

## Unit Outline

### Unit #4: Sorting

CPSC 221: Basic Algorithms and Data Structures

Anthony Estey, Ed Knorr, and Mehrdad Oveis

2016W2

- ▶ Comparing Sorting Algorithms
- ▶ Heapsort
- ▶ Mergesort
- ▶ Quicksort
- ▶ More Comparisons
- ▶ Complexity of Sorting

2 / 33

## Learning Goals

- ▶ Describe, apply, and compare various sorting algorithms.
- ▶ Analyze the complexity of these sorting algorithms.
- ▶ Explain the difference between the complexity of a problem (sorting) and the complexity of a particular algorithm for solving that problem (e.g., Insertion Sort).

3 / 33

## How to Measure Sorting Algorithms

- ▶ Computational complexity (a.k.a. runtime)
  - ▶ Worst case
  - ▶ Average case
  - ▶ Best caseHow often is the input sorted, reverse sorted, or “almost” sorted ( $k$  swaps from sorted where  $k \ll n$ )?
- ▶ Stability: What happens to elements with identical keys? Why do we care?
- ▶ Memory Usage: How much extra memory is used?

4 / 33



## Mergesort

## Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

Mergesort is a “divide and conquer” algorithm.

1. If the array has 0 or 1 elements, it's sorted. Stop.
2. Split the array into two approximately equal-sized halves.
3. Sort each half recursively (using Mergesort).
4. Merge the sorted halves to produce one sorted result:
  - ▶ Consider the two halves to be queues.
  - ▶ Repeatedly dequeue the smaller of the two front elements (or dequeue the only front element if one queue is empty) and add it to the result.

9 / 33

10 / 33

## Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

3	-4	7	5
---	----	---	---

9	6	2	1
---	---	---	---

## Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

3	-4	7	5
---	----	---	---

3	-4
---	----

7	5
---	---

10 / 33

10 / 33

### Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

3	-4	7	5
---	----	---	---

3	-4
---	----

3
---

-4
----

7	5
---	---

9	6	2	1
---	---	---	---

### Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

3	-4	7	5
---	----	---	---

3	-4
---	----

3
---

-4
----

7	5
---	---

-4	3	*
----	---	---

### Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

3	-4	7	5
---	----	---	---

3	-4
---	----

3
---

-4
----

7	5
---	---

9	6	2	1
---	---	---	---

### Mergesort Example

3	-4	7	5	9	6	2	1
---	----	---	---	---	---	---	---

3	-4	7	5
---	----	---	---

3	-4
---	----

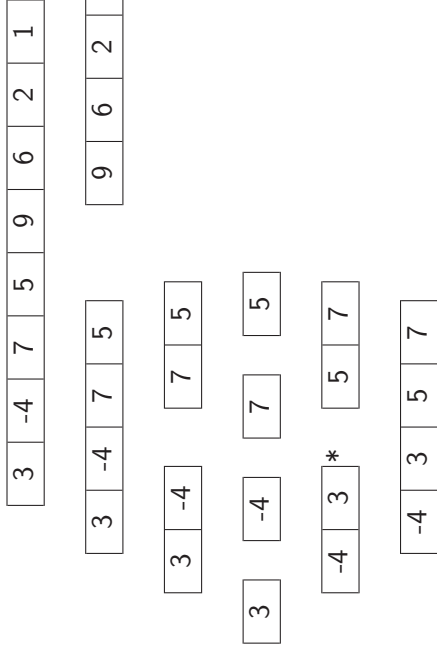
3
---

-4
----

7	5
---	---

-4	3	*
----	---	---

## Mergesort Example



10/33

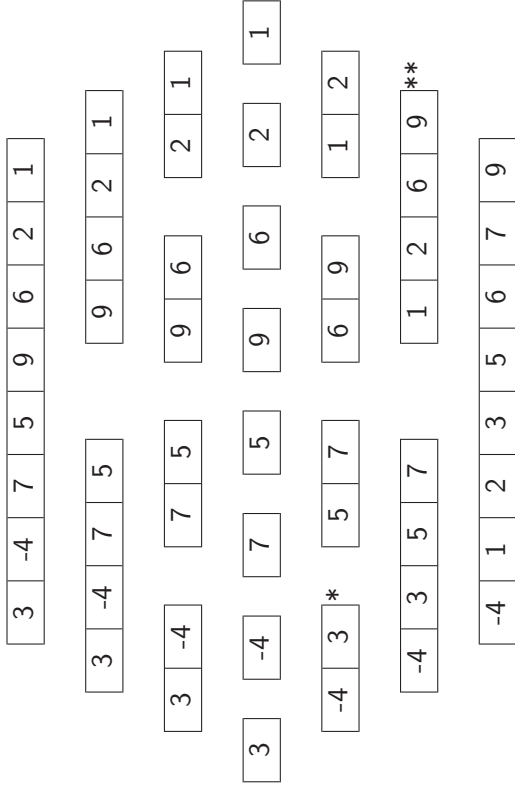
## Mergesort Code

```
void msort(int x[], int lo, int hi, int tmp[]) {
    if (lo >= hi) return;
    int mid = (lo+hi)/2;
    msort(x, lo, mid, tmp);
    msort(x, mid+1, hi, tmp);
    merge(x, lo, mid, hi, tmp);
}

void mergesort(int x[], int n) {
    int *tmp = new int[n];
    msort(x, 0, n-1, tmp);
    delete[] tmp;
}
```

11/33

## Mergesort Example



10/33

## Merge Code

```
void merge(int x[], int lo, int mid, int hi, int tmp[]) {
    int a = lo, b = mid+1;
    for( int k = lo; k <= hi; k++ ) {
        // What's the loop invariant, at this point?
        if( a <= mid && (b > hi || x[a] < x[b]) )
            tmp[k] = x[a++];
        else tmp[k] = x[b++];
    }
    for( int k = lo; k <= hi; k++ )
        x[k] = tmp[k];
}
```

12/33

```
merge( x, 0, 0, 1, tmp ); // step *
```

x :	3	-4	7	5	9	6	2	1
tmp :	-4	3	?	?	?	?	?	?
x :	-4	3	7	5	9	6	2	1

```
merge( x, 4, 5, 7, tmp ); // step **
```

x :	-4	3	5	7	6	9	1	2
tmp :	?	?	?	?	1	2	6	9
x :	-4	3	5	7	1	2	6	9

```
merge( x, 0, 3, 7, tmp ); // is the final step
```

## Stable:

Dequeue from the left queue if the two front elements are equal.

## Memory:

This is not easy to implement without using  $\Omega(n)$  extra space; so, it is not viewed as an in-place sort. Plus there's the cost of the call stack ( $\Omega(\log n)$ ).

# Quicksort (C.A.R. Hoare 1961)

In practice, this is one of the fastest sorting algorithms.

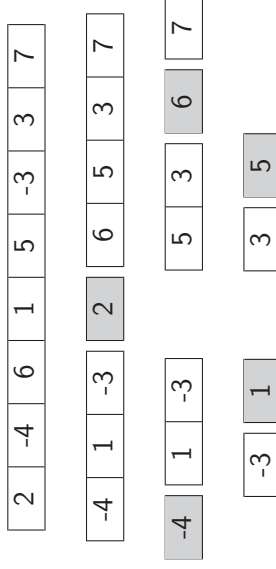
1. Pick a **pivot**
2. Reorder the array such that all elements  $<$  pivot are to its left, and all elements  $\geq$  pivot are to its right.
3. Recursively sort each partition.

2	-4	6	1	5	-3	3	7
---	----	---	---	---	----	---	---

-4	1	-3	2	6	5	3	7
left partition			pivot	right partition			

What's the base case?

# Quicksort: Visually





## Quicksort: Running Time

The running time is proportional to number of comparisons; so, let's count comparisons.

1. Pick a pivot.  
**Zero comparisons**
  2. Reorder (partition) the array around the pivot value.  
Quicksort compares each element to the pivot.  
 **$n - 1$  comparisons**
  3. Recursively sort each partition.  
**The number of comparisons depends on the size of the partitions.**
- ▶ If the partitions have size  $n/2$  (or any constant fraction of  $n$ ), the runtime is  $\Theta(n \log n)$  (like Mergesort).
  - ▶ In the worst case, however, we might create partitions with sizes 0 and  $n - 1$ . When might this occur?

21 / 33

## Quicksort: Worst Case

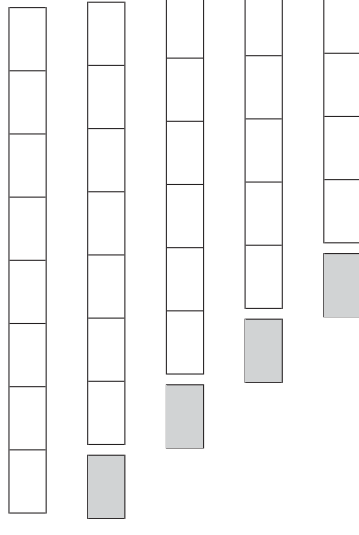
If this happens at every partition...  
Quicksort makes  $n - 1$  comparisons in the first partition and recurses on a problem of size 0 and size  $n - 1$ :

$$\begin{aligned}
 T(n) &= (n - 1) + T(0) + T(n - 1) = (n - 1) + T(n - 1) \\
 &= (n - 1) + (n - 2) + T(n - 2) \\
 &\vdots \\
 &= \sum_{i=1}^{n-1} i = (n - 1)(n - 2)/2
 \end{aligned}$$

This is  $\Theta(n^2)$  comparisons.

23 / 33

## Quicksort: Visually – the Worst Case



22 / 33

## Quicksort: Average Case (Intuition)

- ▶ On an average input (i.e., random order of  $n$  items), our chosen pivot is equally likely to be the  $i$ th smallest for any  $i = 1, 2, \dots, n$ .
- ▶ With probability  $1/2$ , our pivot will be from the middle  $n/2$  elements – a good pivot.



- ▶ Any good pivot creates two partitions of size at most  $3n/4$ .
- ▶ We expect to pick one good pivot every two tries.
- ▶ Expected number of splits is at most  $2 \log_{4/3} n \in O(\log n)$ .
- ▶  $O(n \log n)$  total comparisons. **True, but this intuition is not a proof.**

24 / 33



## Quicksort: Stability & Memory

### Stable:

Quicksort can be made stable – most easily by using more memory.

### Memory:

In-place sort

25 / 33

## A Comparison of Quicksort, Mergesort, Heapsort, and Insertion Sort

### Running Time:

	$\Theta(n)$	$\Theta(n \log n)$	$\Theta(n^2)$
Best case:	Insert	Quick, Merge, Heap	
Average case:		Quick, Merge, Heap	Insert
Worst case:		Merge, Heap	Quick, Insert
“Real” data:		Quick < Merge < Heap < Insert	

Some Quicksort/Mergesort implementations use Insertion Sort on small arrays (base cases).

Some results depend on the implementation. For example, an initial check whether the last element of the left subarray is less than the first of the right can make Mergesort’s best case linear.

27 / 33

## Comparison of Running Times for 100 Samples

n	Insertion		Heap		Merge		Quick	
	avg	max	avg	max	avg	max	avg	max
100,000	11.20s	16.37s	0.04s	0.08s	0.03s	0.04s	0.02s	0.04s
200,000	36.97s	60.01s	0.08s	0.16s	0.06s	0.11s	0.06s	0.16s
400,000	172.36s	505.38s	0.56s	1.74s	0.54s	0.91s	0.46s	0.69s
800,000			0.37s	0.83s	0.21s	0.35s	0.19s	0.32s
1,600,000			0.93s	1.77s	0.52s	1.12s	0.44s	0.78s
3,200,000			2.07s	3.04s	1.01s	1.95s	0.91s	1.44s
6,400,000			4.76s	7.54s	2.18s	3.88s	1.97s	3.45s
12,800,000			10.65s	12.38s	4.56s	7.01s	4.13s	5.94s

The code is from the lecture notes and labs, but it is not optimized.

26 / 33

## A Comparison of Quicksort, Mergesort, Heapsort, and Insertion Sort (cont.)

### Stability:

#### Stable (easy):

Insert, Merge (we prefer the left of the two sorted subarrays when encountering ties)

#### Stable (with effort):

Quick  
Heap

#### Unstable:

### Memory Use:

- ▶ Insert, Heap, Quick < Merge

28 / 33

## Complexity of the Sorting Problem

The **complexity** of a problem is the complexity of the best algorithm for that problem.

How powerful is our computer?

We'll only consider **comparison-based** algorithms.

They can compare two array elements in constant time.

They cannot manipulate array elements in any other way.

For example, they cannot assume that the elements are numbers and perform arithmetic operations (like division) on them.

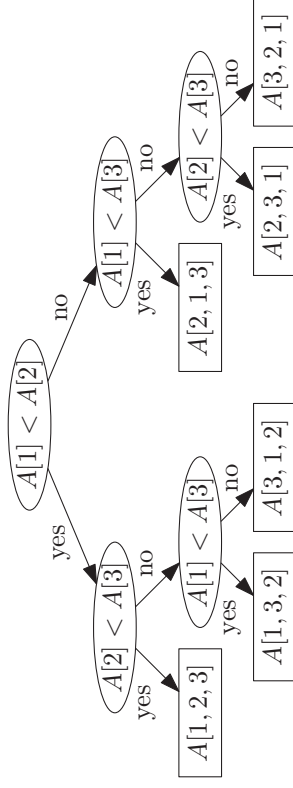
Insertion Sort, Heapsort, Mergesort, and Quicksort are comparison-based.

Radix sort is not.

29 / 33

## Complexity of the Sorting Problem

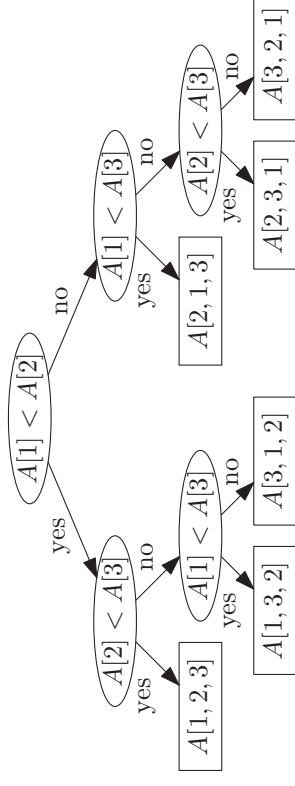
- ▶ This is the decision tree representation of Insertion Sort on inputs of size  $n = 3$ .
- ▶ Each leaf outputs the input array in some particular order. For example,  $A[3, 1, 2]$  means output  $A[3]$ ,  $A[1]$ ,  $A[2]$ .



31 / 33

## Comparison-Based Algorithms Using a Decision Tree Model

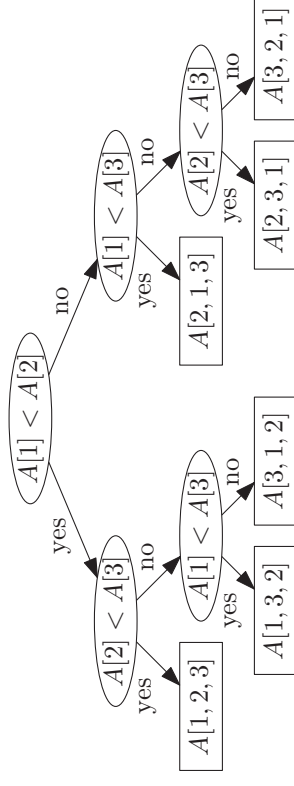
Each comparison is a “choice point” in the algorithm: the algorithm can do one thing if the comparison is true and another if false. So, the algorithm is like a binary tree...



30 / 33

## Complexity of the Sorting Problem

- ▶ There are  $n!$  possible output orderings of an input array of size  $n$ .
- ▶ There must be a leaf for each one; otherwise, the algorithm fails to sort.
  - ▶ For example, if leaf  $A[2, 3, 1]$  doesn't exist then the algorithm cannot sort [cat, ant, bee].



32 / 33

## Complexity of the Sorting Problem

- ▶ The number of leaves is at least  $n!$ .
- ▶ The height of the decision tree is at least  $\lceil \lg(n!) \rceil$ .
- ▶ The number of comparisons made *in the worst case* is at least  $\lceil \lg(n!) \rceil$ .
- ▶ This is true for **any comparison-based sorting algorithm**; therefore, the complexity of the sorting problem is  $\Omega(n \log n)$ .

