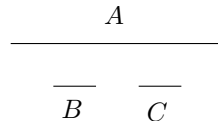


1. Consider the three jobs $A = (0, 5)$, $B = (1, 2)$, and $C = (3, 4)$.



The optimal solution is $\{A\}$ with length 5.

- (a) This version considers the jobs in the order A, C, B . It calculates $\text{last}[1, 2, 3] = [0, 0, 2]$ and $L[0, 1, 2, 3] = [0, 5, 5, 6]$, returning 6.
- (b) This version considers the jobs in the order A, B, C . It calculates $\text{last}[1, 2, 3] = [0, 0, 2]$ and $L[0, 1, 2, 3] = [0, 5, 5, 6]$, returning 6.
- (c) works. See problem set 6 explanation.
2. The cheapest way to get from city 1 to city n must stop at some last city k and then go directly from city k to city n . Let $F[i]$ be the cheapest rental cost to canoe from city 1 to city i .

$$F[i] = \min_{1 \leq k < i} \{F[k] + p_{ki}\} \quad \text{for } i = 2, 3, \dots, n$$

$$F[1] = 0$$

The actual algorithm is a simple for-loop:

```
CheapestCost( $p_{ij}$  for  $1 \leq i < j \leq n$ )
   $F[1] = 0$ 
  For  $i = 2$  to  $n$ 
     $F[i] = \min_{1 \leq k < i} \{F[k] + p_{ki}\}$ 
  Return  $F[n]$ 
```

3. The maximum sum contiguous subsequence that ends at position j is either the empty subsequence or the maximum sum contiguous subsequence that ends at position $j - 1$ plus a_j . Let $S[j]$ be the sum of the maximum sum contiguous subsequence that ends at position j .

$$S[j] = \max\{0, S[j - 1] + a_j\} \quad \text{for } j = 1, 2, \dots, n$$

$$S[0] = 0$$

```
MaxSumSubsequence( $a_1, a_2, \dots, a_n$ )
   $S[0] = 0, m = 0$ 
  For  $j = 1$  to  $n$ 
     $S[j] = \max\{0, S[j - 1] + a_j\}$ 
    If  $S[j] > S[m]$  then  $m = j$ 
   $j = m$ 
  While  $j > 0$  and  $S[j] = S[j - 1] + a_j$ 
     $j = j - 1$ 
  Return  $a_{j+1}, a_{j+2}, \dots, a_m$ 
```