Unit #0: Introduction

CPSC 221: Basic Algorithms and Data Structures

Anthony Estey, Ed Knorr, and Mehrdad Oveisi-Fordoei[1]

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Annotated Slides from Ed's Class

[1]Thanks to Steve Wolfman for the content of most of these slides with additional material over the years 2004-2016 from Will Evans, Alan Hu, Ed Knorr, and Kim Voll.

Unit Outline

- Course logistics
- Course overview
- Fibonacci Fun
- Arrays
- Queues
- Stacks
- Deques

Course Information

Instructors

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Office hours

See `www.ugrad.cs.ubc.ca/~cs221`

Textbooks

- Susanna Epp's *Discrete Mathematics with Applications*
- Elliot Koffman and Paul Wolfgang's *Objects, Abstraction, Data Structures and Design Using C++*

Course Work

No late work; but we can exercise some discretion for medical cases, etc.

10%	Labs
15%	Programming projects (~ 3)
15%	Written homework (~ 3)
20%	Midterm exam
40%	Final exam

You must pass the final exam and the combination of labs/assignments in order to pass the course.

Collaboration

You may work in groups of two people on:

- Labs
- Programming projects
- Written homework

You may also collaborate with others as long as you follow the rules (see the website) and **acknowledge** their help on your assignment.

Don't violate the collaboration policy.

In other words, DON'T CHEAT!

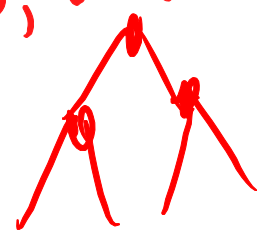
Course Mechanics

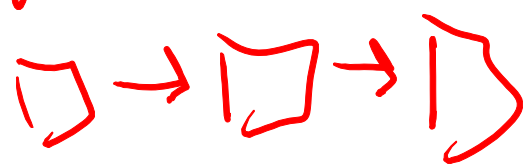
- Web page: `www.ugrad.cs.ubc.ca/~cs221`
- Piazza:
  `https://piazza.com/ubc.ca/winterterm22016/cpsc221/home`
- UBC Connect site: `www.connect.ubc.ca`
- Most/all labs are in ICCS 015 (check your own timetable)
  - Use the `Xshell` program on the lab machines to `ssh` into a undergrad Unix machine (e.g. `lulu.ugrad.cs.ubc.ca`)
- Programming projects will be graded on `UNIX/g++`

What is a Data Structure?

- a storage structure for data
- a way of storing, accessing, organizing, and manipulating data using a set of well-defined operations

e.g., array, vector, linked list (singly, doubly), hash set/table, hash map, dictionary, tree, graph, queue, stack, heap, etc.





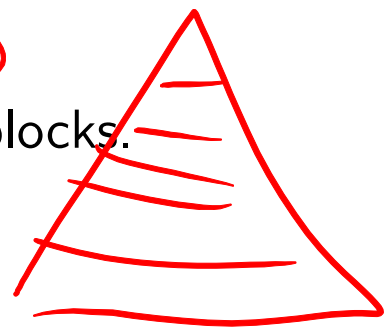
Observation

- All programs manipulate data.
  - Programs process, store, display, and gather data.
  - Data can be information, numbers, images, sound, etc.
- The programmer must decide how to store and manipulate data.
- This choice influences the program in many ways:
  - Execution speed ✓
  - Memory requirements ✓
  - Maintenance (debugging, extending, etc.) ✓

Goals of the Course

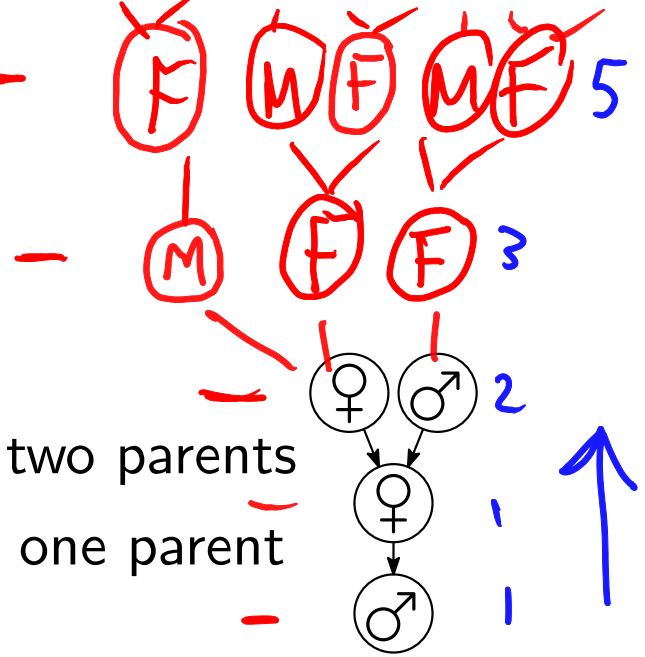
- Become familiar with some of the fundamental data structures and algorithms in computer science.
  - Learn when to use them.
- Improve your ability to solve problems abstractly.
  - Data structures and algorithms are the building blocks.
- Improve your ability to analyze algorithms.
  - Prove correctness.
  - Gauge, compare, and improve time and space complexity.
- Become modestly skilled with C++ and UNIX, but this is largely on your own!

e.g., Bloom's Taxonomy



for learning

Analysis Example: Fibonacci Numbers



Bee Ancestory:

1. Fertilized egg becomes a female bee with two parents
2. Unfertilized egg becomes a male bee with one parent

How many great-grandparents does a male bee have? great-great-grandparents? ...

Fibonacci numbers: 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, ...

First two numbers are 1; each succeeding number is the sum of the previous two numbers.

Rabbits:
Assume: 1. each pair of fertile rabbits produces a new pair of offspring every month
2. rabbits become fertile in their second month
3. old rabbits never die

Recursive Fibonacci

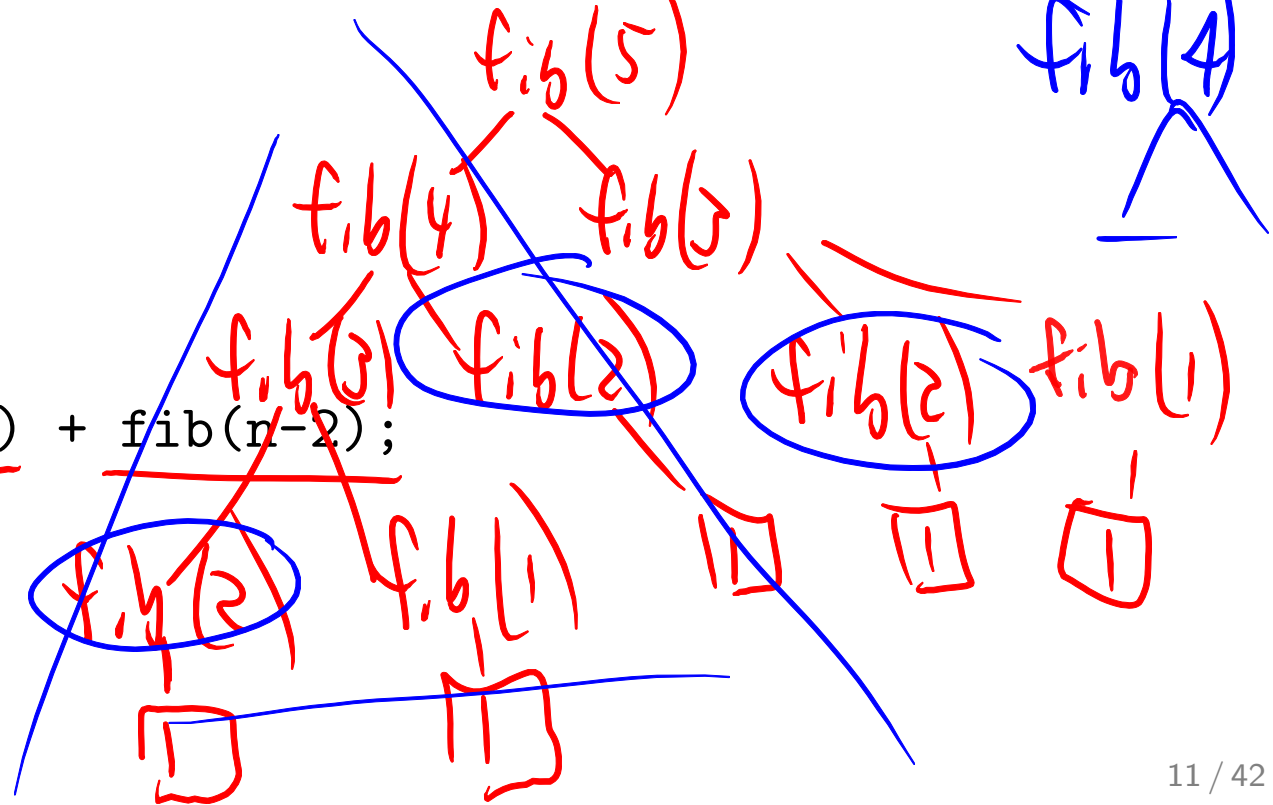
Problem: Calculate the nth Fibonacci number.

Recursive definition:

$$fib_n = \begin{cases} 1 & \text{if } n = 1 \\ 1 & \text{if } n = 2 \\ fib_{n-1} + fib_{n-2} & \text{if } n \geq 3 \end{cases}$$

C++ code:

```
int fib(int n) {
    if (n <= 2)
        return 1;
    return fib(n-1) + fib(n-2);
}
```



Too slow! Why?

Iterative Fibonacci

(2) Let's save our calculation for fib(k) in an array. Then:
- look it up when reusing fib(k) = "memoization"

(3) Dynamic Programming: Bottom-up

Idea: Use an array

```
int fib(int n) {
  int F[n+1];

  F[0]=0; F[1]=1; F[2]=1;
  for( int i=3; i<=n; ++i ) {
    F[i] = F[i-1] + F[i-2];
  }
  return F[n];
}
```

don't need array (last 2 numbers)

a, b, c

$a \leftarrow b$, $b \leftarrow c$, 3b

(We don't really need the array.)

Can we do better?

Fibonacci by Formula

Idea: Use a formula (a *closed form solution* to the recursive definition.)

$$fib_n = \frac{\varphi^n - (-\varphi)^{-n}}{\sqrt{5}}$$

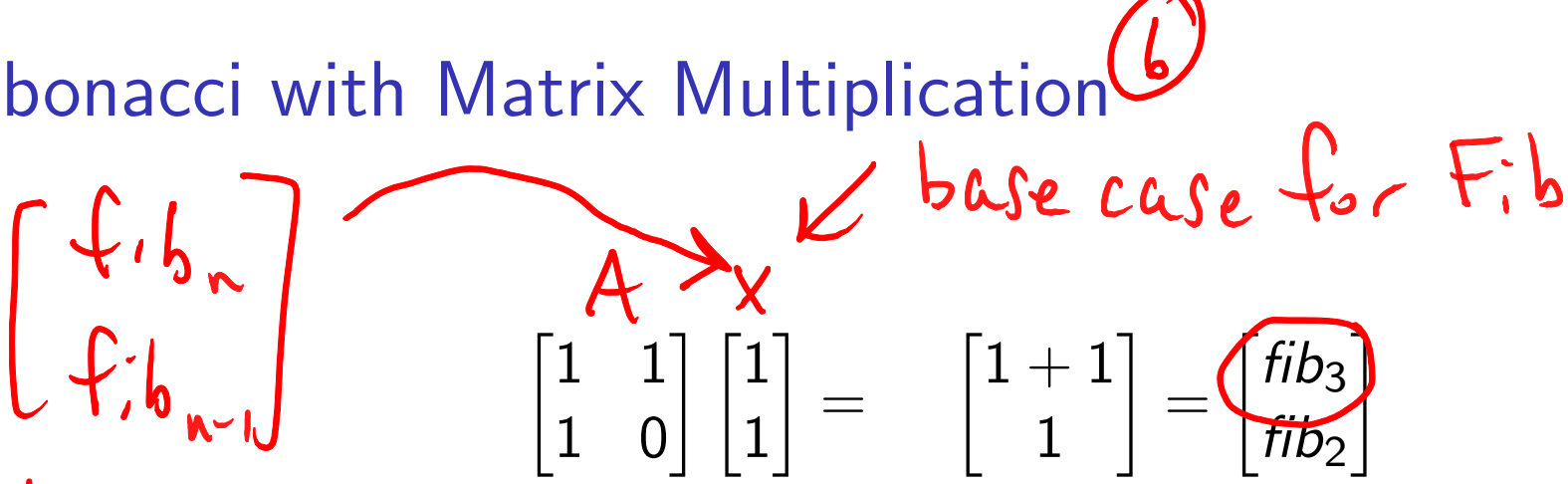
where $\varphi = (1 + \sqrt{5})/2 \approx 1.61803$.

```
#include <cmath>
int fib(int n) {
    double phi = (1 + sqrt(5))/2;
    return (pow(phi, n) - pow(-phi,-n))/sqrt(5);
}
```

Sadly, it's impossible to represent $\sqrt{5}$ exactly on a digital computer.

Can we do better?

Fibonacci with Matrix Multiplication



$$\begin{bmatrix} fib_n \\ fib_{n-1} \end{bmatrix} \quad \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 1+1 \\ 1 \end{bmatrix} = \begin{bmatrix} fib_3 \\ fib_2 \end{bmatrix}$$

matrix mult. is associative

$$(AA)x = A(Ax)$$

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} fib_4 \\ fib_3 \end{bmatrix}$$

but $AB \neq BA$ ∴ NOT COMMUTATIVE

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-2} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} fib_n \\ fib_{n-1} \end{bmatrix}$$

$j \times k \quad k \times m \Rightarrow j \times m$

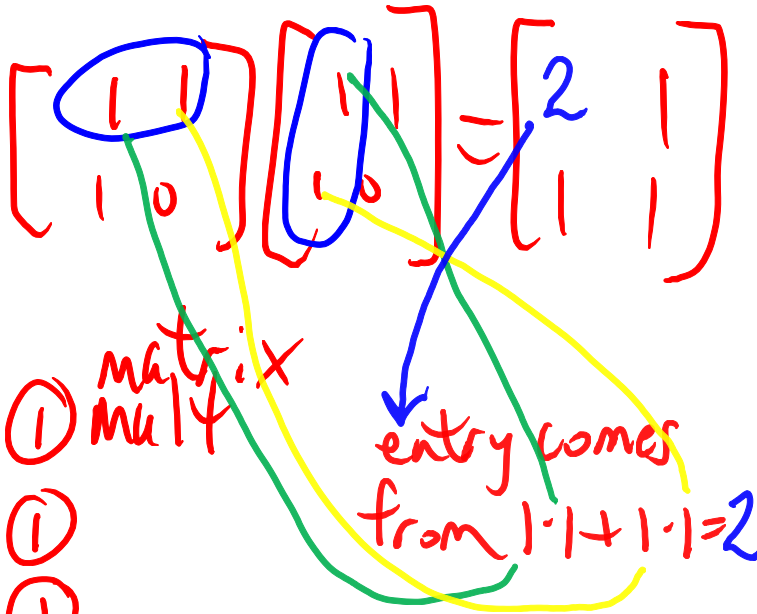
How do we calculate $\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-2}$?

result is a $j \times m$ matrix

Repeated Squaring

$$A \times A \times A \times A = (A \times A) \times (A \times A) = A^2 \times A^2 = A^4$$

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$



$$A \times A = A^2 \quad (1)$$
$$A^2 \times A^2 = A^4 \quad (1)$$
$$A^4 \times A^4 = A^8 \quad (1)$$
$$A^8 \times A^8 = A^{16} \quad (1)$$
$$A^{16} \times A^{16} = A^{32} \quad (1)$$
$$\vdots \quad A^{64} \quad (1)$$

$(100)_{10} = (1100100)_2$ 64 32 16 8 4 2 1

Example: $A^{100} = A^{64} \times A^{32} \times A^4$. 8 instead of 99 multiplications.
Generally, about $\log_2 n$ multiplications.

Is this better than iterative Fibonacci?

and also the little multiplications

Abstract Data Type

Also, See p. 24.

Abstract Data Type

Mathematical description of an object and the set of operations on the object

data

Example: **Dictionary ADT**

- Stores pairs of strings: (word, definition)
- Operations:
  - Insert(word,definition)
  - Delete(word)
  - Find(word)

data

How to implement it is another question.
- hash? - array? - tree?

Another Example: Array ADT

- Store things like integers, (pointers to) strings, etc.
- Operations:
  - Initialize an empty array that can hold n things.

    ```
    thing A[n];
    ```
  - Access (read or write) the ith thing in the array ($0 \leq i \leq n-1$).

    ```
    thing1 = A[i]; Read
    A[i] = thing2; Write
    ```

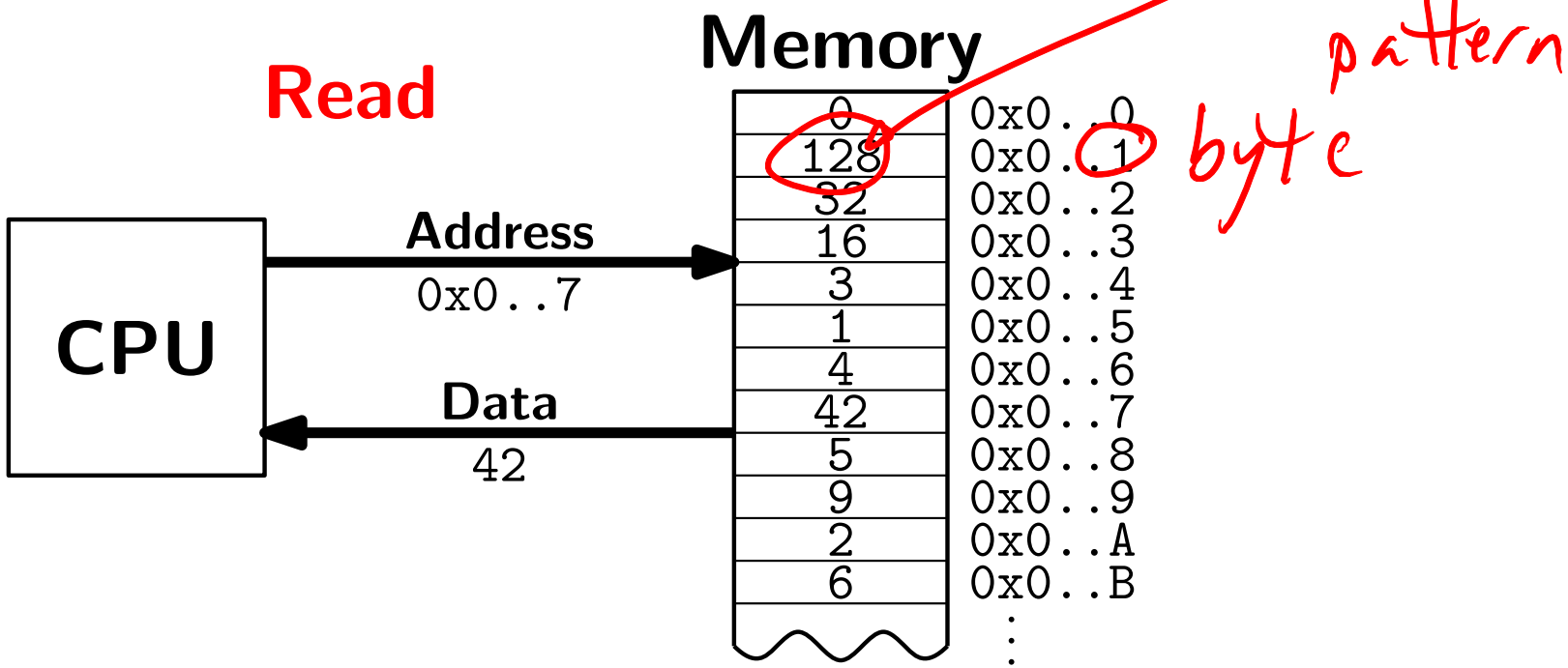
A's base address + i * sizeof object

Why Arrays?

- Computer memory is an array.

  Read: CPU provides address i,

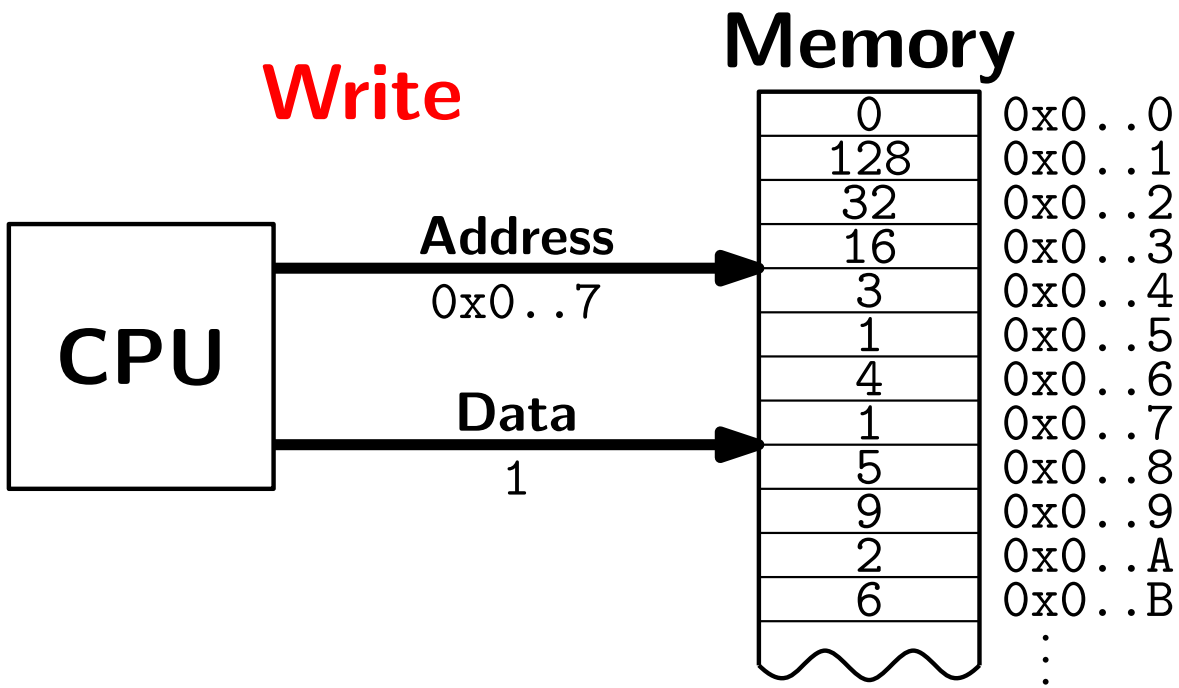
  memory unit returns the data stored at i.



Why Arrays?

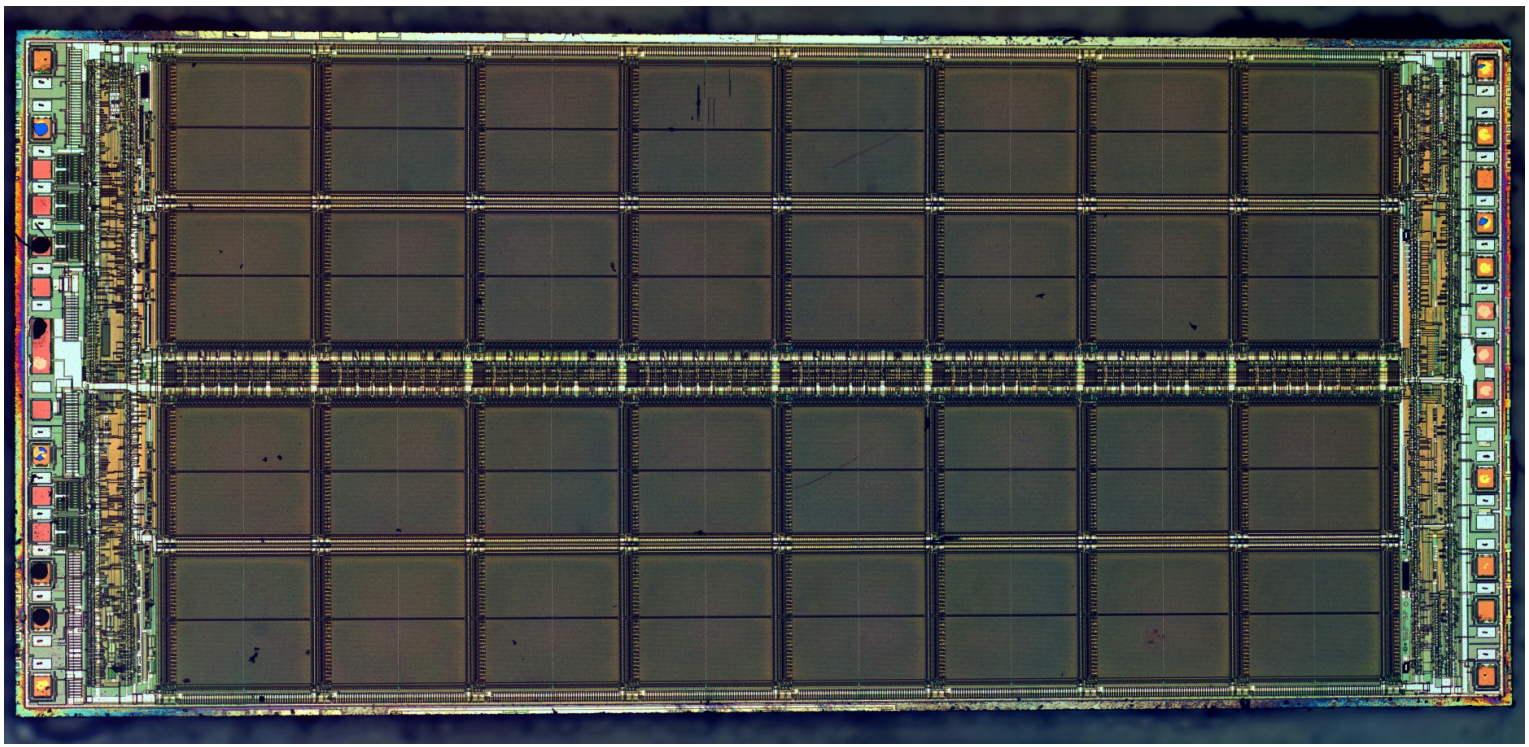
- Computer memory is an array.

  Write: CPU provides address i and data d, memory unit stores data d at i.



Why Arrays?

- Computer memory is an array. Every bit has a physical location.



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Why Arrays?

- Computer memory is an array.
- Simple and fast.
- Used in almost every program.
- Used to implement other data structures.

Array Limitations

- We need to know the size of the whole array when the array is created.

  Fix: Resizeable arrays.
  If the array fills up, allocate a new, bigger array and copy the old contents to the new array. Then, delete the old array.

- Indices are integers 0,1,2,...

  Fix: Hashing. This will give us greater flexibility.
  (more later)

How Would You Implement the Array ADT?

- get a block of memory of sufficiently large size (i.e., whatever the programmer needs)
- note the starting address of the block of memory
- note the size of each array element or object

Where do we find a given element?

$A[k] = A_{base\ address} + k * sizeof(object)$

How Would You Implement the Array ADT?

Arrays in C++

4 bytes

A is the base address location

To Create:

```
int A[100];
```

← declares but doesn't initialize

To Access:

```
for ( int i=0; i<100; i++ )
    A[i] = (i+1) * A[i-1];
```

careful!

- crash?

don't do this

dereferencing operator

```
A[5] = 777;
```

/* 6th location */

same

```
≡ *(A+5) = 777;
```

pointer arithmetic

How Would You Implement the Array ADT?

Arrays in C++

To Create:

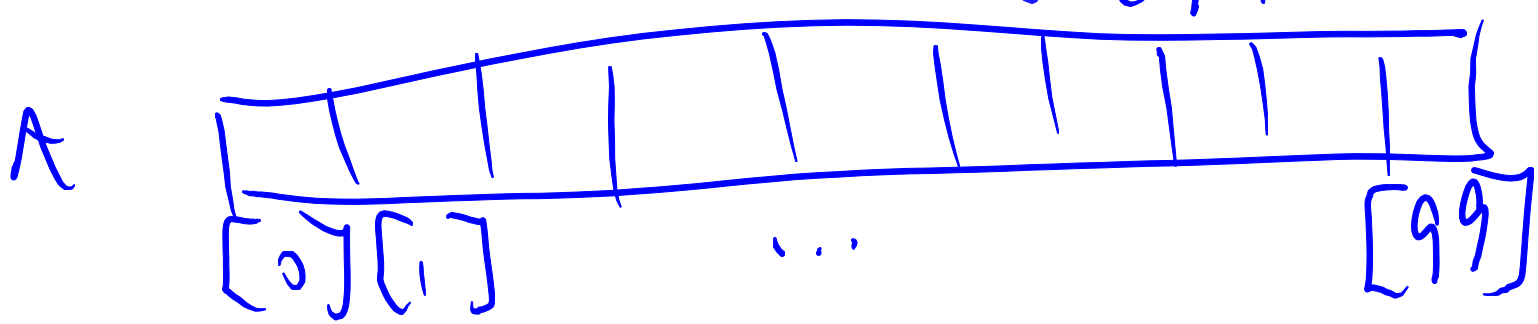
```
int A[100];
```

/* Assume A is initialized. */

To Access:

```
for ( int i=0; i<100; i++ )
     A[i] = (i+1) * A[i-1];
```

Warning No bounds checking! } C++



Data Structures as Algorithms; Abstract Data Types

Algorithm

Example: Sorting

- An algorithm is a high-level, language independent description of a step-by-step process for solving a problem.
- There may be multiple algorithms for solving a problem, and some may be more efficient than others.

many sorting algorithms exist — some better than others

Data Structure

- A data structure provides a way of storing and organizing data so that it can be manipulated by an ADT.

accessing

- An ADT describes *what* is stored, and it defines the *interface* (set of operations).

data

- An ADT is implemented by a data structure which specifies *how* the data is stored, and it provides *algorithms*
- An ADT may use different data structures in its implementation, for each operation.

Why So Many Data Structures?

Ideal Data Structure

- Fast, elegant, memory efficient

Trade-offs

- Time vs. space
- Performance vs. elegance
- Generality vs. simplicity
- One operation's performance vs. another's

Example: Data Structures for a Dictionary ADT

choices

- List
- Skip list
- Binary search tree
- AVL tree
- Splay tree
- B-tree
- Red-Black tree
- Hash table

...

Code Implementation for a Dictionary

Theory

- An abstract base class (interface) describes the ADT.
- Descendents implement the data structures for the ADT.
- Data structures can change without affecting the client code.

Practice

- Different implementations sometimes suggest different interfaces (generality vs. simplicity).
- The performance of a data structure may influence the form of the client code (time vs. space, one operation vs. another).

ADT Presentation Algorithm

1. Present an ADT.
2. Motivate it using some applications.
3. Repeat
   - 3.1 Develop a data structure for the ADT.
   - 3.2 Analyze its properties:
     - Efficiency
     - Correctness
     - Limitations
     - Ease of programming
4. Contrast the data structure's strengths and weaknesses.
   - Understand when to use each one.

Queue ADT

Queue operations

- create
- destroy
- enqueue = put in
- dequeue = take out
- is_empty

Queue property

If x is enqueued before y is enqueued, then x will be dequeued before y is dequeued.

FIFO: First In First Out

Applications of a Queue (Q)

- Holding jobs for a printer
- Storing packets on network routers
- Holding memory "freelists"
- Making wait lists fair
- Performing a breadth-first search (BFS)

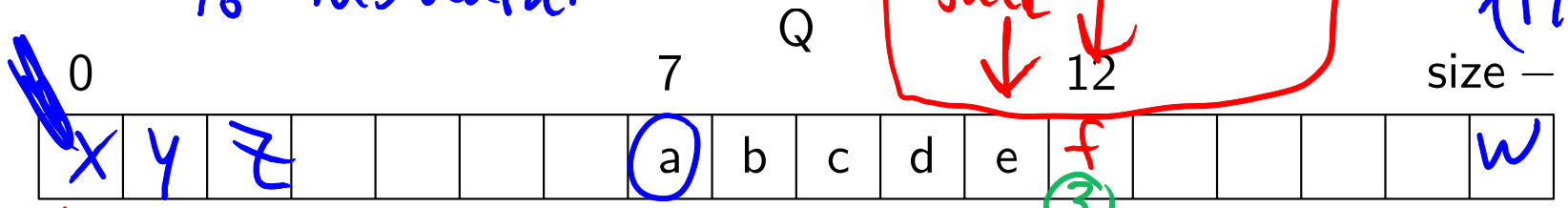
Abstract Q Example

enqueue R
enqueue O
dequeue
enqueue T
enqueue A
enqueue T
dequeue
dequeue
enqueue E
dequeue

In order, what letters are dequeued?

a. OATE
b. ROTA
c. OTAE
d. None of these, but it **can** be determined from just the ADT.
e. None of these, and it **cannot** be determined from just the ADT.

Circular Array Q Data Structure



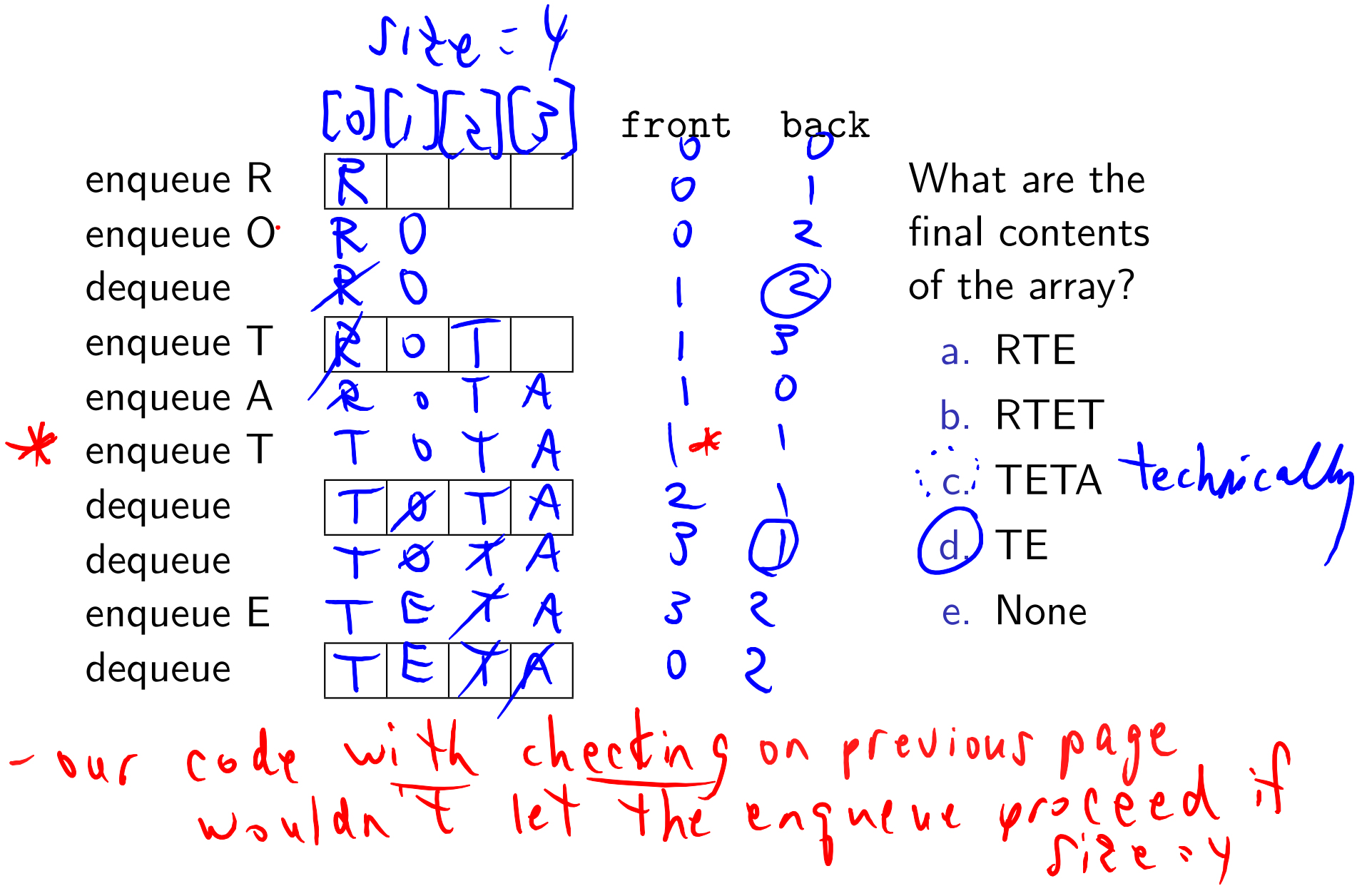

```
void enqueue(Object x) {
  Q[back] = x;
  back = (back + 1) % size;
}
```

```
bool is_empty() {
  return (front == back);
}
```

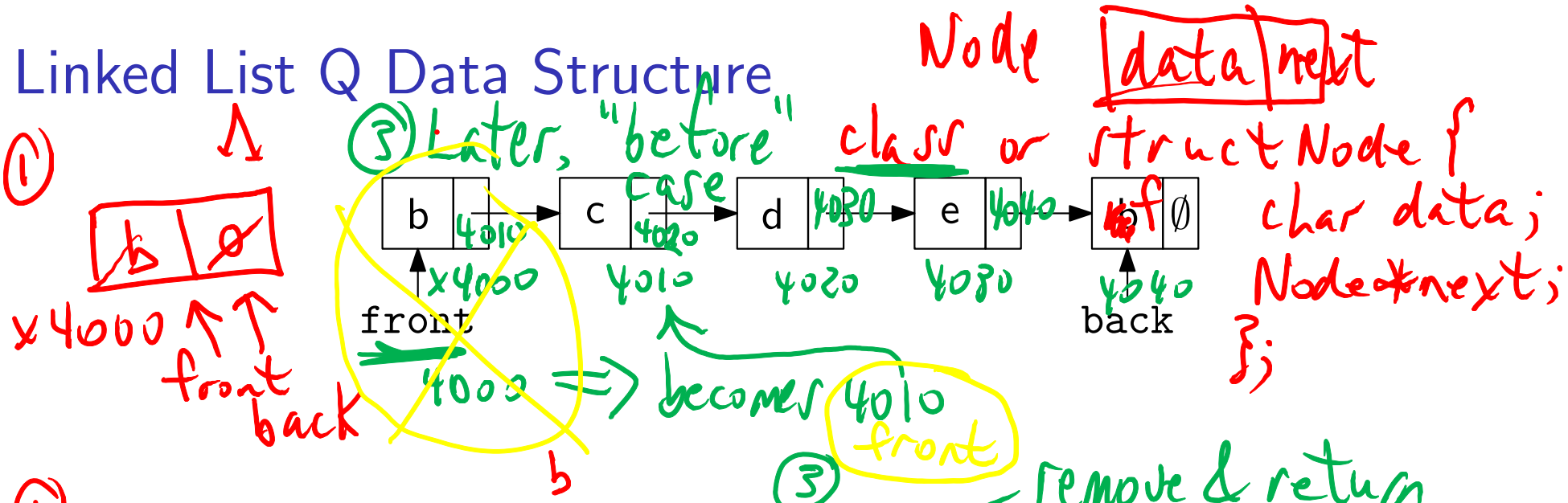
```
Object dequeue() {
  x = Q[front];
  front = (front + 1) % size;
  return x;
}
```

```
bool is_full() {
  return (front ==
    (back + 1) % size);
}
```

Circular Array Q Example



Linked List Q Data Structure




```
① void enqueue(Object x) {
    if (is_empty())
      front = back = new Node(x);
    else {
      back->next = new Node(x);
      back = back->next;
    }
  }
```

```
⓪ bool is_empty() {
    return (front == NULL);
  }
```

```
③ Object dequeue() {
    assert(!is_empty());
    Object ret = front->data;
    Node *temp = front;
    front = front->next;
    delete temp;
    return ret;
  }
```

else run out of memory

DIY memory management

start with front = NULL and back = NULL

Compare: Circular Array vs. Linked List

Ease of implementation

Generality

Speed

Memory use

modulus might be tricky, but often LL's are slightly harder for most programmers

- array may run out of space

- both fast, O(1)

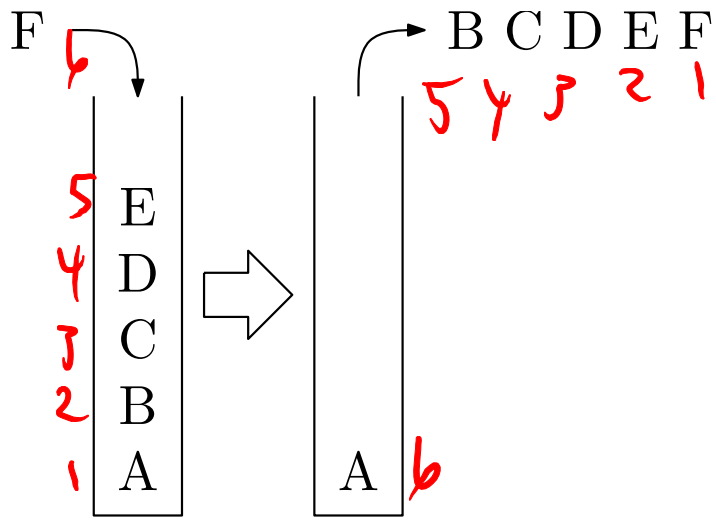
- depends on array's wasted space
- need a ptr for LL node
  e.g., extra 8 bytes/node

- array is better for locality
  (disk pages, RAM → cache)

Stack ADT

Stack operations

- create
- destroy
- push
- pop
- top
- is_empty



Stack property

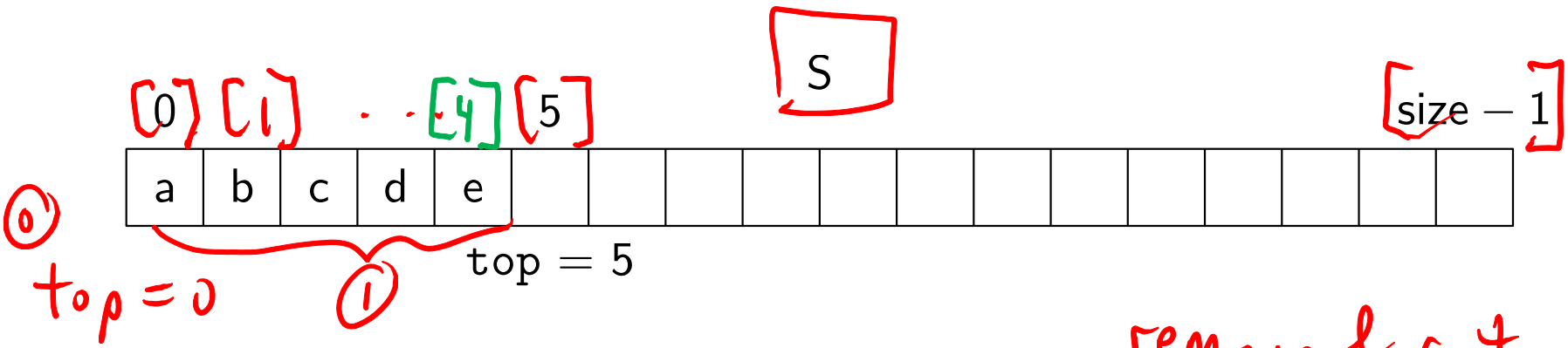
if x is pushed before y is pushed, then x will be popped after y is popped.

LIFO: Last In First Out

Stacks in Practice

- Implementing a function call stack
- Removing recursion
- Balancing symbols (e.g., parentheses)
- Evaluating Reverse Polish Notation (RPN)
- Performing a depth-first search (DFS)

Array Stack Data Structure

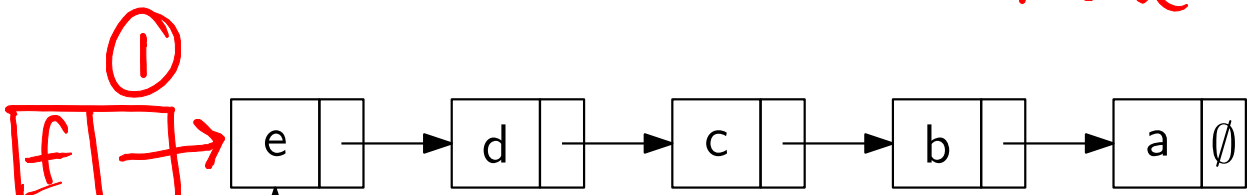



```
void push(Object x) {
  assert(!is_full());
  S[top] = x;
  top++;
}

Object top() {
  assert(!is_empty());
  return S[top-1];
}
```

```
Object pop() {
  assert(!is_empty());
  top--;
  return S[top];
}
bool is_empty() {
  return( top == 0 );
}
bool is_full() {
  return( top == size);
}
```

Linked List Stack Data Structure




```
void push(Object x) {
  Node *temp = top;
  top = new Node(x);
  top->next = temp;
}

Object top() {
  assert(!is_empty());
  return top->data;
}
```

```
Object pop() {
  assert(!is_empty());
  Object ret = top->data;
  Node *temp = top;
  top = top->next;
  delete temp; //free the node
  return ret;
}

bool is_empty() {
  return( top == NULL );
}
```

Deque ADT

Deque (Double-ended queue) operations

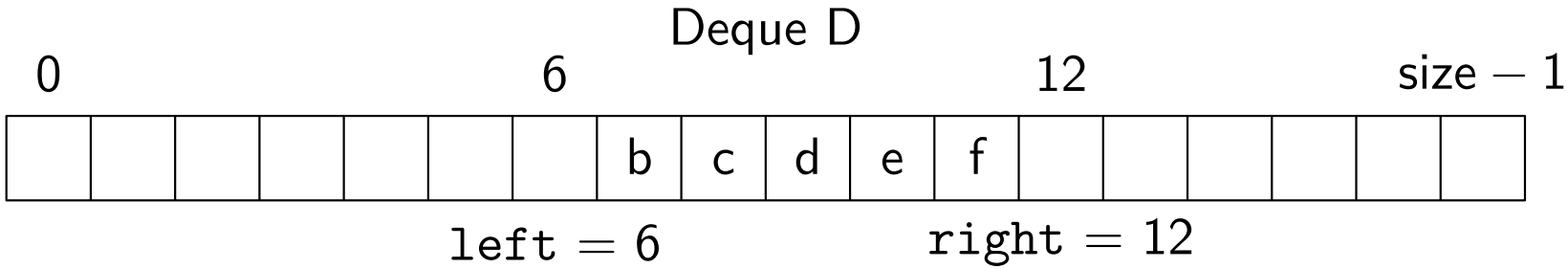
- create/destroy
- pushL/pushR
- popL/popR
- is_empty

Deque property

Deque maintains a list of items.
push/pop adds to/removes from front(L)/back(R) of list.

Circular Array Deque Data Structure

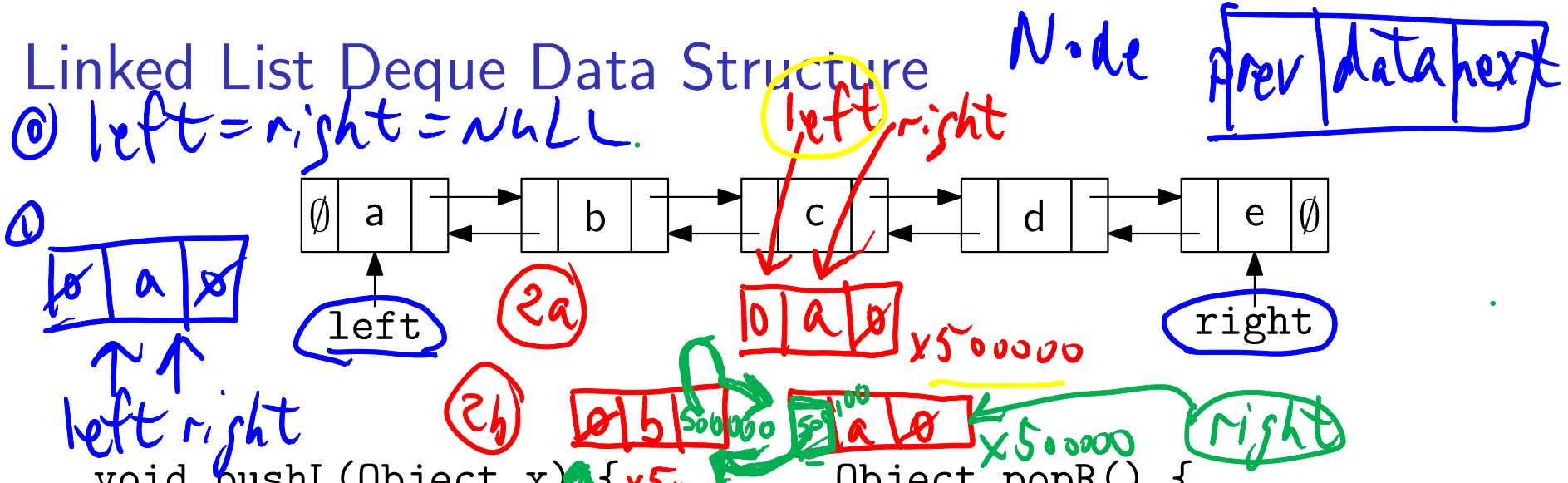
review on your own




```
void pushL(Object x) {
  assert(!is_full());
  D[left] = x;
  left = (left - 1) % size;
}
Object popR() {
  assert(!is_empty());
  right = (right - 1) % size;
  return D[right];
}
...
```

```
bool is_empty() {
  return( left ==
      (right - 1) % size);
}
bool is_full() {
  return( left ==
      (right + 1) % size);
}
```

Linked List Deque Data Structure




```
void pushL(Object x) {
  if ( is_empty() )
    left = right = new Node(x);
  else {
    left->prev = new Node(x);
    left->prev->next = left;
    left = left->prev;
  }
}

bool is_empty() {
   return left==NULL;
}
```

```
Object popR() {
  assert(!is_empty());
  Object ret = right->data;
  Node *temp = right;
  right = right->prev;
  if (right) right->next = NULL;
  else left = NULL;
  delete temp;
  return ret;
}
```

Data Structures You Should Already Know (Somewhat)

- Arrays
- Linked lists
- Trees
- Queues
- Stacks