

1. Let  $F[i, X]$  be the value of the most valuable combination of items of weight at most  $X$  from items  $1, 2, \dots, i$ . The optimal solution to obtain value  $F[i, X]$  either (1) includes item  $i$  (if it fits) and an optimal solution using items  $1, 2, \dots, i - 1$  with knapsack capacity  $X - w_i$ , or (2) excludes item  $i$  and is an optimal solution using items  $1, 2, \dots, i - 1$  with knapsack capacity  $X$ .

$$F[i, X] = \begin{cases} \max\{F[i - 1, X - w_i] + v_i, F[i - 1, X]\} & \text{if } w_i \leq X \\ F[i - 1, X] & \text{otherwise} \end{cases}$$

$F[0, X] = 0$  for integers  $X$  between 0 and  $W$ .  $F[i, 0] = 0$  for integers  $i$  between 0 and  $n$ .

Knapsack( $v_1, \dots, v_n, w_1, \dots, w_n, W$ )

for  $X := 0$  to  $W$   $F[0, X] := 0$

for  $i := 0$  to  $n$   $F[i, 0] := 0$

for  $i := 1$  to  $n$

for  $X := 0$  to  $W$

$F[i, X] := F[i - 1, X]$

if  $w_i \leq X$  and  $F[i, X] < F[i - 1, X - w_i] + v_i$  then

$F[i, X] := F[i - 1, X - w_i] + v_i$

$X := W$

for  $i := n$  to 1

if  $F[i, X] > F[i - 1, X]$  then

print  $i$

$X := X - w_i$

2. For  $i \leq j$ , let  $A[i, j]$  be the length of a longest palindromic subsequence of  $X[i \dots j]$ . If  $X[i] = X[j]$  then such a palindromic subsequence is  $X[i]$  followed by a longest palindromic subsequence of  $X[i + 1 \dots j - 1]$  followed by  $X[j]$ . If  $X[i] \neq X[j]$  then such a palindromic subsequence is either a longest palindromic subsequence of  $X[i + 1 \dots j]$  or of  $X[i \dots j - 1]$  since both  $X[i]$  and  $X[j]$  cannot contribute. For  $i < j$ ,

$$A[i, j] = \begin{cases} A[i + 1, j - 1] + 2 & \text{if } X[i] = X[j] \\ \max\{A[i + 1, j], A[i, j - 1]\} & \text{otherwise} \end{cases}$$

$A[i, i] = 1$  for all  $i$  between 1 and  $n$ .  $A[i, i - 1] = 0$  for all  $i$  between 2 and  $n$ .

PalindromicSubseq( $X[1 \dots n]$ )

for  $i := 1$  to  $n$   $A[i, i] := 1$

for  $i := 2$  to  $n$   $A[i, i - 1] := 0$

for  $d := 1$  to  $n - 1$

for  $i := 1$  to  $n - d$

$j := i + d$

if  $X[i] = X[j]$  then  $A[i, j] := A[i + 1, j - 1] + 2$

else  $A[i, j] := \max\{A[i + 1, j], A[i, j - 1]\}$

return  $A[1, n]$