n) = cn$

Ind Hyp: ~~$T(k) = ck$~~ $\forall i < k, T(i) = ci$

Ind Step: Show $T(k+1) = c(k+1)$ \Leftarrow

$$T(k+1) = T(\lfloor \frac{k+1}{2} \rfloor) + T(\lceil \frac{k+1}{2} \rceil) + k+1$$

$$= c \lfloor \frac{k+1}{2} \rfloor + c \lceil \frac{k+1}{2} \rceil + k+1 \quad \text{by IH}$$

$$= c(k+1) + k+1$$

$$\textcircled{2} = (c+1)(k+1) \leftarrow \text{let } d = c+1$$

CANNOT RE-SELECT YOUR CONSTANT

For ints n :
 $\lfloor \frac{n}{2} \rfloor + \lceil \frac{n}{2} \rceil = n$

FAILED PROOF \rightarrow

$$T(1) = 1$$

$$T(n) = T(\lceil \frac{n}{2} \rceil) + T(\lfloor \frac{n}{2} \rfloor) + 1$$

GUESS + TEST $T(n) \in O(n)$

WARNING #2: not all $O(n)$ functions look like cn .

GUESS: $T(n) = cn$

IH: $T(i) = ci$ for all $i < k$

IS: Show $T(k) = ck - 1$

$$\begin{aligned} T(k) &= T(\lceil \frac{k}{2} \rceil) + T(\lfloor \frac{k}{2} \rfloor) + 1 \\ &= c \lceil \frac{k}{2} \rceil + c \lfloor \frac{k}{2} \rfloor + 1 \quad \text{by IH} \\ &= ck - 1 \end{aligned}$$

TRY $T(n) = cn - 1$

BC: $T(1) = c \cdot 1 - 1 = 2 \cdot 1 - 1 = 1 \checkmark$

NEED $T(1) = 1$

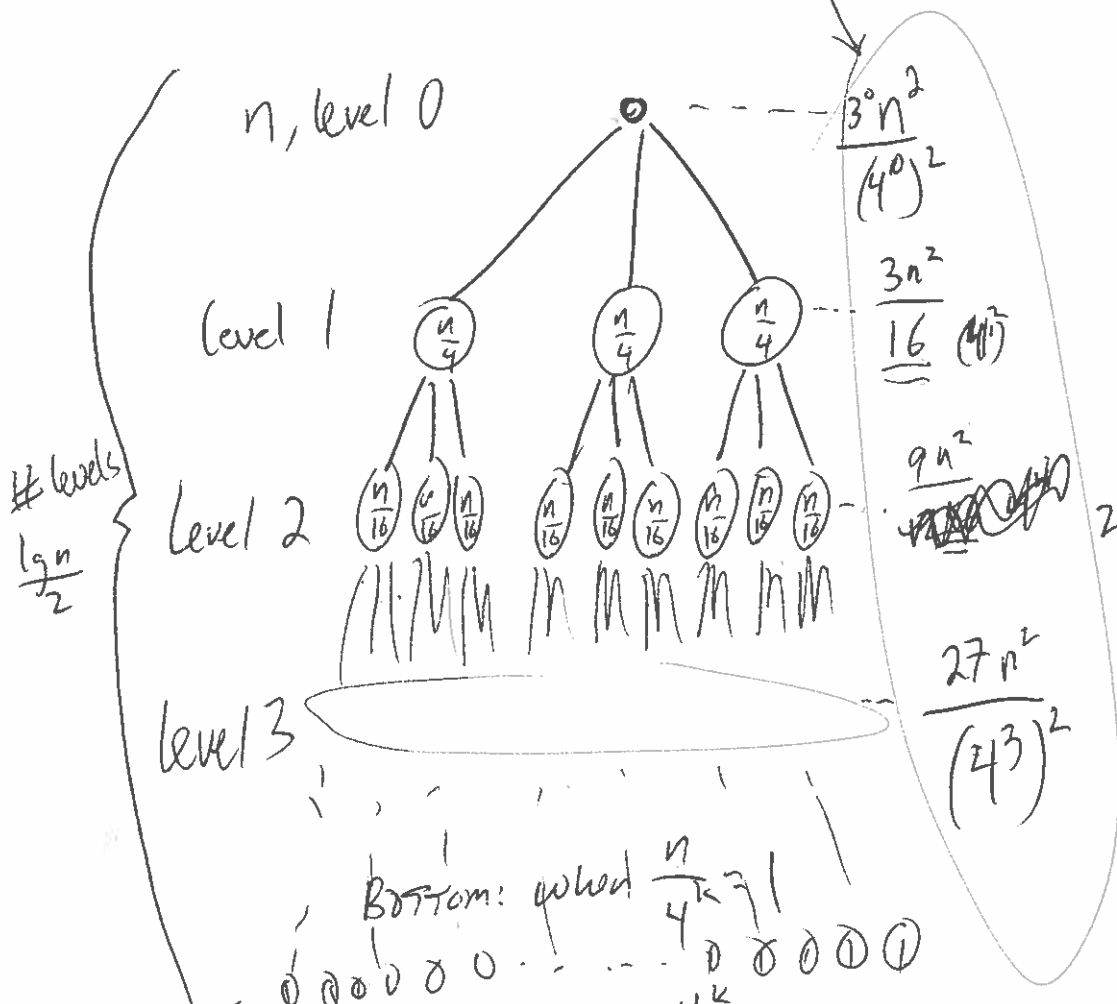
Let $c=2$

$T(1) = 1$

$\in O(n^2)$ by below $2\phi/4/21$

$T(n) = 3T(\frac{n}{4}) + n^2 \in \Omega(n^2)$ clearly

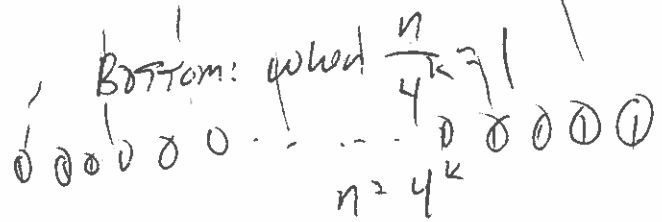
Assume n is a power of 4



level k
 $\frac{3^k n^2}{(4^k)^2} = \frac{3^k n^2}{(4^2)^k} = \left(\frac{3}{16}\right)^k n^2$

$0 < a < 1$

$\sum_{i=0}^{\infty} a^k = \frac{1}{1-a}$



$3^{\frac{\lg n}{2}}$ nodes
 $\lg n = \lg 4^k = k \lg 4 = 2k$
 $k = \frac{\lg n}{2}$

$T(n) = \left(\sum_{i=0}^{k-1} \left(\frac{3}{16}\right)^k n^2\right) + 3^{\frac{\lg n}{2}}$

$= n^2 \left(\sum_{i=0}^{k-1} \left(\frac{3}{16}\right)^k\right) + 3^{\frac{\lg n}{2}}$
 $\leq n^2 \left(\sum_{i=0}^{\infty} \left(\frac{3}{16}\right)^k\right) + 3^{\frac{\lg n}{2}} = n^2 \left(\frac{1}{1-\frac{3}{16}}\right) + 3^{\frac{\lg n}{2}}$
 $= n^2 C + 4^{\frac{\lg n}{2}} = n^2 C + 2^{\lg n} = n^2 C + 2^{\frac{\lg n}{2}}$